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Note

No Antitwins in Minimal Imperfect Graphs

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It is customary to call vertices x and y twins if every vertex distinct from x and y is adjacent either to both of them or to neither of them. By analogy, we shall call vertices x and y antitwins if every vertex distinct from x and y is adjacent to precisely one of them. Lovász proved that no minimal imperfect graph has twins. The purpose of this note is to prove the analogous statement for antitwins. © 1988 Academic Press, Inc.

Claude Berge proposed to call a graph G *perfect* if for every induced subgraph H of G , the chromatic number of H equals the largest number of pairwise adjacent vertices in H . At the same time he conjectured that a graph is perfect if and only if its complement is perfect. Lovász [2] proved this conjecture which is known as the Perfect Graph Theorem.

It is customary to call vertices x and y *twins* if every vertex distinct from x and y is adjacent either to both of them or to neither of them. By analogy, we shall call x and y *antitwins* if every vertex distinct from x and y is adjacent to precisely one of them.

A graph G is *minimal imperfect* if G itself is imperfect but every proper induced subgraph of G is perfect. Lovász [2] proved that no minimal imperfect graph has twins. The purpose of this note is to prove the analogous statement for antitwins.

THEOREM. *No minimal imperfect graph contains antitwins.*

Proof. Assume the statement false: some minimal imperfect graph G contains antitwins u and v . Let A denote the set of all neighbours of u other than v , and let B denote the set of neighbours of v other than u .

As usual, a *clique* is a set of pairwise adjacent vertices, and a *stable set* is

a set of pairwise non-adjacent vertices; we let α and ω denote the largest size of a stable set and a clique, respectively, in G . We claim that

B contains a clique of size $\omega - 1$ that extends into no clique of size ω in $A \cup B$. (1)

To justify this claim, colour $G - v$ by ω colours and let S be the colour class that includes u . Since $G - S$ cannot be coloured by $\omega - 1$ colours, it contains a clique of size ω ; since $G - S - v$ is coloured by $\omega - 1$ colours, we must have $v \in C$. Hence $C - v$ is a clique in B of size $\omega - 1$. If a vertex x extends $C - v$ into a clique of size ω then $x \notin A$ (since $A \cap S = \emptyset$ and $G - S - v$ is coloured by $\omega - 1$ colours) and $x \notin B$ (since otherwise x would extend C into a clique of size $\omega + 1$). Thus (1) is justified.

The Perfect Graph Theorem guarantees that the complement of G is minimal imperfect; hence (1) implies that

A contains a stable set of size $\alpha - 1$ that extends into no stable set of size α in $A \cup B$. (2)

Now let C be the clique featured in (1) and let S be the stable set featured in (2); let x be a vertex in C that has the smallest number of neighbours in S . By (2), x has a neighbour z in S ; by (1), z is non-adjacent to some y in C . Since y has at least as many neighbours in S as x , it must have a neighbour w in S that is non-adjacent to x . Now u, z, x, y, w induce in G a chordless cycle. Since this cycle is imperfect, G is not minimal imperfect, a contradiction.

Chvátal *et al.* [1] call a graph G an (α, ω) -graph if it satisfies the following conditions:

- (i) G contains exactly $\alpha\omega + 1$ vertices.
- (ii) For every vertex w of G , the vertex-set of $G - w$ can be partitioned into α disjoint cliques of size ω and into ω disjoint stable sets of size α .
- (iii) Each vertex of G is included in precisely α stable sets of size α and in precisely ω cliques of size ω .
- (iv) Each stable set of size α is disjoint from precisely one clique of size ω and each clique of size ω is disjoint from precisely one stable set of size α .

Padberg [3] proved that every minimal imperfect graph is an (α, ω) -graph. However, there exist (α, ω) -graphs that contain antitwins. One such graph is featured in Fig. 1: the vertices 0 and 5 are antitwins.

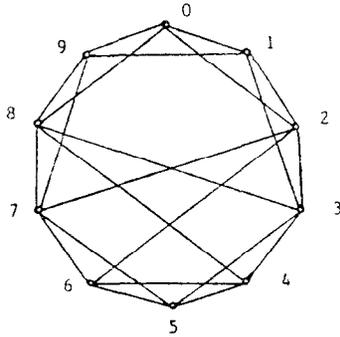


FIGURE 1

ACKNOWLEDGMENT

The author is indebted to Vašek Chvátal for the example in Fig. 1 and for many inspiring ideas.

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