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# A Flowchart Structure for Modification of a MODSIM Process Model

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**A FLOWCHART STRUCTURE FOR  
MODIFICATION OF A MODSIM PROCESS MODEL**

by

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B.Tech. July 1996, Jawaharlal Nehru Technological University, Hyderabad, India

A Thesis Submitted to the Faculty of  
Old Dominion University in Partial Fulfillment of the  
Requirement for the Degree of


MASTER OF SCIENCE

COMPUTER ENGINEERING

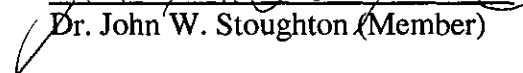
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August 2000

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# **ABSTRACT**

## **A FLOWCHART STRUCTURE FOR MODIFICATION OF A MODSIM PROCESS MODEL**

Murali K Adatrao  
Old Dominion University, 1997  
Director: Dr. James F. Leathrum, Jr.

There are many software processes and software development models that support the development of a software model prior to implementation. However, more often than not, these practices are not followed resulting in a lack of documentation of the intended functionality of the software. As a result, software often becomes a black box for later developers. Even a simple bug fix can turn into an exhaustive task for the developers, as they must attempt to infer the intended system behavior. New designers cannot make any changes to the software or extend the software behavior, as the underlying model within the software itself is unclear. An understanding of the underlying model helps both the designers and programmers to do their job quickly and effectively.

This thesis proposes a model to describe an existing undocumented software system developed in MODSIM III to facilitate future development efforts. The model provides a way to understand the underlying model of the software system by first documenting the software using flowchart constructs. The model also ensures designers and programmers make respective changes in the flowcharts before implementing the proposed changes. This helps document the recent architecture of the software process in the form of flowcharts. The model is described in this thesis and demonstrated on the port simulation PORTSIM.

This thesis is dedicated to my mother.

## **ACKNOWLEDGMENTS**

I would like to thank my committee members for their patience in guiding my research and editing of this manuscript. The untiring efforts of my major advisor who helped me in all stages of this thesis deserve special recognition.

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# SECTION 1

## INTRODUCTION

### 1.1 Overview

This thesis proposes a systematic approach that can be followed during the modification and maintenance of an undocumented, complex software system and is applied to simulations developed in MODSIM III [13]. The process is demonstrated on a port simulation, PORTSIM [12]. The objectives of this process are to abstract documentation of the software in the form of process flow in order to understand the underlying model and use it while developing the software. The documentation should provide required information for the designers and the programmers for modeling and debugging. Here debugging represents removing modeling errors but not coding errors. The document should, however, help to isolate code containing coding errors. The model is specific to MODSIM in order to capture the language constructs inherent in MODSIM, but new models can be developed for other languages in a similar manner.

The main problem in working on undocumented software is that the tasks it performs are identifiable, but the process of accomplishing those tasks is not known. In order to make any changes in the software, it is necessary to understand the way in which the software is accomplishing its tasks.

In software development models, the software system is first analyzed and designed from the specifications before going into implementation phase. So by the time the software

reaches the implementation phase the designers have a good understanding of what is happening in the system. Also, the programmers, if they are not the designers, can learn about the system by referring to design documents. The design documents are also helpful during debugging. By understanding the kind of bug the software is suffering from, one can make a good guess on the part of the code that is causing the problem. Hence there is a great need to document the undocumented software to continue to develop the software effectively.

In most cases, designers design the system in accordance to the programming environment that will be used to develop the system. Knowledge of the programming environment enables one to adopt design strategies relative to the programming environment. The way one documents the software should reflect the environment of the programming language used in developing the software. By doing this, any changes proposed by the designers through the documentation can be easily implemented by the programmers. One should take care to make the software's document represent its underlying model instead of the sequence of executable statements. If possible the system should be divided into different logical groups while documenting. This helps to track a bug during debugging once we understand the type of bug and also to speed up the process of modification by only concentrating on the enclosing logical unit while making changes to the system.

## **1.2 How the Objectives are Achieved**

In the proposed model, the above mentioned requirements are met by first developing a set of pictorial representations of MODSIM III constructs and using them in the flowcharts to represent the program flow within the software. By doing this, the flowcharts themselves provide the software's programming environment to the designers. The programmers can easily convert any modifications in the flowcharts into MODSIM III code as long as the designers design their modification within the boundary of the pictorial set. A logical unit is constructed by referring to the related flowcharts from within a flowchart.

## **1.3 Section Overview**

Section 2 provides a brief introduction to the programming language, MODSIM III, used to develop the software, PORTSIM, and to the software itself. It also provides an overview of well-known software models and discusses whether PORTSIM can follow any of the models. In the conclusion, the importance of flowcharts in documenting PORTSIM is discussed.

Section 3 discusses in-depth the important MODSIM III language constructs and how they can be represented pictorially. It later discusses how the pictorial representation of MODSIM III constructs can be used in developing flow diagrams representing the program flow within PORTSIM and how the resulting flow diagrams can be used in the development process of the PORTSIM.

Section 4 provides examples on how to use the flowcharts for debugging and for modeling the code to make improvements.

Section 5 gives the summary of the benefits and the drawbacks and possible solutions of the proposed model.

## **SECTION 2**

### **BACKGROUND**

This section provides an introduction to MODSIM III programming language, PORTSIM, a simulation tool developed to simulate military cargo flow through commercial ports, and to the software engineering models that are used to develop software systems. Finally, this section discusses how the present models are inefficient to use in the development cycle of PORTSIM.

#### **2.1 MODSIM III**

MODSIM III is a modular, object-oriented, strongly typed, block structured, discrete-event simulation language [5]. Every MODSIM III program contains a main module. Programs can be divided into modules with each module supporting some particular functionality. Any module can be compiled separately, making the maintenance of the code easy. Each module may contain a number of objects. An object is an encapsulation of data items which describe the state of the object and methods which describe the interface of the object. MODSIM III also supports inheritance and polymorphism. Every expression and method parameter is checked at compilation time for type consistency. This strongly typed feature helps to identify errors at compilation time rather than at run time. The scope and visibility of variables is restricted to the block in which they are declared. The library modules and the language features provide all the capabilities for developing discrete-event simulation models. MODSIM III demands that simulation models be developed in terms of process. The process is capable of performing multiple concurrent activities. The activities can be modeled to operate independently or in a

synchronized fashion. This process model helps to code a related group of activities in one routine.

## **2.2 PORTSIM**

PORTSIM is a database-driven, object-oriented, computer simulation [12]. The software is being developed in MODSIM III. PORTSIM is useful for force deployment through a given commercial port [14]. It is useful in estimating the time of embarkation and disembarkation of a force, given the parameters of the port and the force [11]. The purpose of this software is to study the effects of initial parameter values on the end results. The software can be used to simulate either port of embarkation (POE) or port of disembarkation (POD) operations. In either mode, we can either create a new scenario in which PORTSIM relies heavily on the underlying database to get required data or load an existing scenario for which PORTSIM reads data from the scenario file.

## **2.3 Software Models**

There are many systematic development models that are being used for developing a software model. The first model that was used to develop a software model is the Waterfall model. The various stages involved in the Waterfall model is shown in Fig. 1 [1]. Though the Waterfall model does have disadvantages, it showed that a systematic approach in developing the software does make improvement in the quality of the code and the time taken to develop the code. Other significant models which were developed later are:

1. Spiral model [1], [2].
2. Booch method.
3. Object Modeling Technique (OMT) [10].
4. Object-Oriented Software Engineering (OOSE) [9].
5. Unified Modeling Language (UML) [7], [8].

UML is the result of combining the Booch, OMT and OOSE methods [3].

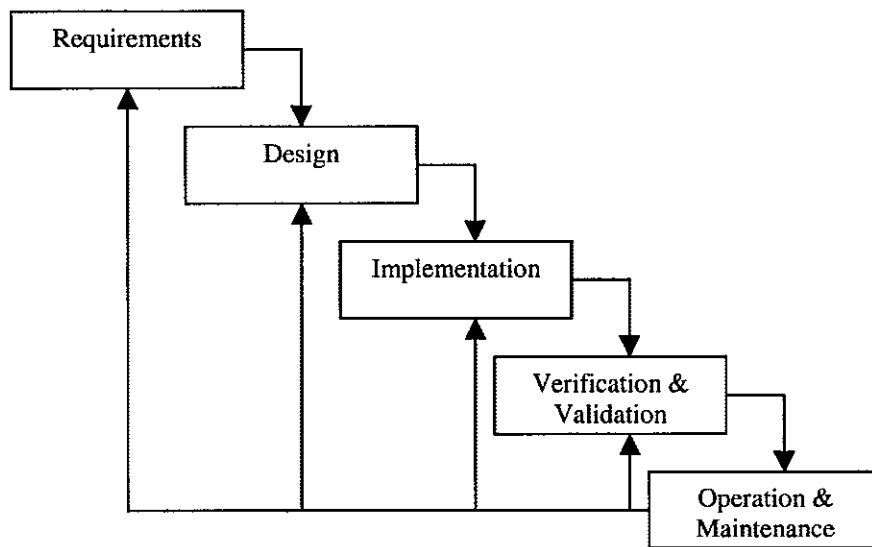


Fig. 1 Waterfall model



### 2.3.1 Introduction to the Booch Model

This section discusses the Booch model as an example which uses Object-Oriented Analysis and Object-Oriented Design [4]. The Booch model divides the process of software development into two processes, namely:

1. Macro process.
2. Micro process.

*Macro process:* Booch's Macro development process [4] is shown in Fig. 2, which describes the overall process of software development.

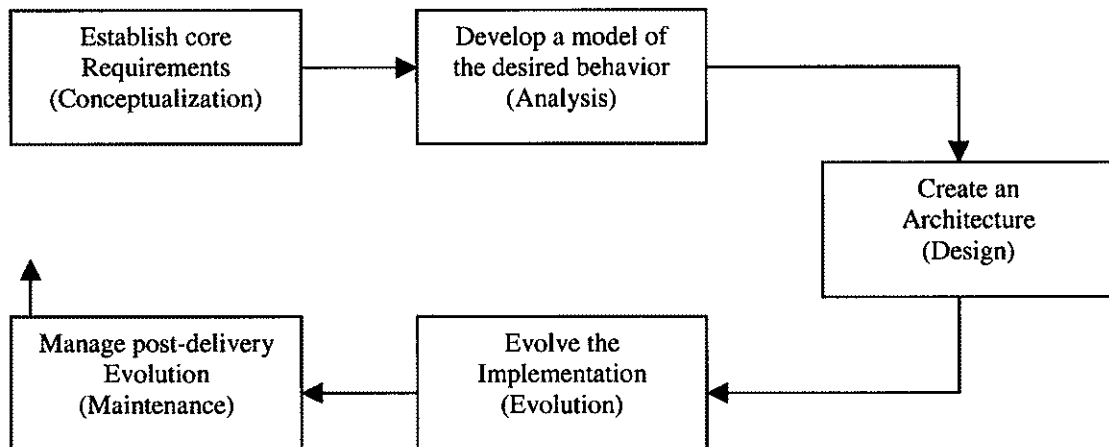


Fig. 2 Booch's Macro development process

*Micro process:* Booch's Micro development process is shown in Fig. 3 [4], which is used to carry out iteratively the activities in the Macro process.

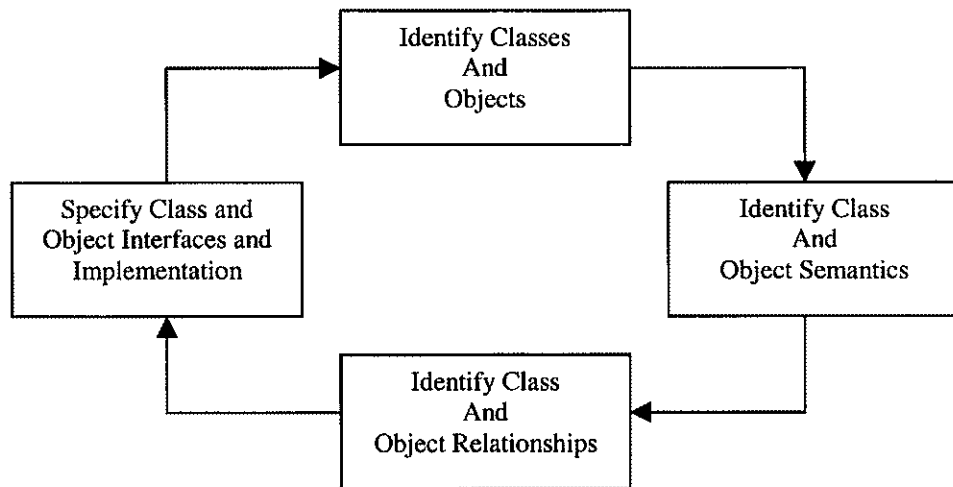


Fig. 3 Booch's Micro development process

Actually, Booch's Macro process can be represented as shown in Fig. 4, where each small circle represents a micro process. For each activity in macro process emphasis will be given on different activity in the corresponding micro process.

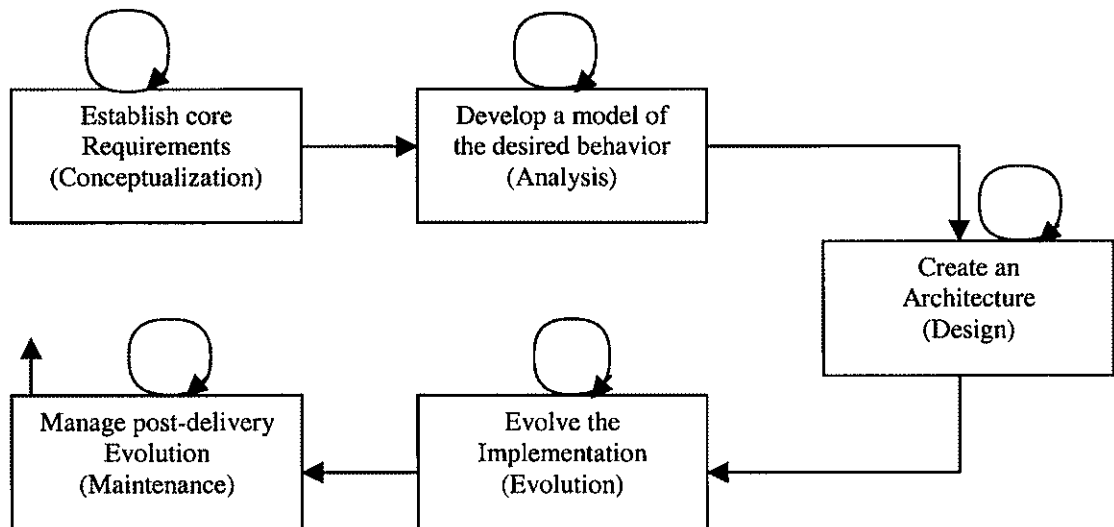


Fig. 4 Alternate representation of Booch's Macro development process

In both the Booch method and the UML, the class diagram is used to specify static relationships between classes in the design phase. Class diagrams do not represent dynamic relationships such as when objects are created or invoke services of other objects.

### *2.3.2 Software development Models and PORTSIM*

From the above discussion, we can see that the present software development models are useful for developing software from software's overall specification to final product. The models actually start from a higher level model of the software and go down to implementation or a lower level model in the later stages. Also, in any given intermediate state of development the models seem to have a proper documentation for the software under development. The absence of any such documentation for PORTSIM and the software structure suggest that PORTSIM did not follow any of the development methods mentioned above. So, to successfully use any of the development methods for PORTSIM, we need to document the software by developing the higher level model of the software from its implementation. As the software is already in an implementation stage, deciding on which development method to use is difficult. Also, the higher level model developed from the software may not be consistent with the higher level model that may be developed if the chosen development method is actually used from the beginning. So it is not wise to think of using any of the above mentioned development methods in the modeling of PORTSIM. The only way left to document PORTSIM is to represent the present architecture in such a way that it will be useful to identify the

underlying model, to identify the dynamic relationships between objects and also to make future enhancements easily and successfully.

## **2.4 Flowcharting**

Flowcharts help to represent the underlying model and the dynamic relationships between objects, thus eliminating any need to develop a higher level model to make future enhancements. Since PORTSIM is implemented in MODSIM III, it is necessary to develop flowchart constructs to represent MODSIM III language constructs. By doing this, the resulting flowchart representation of PORTSIM helps both designers and programmers. Designers can now make changes in the flowcharts using the standardized representation and programmers can convert the changes into MODSIM III constructs. That is, flowcharts developed with a one-to-one relation between the flowchart constructs and programming language constructs act like a common interface between designers and programmers.

## **SECTION 3**

### **PROCESS FOR MODIFICATION OF EXISTING CODE USING FLOW DIAGRAMS**

The modification process of existing software can be classified into bug fixes in the software and modified behavior. Though both of these processes produce new problems in the code if not done properly, the latter one is more prone to this effect, as care should be taken to leave the functionality of the rest of the code unchanged. The first thing to do when modifying existing code is to classify the purpose of modification into the modification process for 1) bug fixing or for 2) better performance. This section presents the process of modeling existing MODSIM software for modification.

#### **3.1 MODSIM III Constructs and Flow Diagram Constructs**

For both modification processes, we need to understand the program flow within the software code. An appropriate way to represent the program flow is by means of flow diagrams. The complexity of the flow diagrams depends on the complexity of the programming language in which the software is developed. Therefore, it is always a good idea to understand the capabilities of the host language and adopt a notation that reflects its capabilities as closely as possible. This helps to translate code into flow diagrams and flow diagrams into code with consistency. PORTSIM, which will be used to demonstrate the modification process, was developed using MODSIM III. The following section discusses the MODSIM III constructs that need to be represented properly to reflect the functionality of the code.

### 3.1.1 MODSIM III constructs

MODSIM III supports process-oriented simulation instead of event-oriented simulation [5]. In case of an event-oriented simulation, passage of time is handled by scheduling the next event for the object currently being processed. This type of timing model is well suited to small models, but in the case of large models, where there is a lot of unrelated event routines in-between, following a single object's behavior is more difficult. MODSIM III simplifies larger models by allowing many aspects of an object's behavior in a model to be described in one method that allows for the passage of time at one or more points in its code. To support this concept, MODSIM utilizes several constructs to handle timing issues between concurrent processes, which are not present in all languages. There are also libraries to support capabilities like a Graphical User Interface (GUI) and methods like **NEW** and **DISPOSE** for creating and destroying objects for which it is undesirable to model the internal components in detail. The following section discusses the constructs that MODSIM III contains to build simulation models.

#### 3.1.1.1 MODSIM III constructs from the programmer's perspective

##### DEFINITIONS:

**Activity:** What occurs in the model as time elapses. No code is executed

**Event:** A point in time at which the state of the model changes in some way. Code is executed, but time does not elapse.

In MODSIM III, a programmer has complete control of the simulation process with the help of ASK, TELL and WAITFOR methods [5]. A call to an ASK method is similar to that of a procedure call, i.e., when an ASK method is invoked, the calling code waits until the method is executed completely. Simulation time is not allowed to elapse in an ASK method. So an ASK method symbolizes the occurrence of an **event**. A call to a TELL method is similar to that of an asynchronous call, i.e., when a TELL method is invoked, the calling code does not wait for the completion of the method. Simulation time is allowed to elapse in a TELL method using a WAIT statement, so a Wait statement symbolizes the occurrence of an **activity**. A TELL method can also be scheduled to execute at some time in the future. A call to a WAITFOR method is similar to that of a procedure call with the ability to elapse simulation time, i.e., WAITFOR method has both the required properties of TELL and ASK methods. One important use of TELL and WAITFOR methods is that a programmer can achieve arbitrary synchronization by the use of TRIGGER objects within them. Triggers are used along with WAIT statements. TRIGGERS can be used when some processes need to wait until a specific condition occurs. A TRIGGER object can have any number of processes waiting for it. The processes will continue to wait till the trigger object's TRIGGER method is invoked by some other process.

Another important programming construct that needs to be represented carefully is the selection construct. This construct is important because the program control may follow any of the associated branches, depending on the result of the selection condition at a given time. So if an activity, event or logical process occurs as a branch in the selection

construct, the program control may or may not go through that branch depending on the result of the selection condition. Hence the branches could be alternate activities, alternate events or alternate logical processes respectively.

### 3.1.1.2 MODSIM III constructs from the GUI developer's perspective

SIMGRAPHICS II is the MODSIM III'S graphics package [6]. Since MODSIM III is an Object Oriented language, the graphical interface is also implemented using objects. Many basic objects have already been provided. The most important ones are shown in TABLE 1 below [6].

TABLE 1

Basic objects and their properties

OBJECT NAME	PROPERTIES
WindowObj	Standard system window which acts as a container for all graphical objects.
ImageObj	Basic object used for static icons and backgrounds.
DynImageObj	Basic graphic object used for animation.
DialogBoxObj	Receives various types of input from the user. Controls, such as buttons, check boxes, list boxes etc. can be part of it.
PaletteObj	Receives input from two-state palette buttons.
MenuBarObj	Receives simple menu selections.



An instance of any of these objects can be made visible by calling their respective DRAW methods and can be erased by calling their respective ERASE methods. In SIMGRAPHICS II, a developer can have complete control on the behavior of graphical objects with the help of the BeSelected, AcceptInput and BeClosed methods. The BeSelected method is automatically invoked when a graphical object is clicked on. This method can be overridden to receive asynchronous input. Within a BeSelected method of a graphical object, the LastPicked field can be checked to determine which control was last clicked on. Depending on the control, a programmer can model the behavior of the graphical object. AcceptInput method is used to retrieve synchronous input from a graphical object. AcceptInput returns the contents of the LastPicked field after the selection has been made. However, if the graphical instance is erased or disposed, Acceptinput returns NILOBJ [6]. BeClosed method is automatically invoked whenever a user chooses to close a graphical object. This method can be overridden to change its default behavior associated with a graphical object.

Of all the three methods, the behavior of the Acceptinput method cannot be changed. This very fact makes the representation of data flow difficult if a dialog box uses both Acceptinput and BeSelected methods in its life cycle. We cannot tell whether the selected input is returned to the AcceptInput method first or to the BeSelected method. No matter which of the two methods receive the selected input first, by the time the statement after a call to AcceptInput method is executed we can say that both the BeSelected and AcceptInput methods are returned. So the point after a call to AcceptInput method can be represented as the point of return for both AcceptInput and BeSelected methods. So in






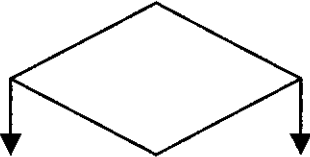
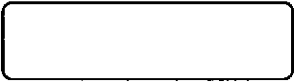
any situations where the actual implementation is hidden in the library routines, we should look for the first point at which we are absolutely sure about the cumulative result, i.e., we need to concentrate on final result rather than the actual sequence of data flow. These issues complicate the modeling process.

#### 3.1.1.3 MODSIM III constructs from the designer's perspective

From the designer's point of view, a call to any type of method is a way of passing messages from one object to another object or to itself. Objects respond to messages by executing the associated code, which may in turn send messages. In some cases this message sending will be so nested that it is a good idea to subdivide the code associated with a message into logical processes where each logical process represents a set of closely related messages. Closely related messages are messages that are used together to accomplish a single target. In fact, any involved message can also be treated like a logical process, so any system can be represented as sequence of logical processes that are to be completed to get the results. The following section shows the notations for various MODSIM III constructs.

### 3.1.2 Flow Diagram constructs for MODSIM III constructs

This section shows the flowchart constructs that can be used to represent program flow within the software developed in MODSIM III.

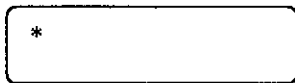
Flowchart Construct	Purpose
	Begin/End of logical process.
	Represents an event within a logical process.
	Represents an activity within a logical process.
	Represents a logical process within a logical process.
	Represents a logical process containing an activity.
	Represents a condition within a logical process.
	Represents an alternate event as a result of condition or as a result of selection.



Represents an alternate activity as a result of condition or as a result of selection.



Represents the point of return of multiple methods irrespective of order of completion of the methods.



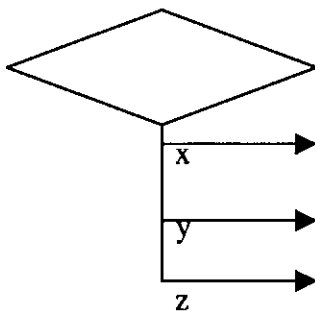
Represents an alternate logical process within a logical process.



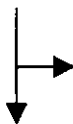
Represents an alternate logical process containing an activity.



Represents the direction of flow of program control.



Represents that the program control will follow any of the directed paths depending on the outcome of the selection statement.



Represents an asynchronous call and hence the program control continues to flow down without waiting for the completion of process started by the branched arrow.

**Flowchart Construct****Purpose**

Represents a terminator.

### 3.1.3 Flow Diagram Examples for MODSIM III constructs

This section shows how the above notations can be used to represent different constructs in MODSIM III.

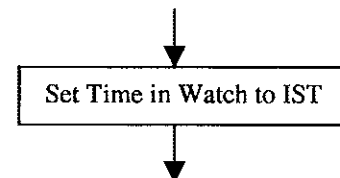
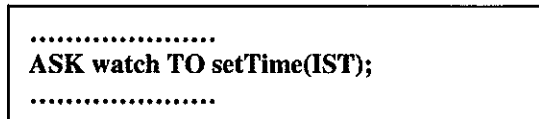


Fig. 5 Example 1: An Event

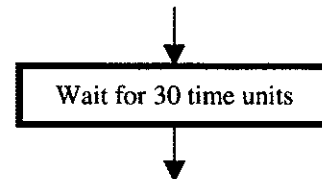
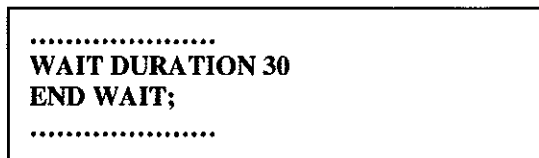


Fig. 6 Example 2: An Activity

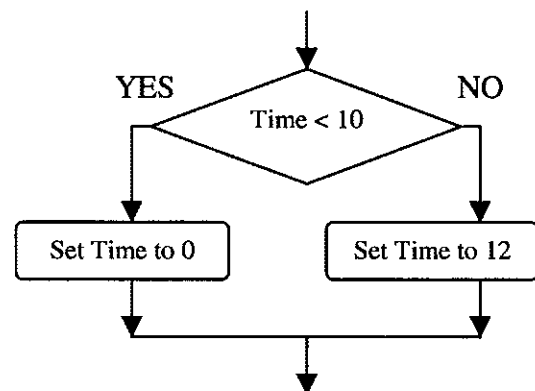
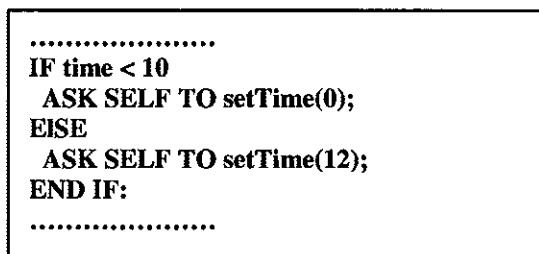


Fig. 7 Example 3: A Condition

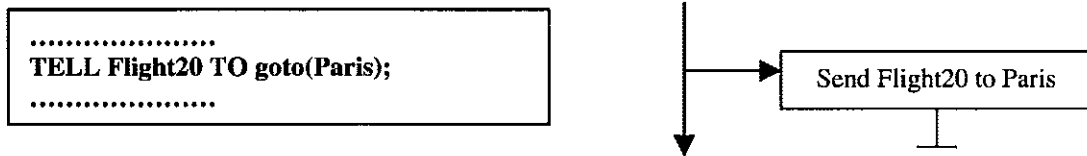


Fig. 8 Example 4: A Asynchronous call

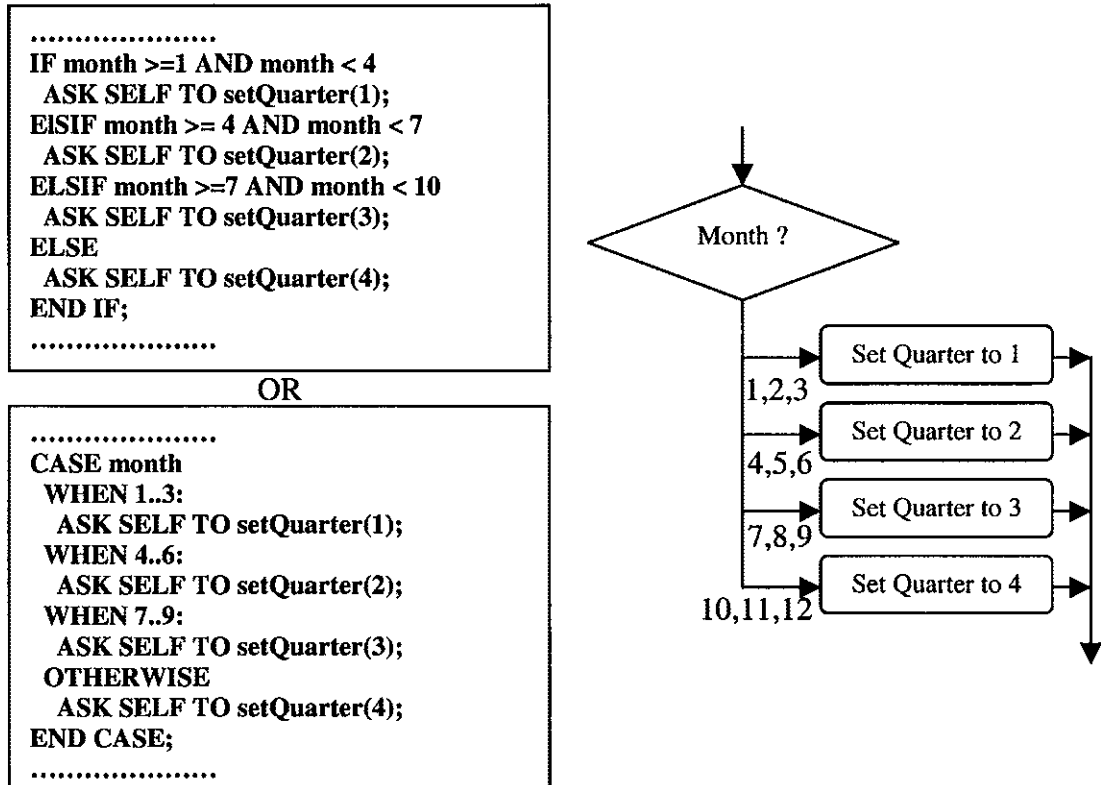


Fig. 9 Example 5: A Selection

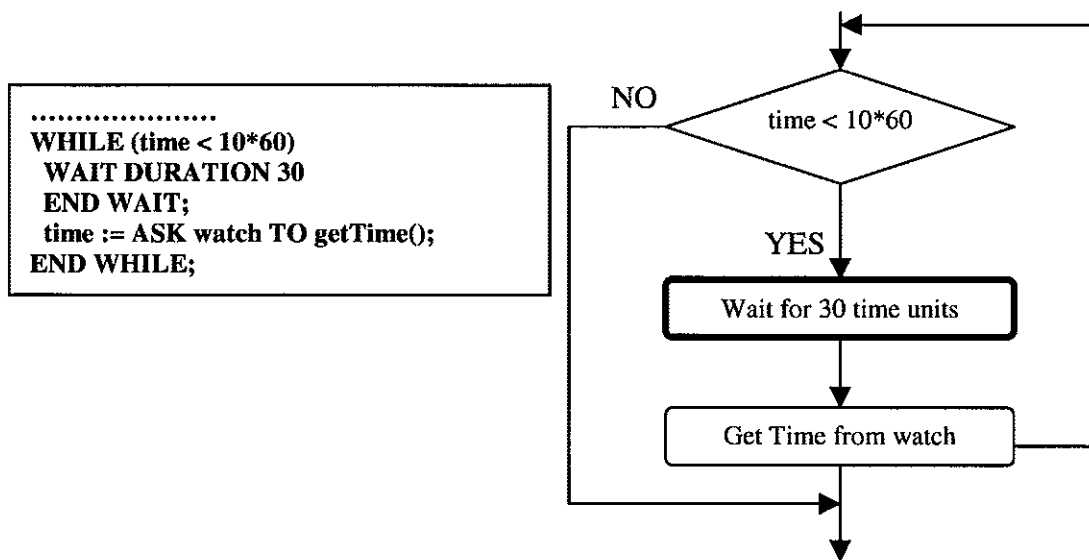


Fig. 10 Example 6: Loop with an activity and an event

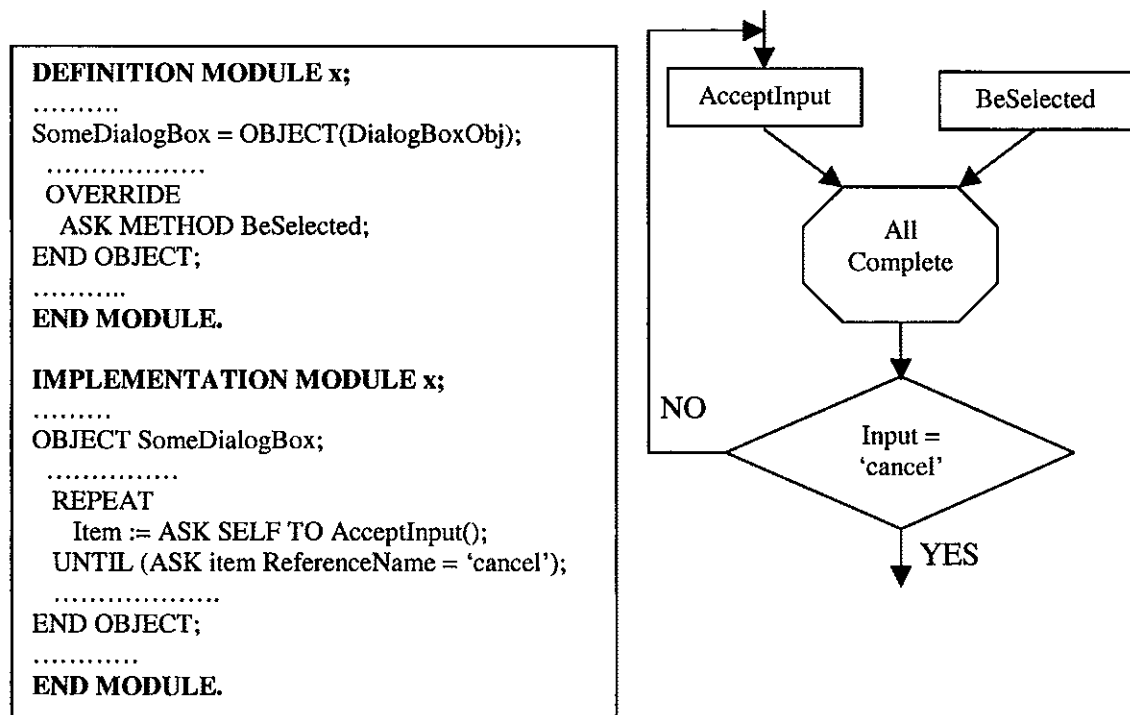


Fig. 11 Example 7: Representation of code where order of data flow is hidden in library routines



## **3.2 Developing Flow Diagrams**

This section describes the method of partitioning the code and representing it as a flow diagram by using the code associated with a method from PORTSIM as an in depth example. Two other examples can be found in Section 4.

### *3.2.1 Method*

The following steps describe the way to develop flow diagrams from MODSIM III code.

- 1) Define the functionality of the method to be decoded by its object type, name, in-put parameters, out put parameters and by the context of the code from which the method is being called.
- 2) Go through the code and mark off the method calls;
- 3) Mark off the activities, events and the asynchronous calls from method calls.
- 4) Identify sets of method calls where available that are used to perform a single task and name that set of methods depending on the task they are performing.
- 5) Identify any method calls that need to be further expanded (involved messages).
- 6) Construct the flow diagrams using the notations developed previously.

### 3.2.2 PORTSIM Example

This following section shows the result of applying the above mentioned procedure of converting code into flow diagrams to the method called “callForwardFlatcars” in PORTSIM'S code.

#### Step 1

The characteristics of the method and their values are given in TABLE 2

TABLE 2  
Characteristics of “CallForwardFlatcars” method

CHARACTERISTIC OF METHOD	VALUE
Object Type	Spur
Name of Method	CallForwardFlatcars
Input Parameters	Interchange Yard Object & Random Object
Output Parameters	None
Context	The method is being called by an instance of port object while initializing spurs.

*Functionality:* This method is used for getting flatcars from the interchange yard to the instance of a spur object to which this message is passed.

### Steps 2 through 5

Fig. 12 shows the result of applying the steps 2 through 5 on the code shown in Fig. C.1.

The following comments are made for Fig. 12 :

- 1) The statements in boxes represents synchronous method calls out of which statements containing "Wait" represent activities and remaining statements represent events.
- 2) Asynchronous calls are represented by bold text.
- 3) Sets of method calls performing a single task are marked along with their names.
- 4) Statements in boxes that are followed by "expand" represents method calls that can be treated as logical processes.

```

TELL METHOD callForwardFlatcars(IN interchangeYard :interchangeYardObj;IN ranGen :RandomObj);
BEGIN
  WHILE (SimTime() <= (simBeginTime + timeToSimulate))
    IF (ordering = "Vehicles/Unit Equipment Only")
      In IY Count Flatcars with Vehicles as their First Cargo
      WHILE (numFlatcarsCarryingVehicles = 0)
        Wait for Train Classified Trigger; In IY Count Flatcars with Vehicles as their First Cargo
      END WHILE;
    ELSIF (ordering = "Containers Only")
      In IY Count Flatcars with Container as their First Cargo
      WHILE (numFlatcarsCarryingContainers = 0)
        Wait for Train Classified Trigger; In IY Count Flatcars with Container as their First Cargo
      END WHILE;
    END IF;
    IF (status = "idle")
      Change Status to Busy
      WHILE (stringOffFlatcars.numberIn = 0)
        Build List of Flatcars at IY (expand)
        IF (stringOffFlatcars.numberIn = 0)
          Wait for 30 Minutes
        END IF;
      END WHILE;
      TELL IYTrigger TO Trigger;
      Wait for Locomotive; Wait for Locomotive arrival at IY; Wait for Coupling Flatcars at IY
      IF (type = "OPEN STAGING")
        Wait for Flatcars Switch to Spur
      ELSIF (type = "COVERED STAGING")
        Wait for Flatcars Switch to Dock
      ELSE
        Wait for Flatcars Switch to Berth
      END IF;
      Wait to Uncouple Flatcars; Return Locomotive; Wait to Discharge all Flatcar contents(expand)
      Change Status to Idle
      TELL spurDoneTrigger TO Trigger;
    ELSE
      Wait for Spur Done Trigger
    END IF;
  END WHILE;
END METHOD {callForwardFlatcars};

```

Start "Flatcars Processing"

End "Flatcars Processing"

Fig. 12 Result of applying steps 2 through 5 to "callForwardFlatcars" method

**Step 6**

In this step, Fig. 12 is converted into flow diagrams using the previously adopted notations. Fig. 13 represents the program flow within the “callForwardFlatcars” logical process, with the identified logical processes in it marked properly. Fig. 14 represents the program flow within the “Flatcar Processing” logical process, with the identified logical processes in it marked properly. Every logical process in Fig. 13 and in Fig. 14 is followed by a figure number, which represents the program flow within that particular logical process.



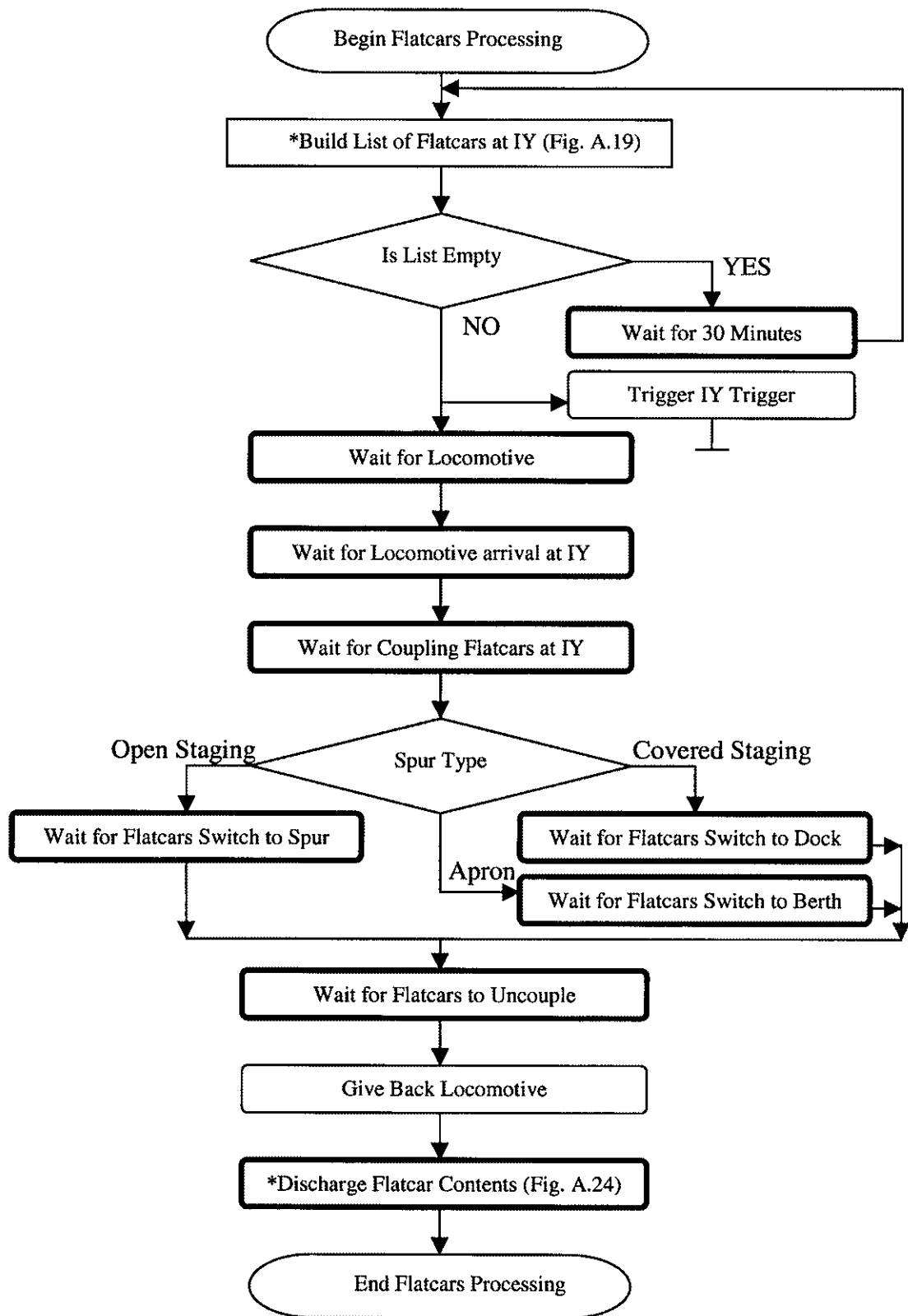


Fig. 14 Program flow within "Flatcars Processing"

### 3.3 Modifying Code using Flow Diagrams

This section shows how flow diagrams developed in the above mentioned fashion can be used to modify the code to fix a known bug.

#### *3.3.1 Method for modifying code using Flow Diagrams*

The following steps describe the way to use flow diagrams to fix bugs in MODSIM III code.

- 1) Define and analyze the problem.
- 2) Identify the set of flow diagrams representing the code that may contain the bug.
- 3) Go through the flow diagrams keeping in mind the type of problem and identify the problem causing segments within the flow diagrams.
- 4) Define the condition that needs to be satisfied to prevent the problem.
- 5) Change the flow diagrams so that the problem is rectified. The change may span more than one flow diagram as we have divided the system into logical processes.
- 6) Run through the flow diagrams again and make sure that the original functionality is not disturbed because of the changes made in the flow diagrams. Go to step 5 if the original functionality is changed.
- 7) Now convert the modified flow diagrams into MODSIM III constructs using the developed notations.
- 8) Execute the code with modifications and check if the problem is solved. If the problem persists, go to step 3.

The following section shows how the above mentioned procedure is applied in solving a problem.



### *3.3.2 PORTSIM Example: Call Forward Flatcars Problem*

#### **Step 1**

Problem definition: Only the first available interchange yard is being serviced, and all other available interchange yards are not serviced even though they have flatcars in them and spurs are available for use.

From the problem definition, we may surmise that the first serviced interchange yard is not allowing other interchange yards to get service or is forcing the code into an infinite loop.

#### **Step 2**

The flow diagrams involving interchange yards are shown in Fig. A.14 and Figures A.16 through A.21.

#### **Step 3**

- 1) From step 1, we decide that we need to concentrate on loops in the flow diagrams.
- 2) From Fig. A.14, we can see that each available interchange yard is trying to get served by a spur. But in order to understand the inner loop completely, we need to understand A.16.
- 3) From Fig. A.16, we can see that each available spur is calling flatcars from the interchange yard. Fig. A.16 points us to Fig. A.17, which expands the “call Forward Flatcars” logical process.

4) From Fig. A.17, we can see that:

- If the ordering is either “vehicles only” or “containers only,” the program control loops until there are flatcars in the interchange yard. However, we cannot say that this flow diagram segment is the one that is causing the problem because we need to solve the problem for interchange yards that have flatcars in them.
- If the spur is idle, then its status is changed to “busy” before executing the “flatcars processing” logical process and is changed back to “idle” after “flatcars processing” logical process. The interchange yard then sends “spur done” signal.
- If the spur is not idle, then the interchange yard waits until it receives “spur done” signal.

5) From Fig. A.18, which expands the “flatcars processing” logical process, we can see that:

- If there are no more flatcars left within the interchange yard, the program control loops until more flatcars are available in the interchange yard. We can say that this segment of the flow diagram is the one that is causing the problem. By the time the program control reaches the loop, the status of the spur is already marked as “busy,” making it unavailable to other interchange yards. *There is no way of getting out of this loop after all the flatcars in the interchange yard are processed.*

**Step 4**

The condition that needs to be satisfied to solve the problem is that an interchange yard should make the spur, to which it is sending flatcars, available to other waiting interchange yards after all of the flatcars in it are processed.

**Step 5**

We can remove the problem causing loop from the "Flatcars processing" logical process and add it before the spur's status check in the "call Forward Flatcars" logical process. This makes sure that the spur's status is marked "busy" only when there are flatcars in the interchange yard. Figures 15 and 16 show the modified parts of the flow diagrams shown in Figures 13 and 14.

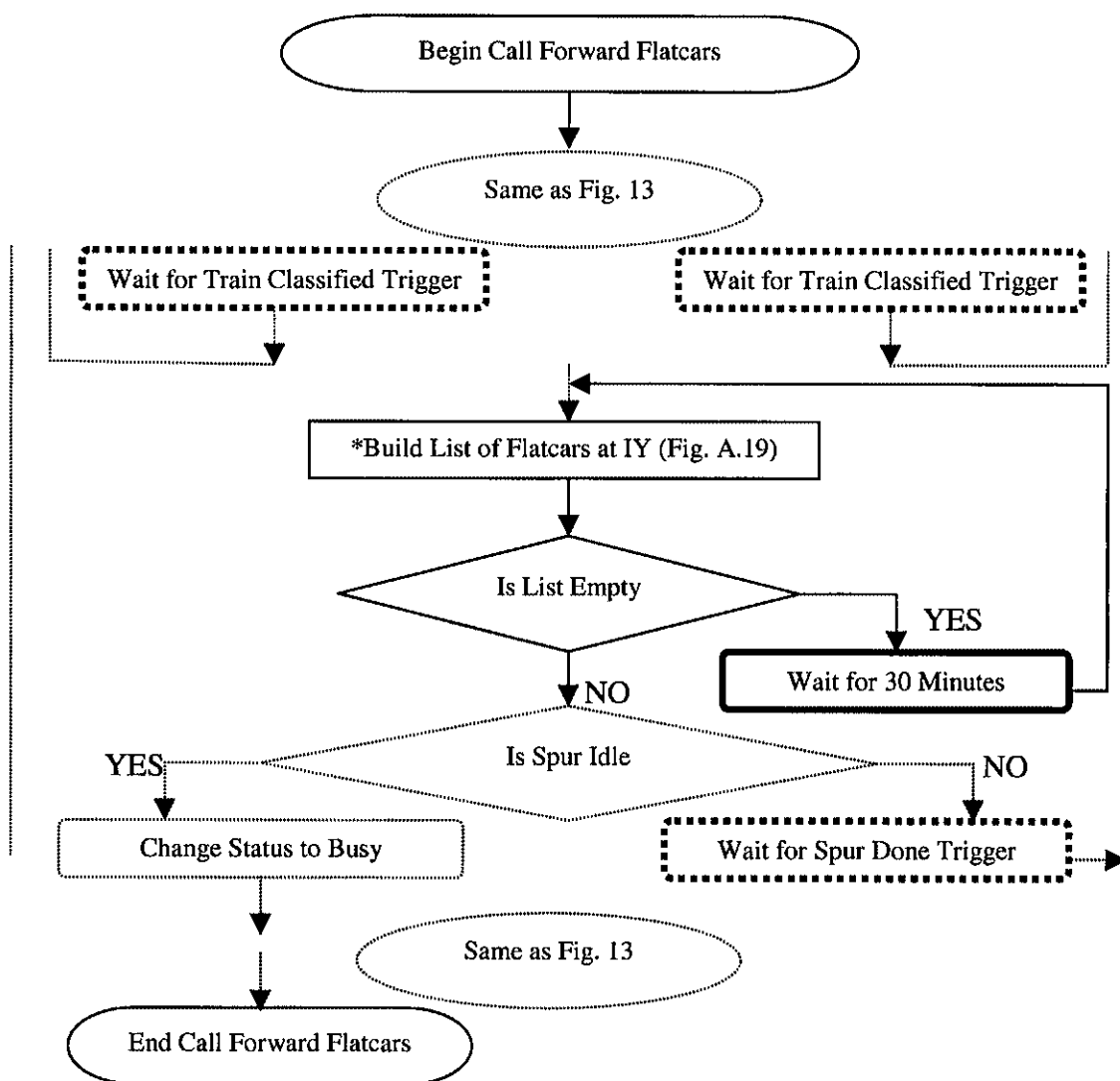


Fig. 15 Modified representation of Fig. 13

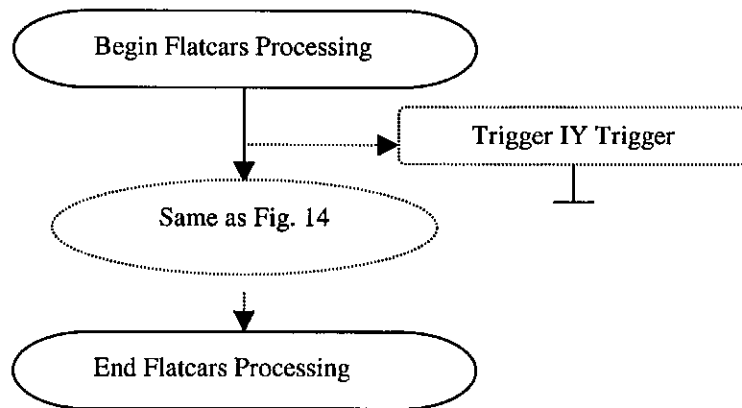


Fig. 16 Modified representation of Fig. 14

### Step 6

The model developed above satisfies the condition mentioned in step 4 by not allowing any interchange yard to monopolize the use of the spurs. However, to understand whether the above model leaves the original functionality of the "call Forward Flatcars" logical process unchanged, it is necessary to understand the effect of shifting the position of "Build List of Flatcars at IY" logical process along with the loop. From Fig. A.19, which expands "Build List of Flatcars at IY" logical process, we can see that it is necessary to understand the "Adding Flatcars with Vehicles" logical process. From figure A.20, which expands the "Adding Flatcars with Vehicles" logical process, we can see that

- 1) While building the list of flatcars, flatcars are removed from the interchange yard, and the capacity of the interchange yard is incremented by their length.

If after building a list of flatcars the spur is found to be busy, the interchange yard will wait for a spur done signal. After receiving the signal it again loops back and builds

another list of flatcars, so the flatcars in the previous list were destroyed without being processed.

### **Step 5 Revisited**

We can change the condition for marking the spur's status as "Busy," without shifting the loop from the "Flatcars Processing" logical process to the "call Forward Flatcars" logical process, in such a way that the spur is marked busy only when there are flatcars in an interchange yard. This can be done by marking the status of the spur as "Busy" after the problem causing loop in the "Flatcars processing" logical process instead of in the "call Forward Flatcars" logical process. Now, if the program control loops indefinitely when there are no flatcars left to process in the interchange yard, the status of the spur will remain as "Idle" making it available for other interchange yards. Also, we need to check the status of the spur in the loop before building a list of flatcars from an interchange yard to make sure that the spur is still available after waiting for thirty minutes for more flatcars in the interchange yard. Figures 17 and 18 show the modified parts of the flow diagrams shown in Figures 13 and 14.

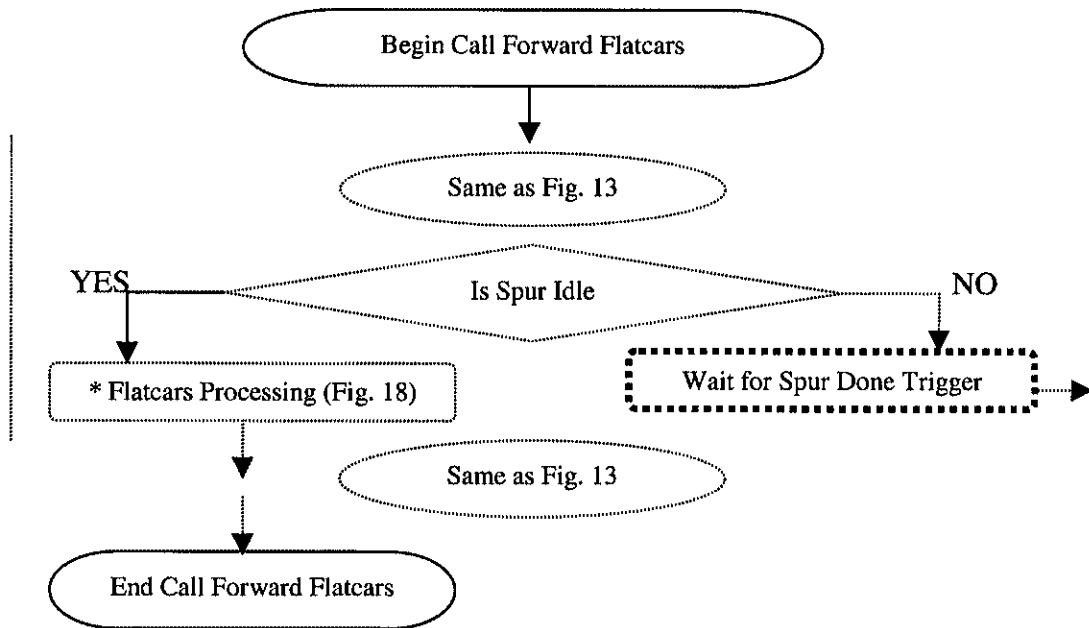


Fig. 17 Revised representation of Fig. 13

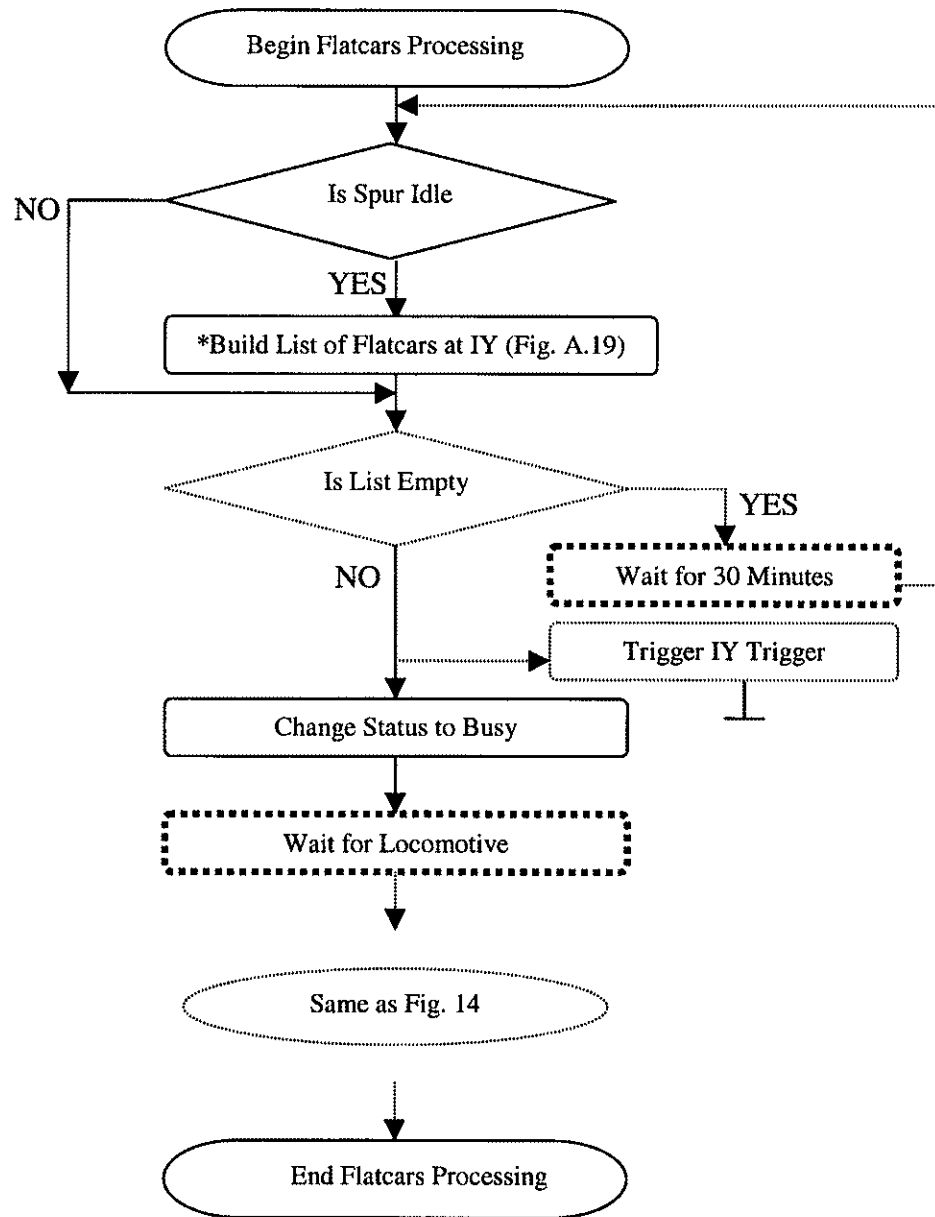


Fig. 18 Revised representation of Fig. 14



### **Step 6 revisited**

The alternate model developed above functions exactly like the original model when there are flatcars in an interchange yard to process, and it also satisfies the condition mentioned in step 4.

### **Step 7**

Now we need to convert the flow diagrams shown in Figures 17 and 18 into MODSIM III constructs using the notations developed. Fig. 19 shows the MODSIM III constructs where the changes are indicated by bold italic text. The modified code is shown in Fig. C.2.

### **Step 8**

The modified code was executed and found that the available spurs are servicing all interchange yards with available flatcars.

```

TELL METHOD callForwardFlatcars(IN interchangeYard :interchangeYardObj;IN ranGen :RandomObj);
BEGIN
  WHILE (SimTime() <= (simBeginTime + timeToSimulate))
    IF (ordering = "Vehicles/Unit Equipment Only")
      In IY Count Flatcars with Vehicles as their First Cargo
      WHILE (numFlatcarsCarryingVehicles = 0)
        Wait for Train Classified Trigger; In IY Count Flatcars with Vehicles as their First Cargo
      END WHILE;
    ELSIF (ordering = "Containers Only")
      In IY Count Flatcars with Container as their First Cargo
      WHILE (numFlatcarsCarryingContainers = 0)
        Wait for Train Classified Trigger; In IY Count Flatcars with Container as their First Cargo
      END WHILE;
    END IF;
    IF (status = "idle")
      Removed the call to spur status change in the IF construct
      WHILE (stringOfFlatcars.numberIn = 0)
        Start "Flatcars Processing"
        IF (status = "idle")
          Build List of Flatcars at IY (expand)
        END IF;
        IF (stringOfFlatcars.numberIn = 0)
          Wait for 30 Minutes
        END IF;
      END WHILE;
      TELL IYTrigger TO Trigger;
      Change Status to Busy
      Wait for Locomotive; Wait for Locomotive arrival at IY; Wait for Coupling Flatcars at IY
      IF (type = "OPEN STAGING")
        Wait for Flatcars Switch to Spur
      ELSIF (type = "COVERED STAGING")
        Wait for Flatcars Switch to Dock
      ELSE
        Wait for Flatcars Switch to Berth
      END IF;
      Wait to Uncouple Flatcars; Return Locomotive; Wait to Discharge all Flatcar contents(expand)
      Change Status to Idle
      End "Flatcars Processing"
      TELL spurDoneTrigger TO Trigger;
    ELSE
      Wait for Spur Done Trigger
    END IF;
  END WHILE;
END METHOD {callForwardFlatcars};

```

Fig. 19 MODSIM III constructs showing the changes made in Figures 13 & 14

## **SECTION 4**

### **DEMONSTRATION OF USING FLOW DIAGRAMS FOR CODE MODIFICATION IN PORTSIM**

This section demonstrates the use of the model developed in the previous section to fix bugs and to modify the code using flowcharts. The “Ship Exiting Problem” as described below demonstrates a second use of the model in fixing bugs similar to the “Call Forward Flatcars” problem detailed in Section 3. The “Dialog Boxes Closing Problem,” as described later, demonstrates the use of the model in making improvements in the software.

#### **4.1 Ship Exiting Problem**

##### **Step 1**

Problem definition: Ships are not leaving a port if they have remaining cargo space and all cargo has been loaded onto a ship.

From the problem definition, we may surmise that this problem is due to either an infinite loop or an inappropriate condition that allows a ship to leave the port.

##### **Step 2**

The flow diagrams involving ship life cycle are shown in Figures A.46 through A.78.

### Step 3

1) From step 1, we decide that we need to concentrate on loops and conditions in the flow diagrams.

2) From Fig. A.46 we can see that

- The time the last cargo was loaded onto the ship is set to the present simulation time.
- A ship in the arrival list is then sent to dock at a berth.

In order to understand the docking process, we need to understand Fig. A.47.

3) From Fig. A.47 we can see that

- If none of the berths are available for military use then the ship is added to ship queue.
- If any of the berths are available, berthing is done depending on the ship type.

In order to understand the berthing process for various types of ships, we need to understand Figures A.48 through A.51.

4) From Figures A.48 to A.51, we can see that to understand these figures we need to understand Fig. A.52.

5) From Fig. A.52, we can see that if a berth is available and suitable for the ship, then

- The berth is marked busy.
- The ship is served or loaded at that berth.

In order to understand the process of serving a ship, we need to understand Fig. A.53.

6) From Fig. A.53, we can see that

- If space on the ship is not available, the ship is released from berth.

- If space on the ship is available, then
  - Staging areas with maximum number of vehicles, containers and pallets ready to load are selected.
  - If none of the cargo items are available then the program waits for sixty minutes.
  - If any of the cargo items are available, those cargo items are call forwarded from their respective staging areas.
  - Cargo items are loaded onto the ship.
  - If any of the cargo items are left on berth, those cargo items are sent back to staging areas.
  - If the difference between present time and the time the last cargo item was loaded is less than the maximum wait time for cargo of that ship, the above process is repeated. Otherwise, the ship is released from the berth. This implies that a ship is released from berth if none of the cargo items are available within the maximum allowed wait time for cargo of the ship.

But in order to understand the processes of call forwarding cargo, loading cargo and ship release, we need to understand Figures A.54, A.66 and A.78, respectively.

- 7) From Fig. A.54, we can see that a call forward list is built depending on the ship type, so to understand Fig. A.54 we need to understand Figures A.55, A.56 and A.63.

8) From Figures A.55, A.56 and A.63 we can see that

- If the cargo items and the port assets required to load them onto the ship are not available, the program again waits for sixty minutes and checks for the availability of cargo items and the port assets required by them. This loop continues till both the cargo items and the port assets required by them are available.

We can say that this segment of the flow diagrams is the one that is causing the problem. If there are no more cargo items left to occupy the present staging area, from which the call forward list is being built, there is no way for the program control to get out of this loop. So, the ship cannot be released from the berth once the program control is caught in this loop.

#### **Step 4**

The condition that needs to be satisfied to solve the problem is that a ship should be released from a berth when all of the cargo items in a staging area are loaded onto the ship.

#### **Step 5**

We can remove the defective loop from the Figures A.55, A.56 and A.63 and convert it into a condition statement as shown in Figures 20, 21 and 22, respectively. This will allow a ship to be released from the berth when all the cargo items in a staging area are loaded onto it by making the "Call Forward Cargo" process behave as a WAIT statement, as shown in Fig. 23.

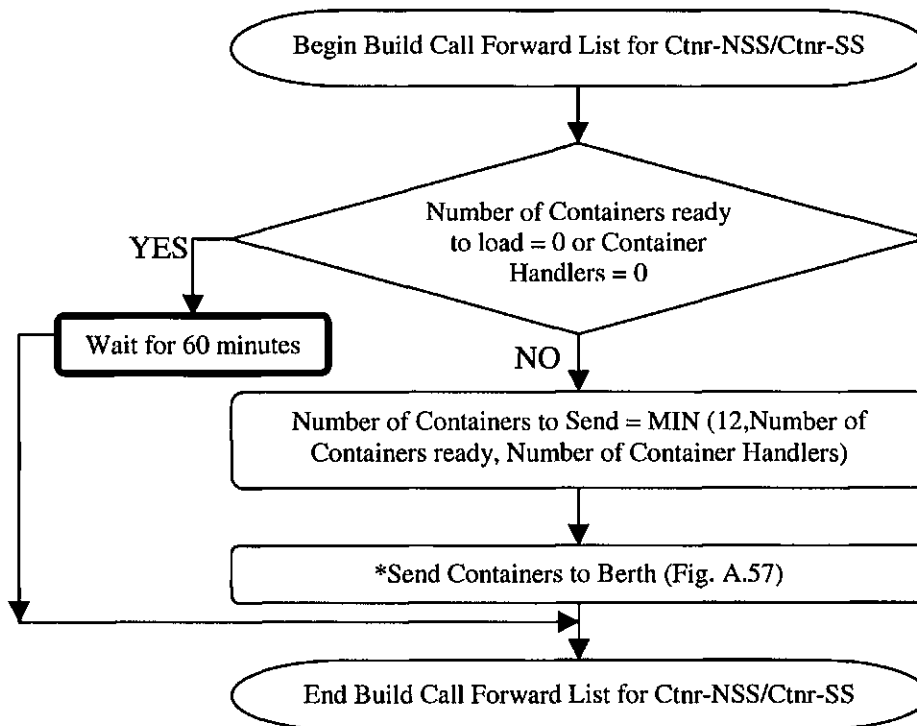


Fig. 20 Modified representation of Fig. A.55

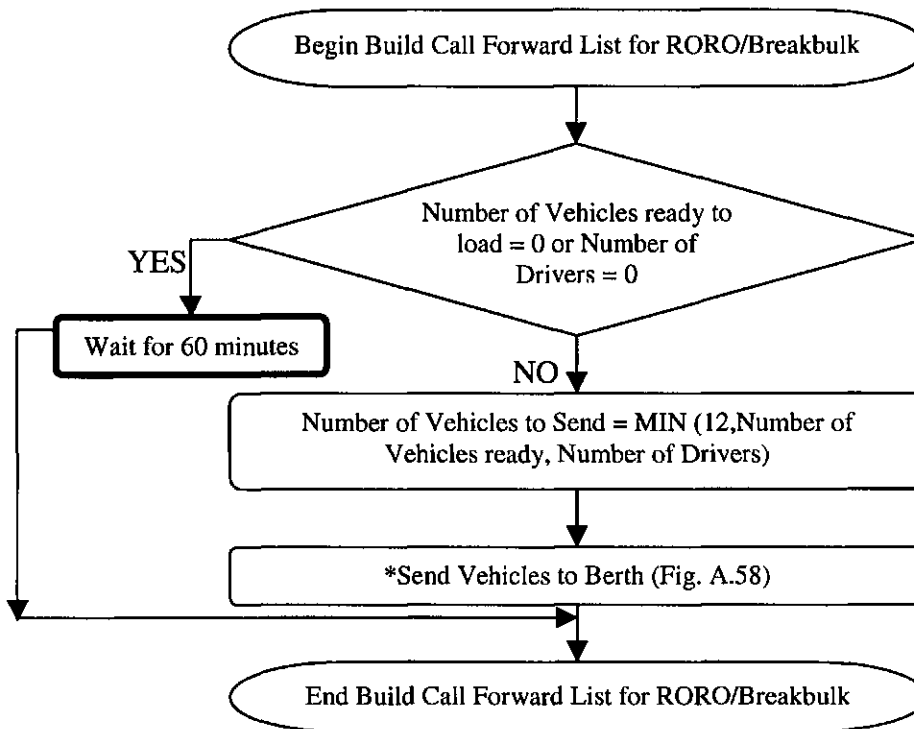


Fig. 21 Modified representation of Fig. A.56

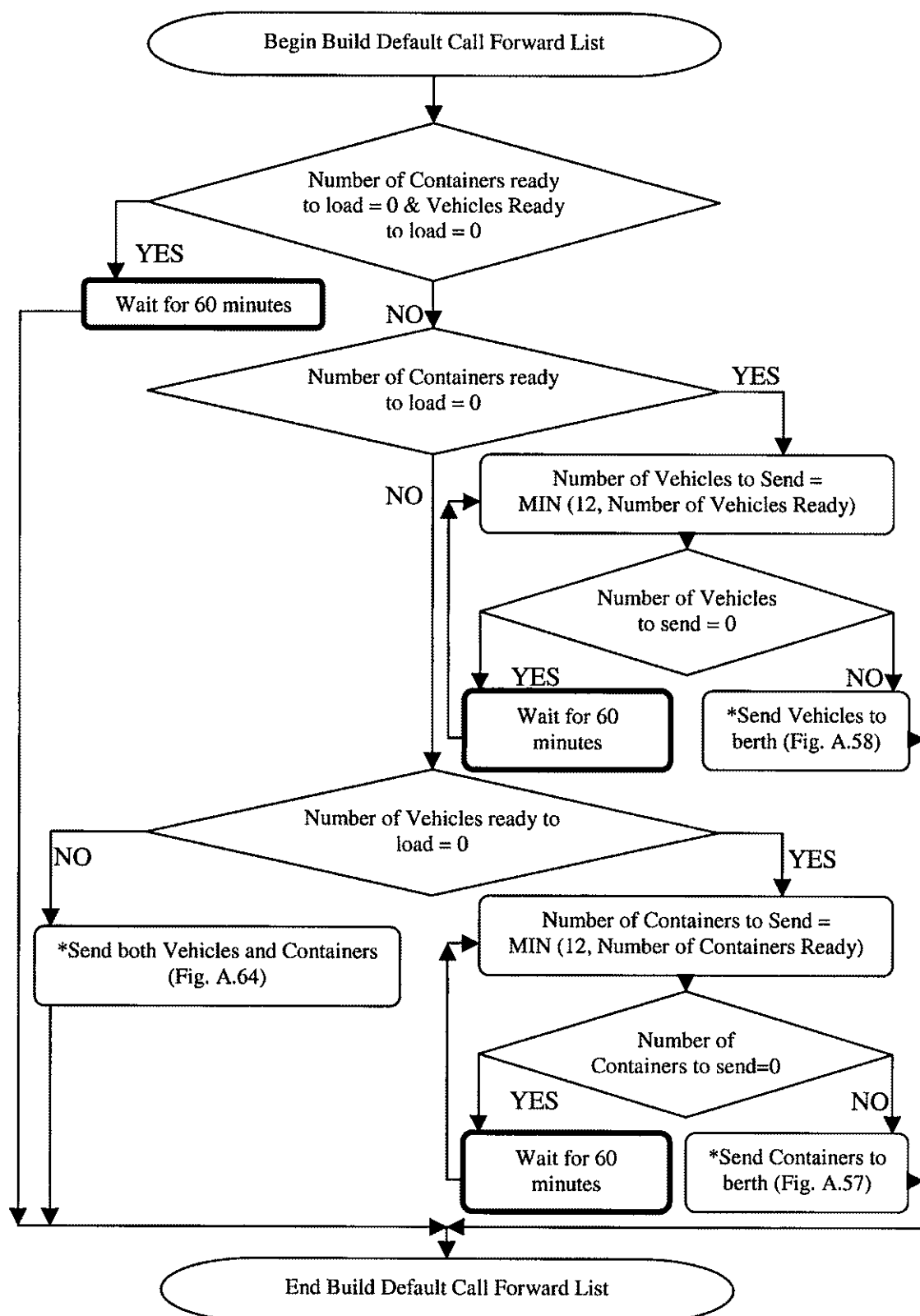


Fig. 22 Modified representation of Fig. A.63



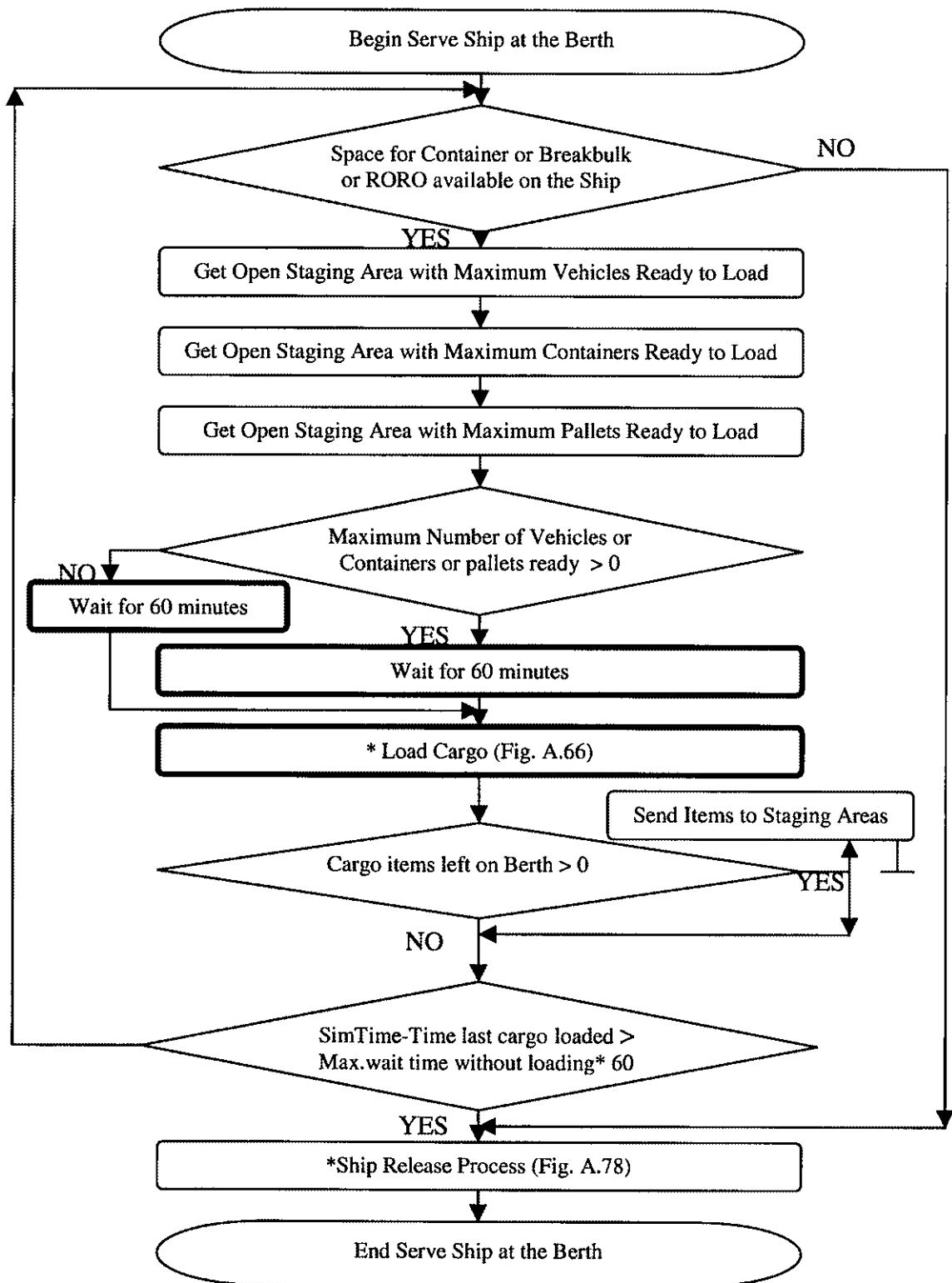


Fig. 23 Logical representation of Fig. A.53 when there is no cargo left to load

**Step 6**

From the understanding of Fig. A.53, we can see that we made the berth release the ship when there are no cargo items left to load. However, to see whether the above model leaves the original functionality of the call forward cargo process unchanged, it is necessary to understand the purpose of using the loop during the preparation of call forward list. From Figures A.55, A.56 and A.63, we can see that the loop is making sure that the program waits for at least sixty minutes for more cargo to arrive at the staging area. The model developed above is not supporting this minimum waiting period functionality.

**Step 5 Revisited**

We can remove the loop and wait for sixty minutes if there are no cargo items left in the staging area. Once the program control returns from the wait statement, we can check for new cargo items in the staging area again. Figures 24, 25 and 26 show this new model.

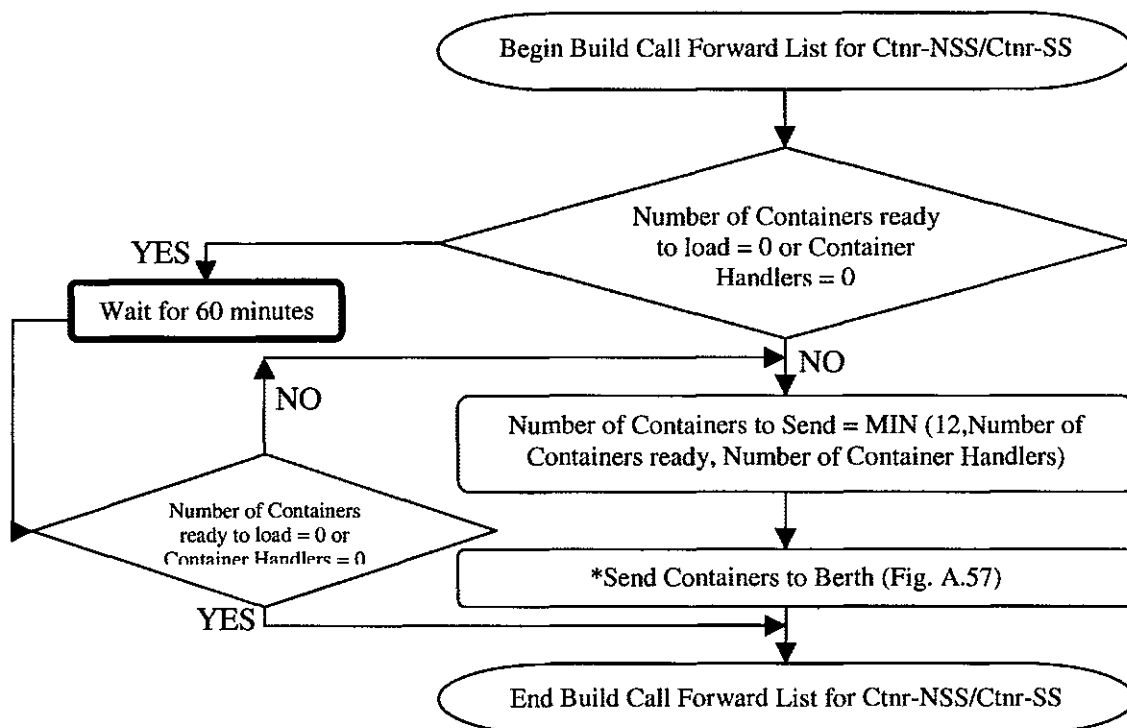


Fig. 24 Modified representation of Fig. 20

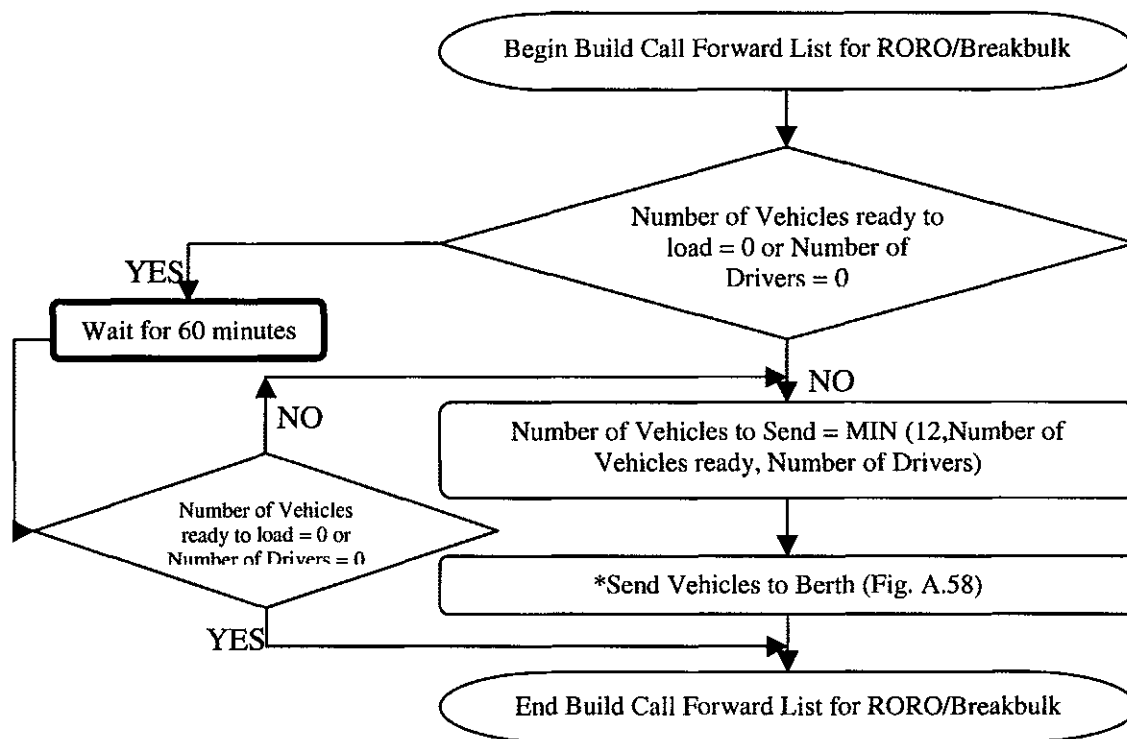


Fig. 25 Modified representation of Fig. 21

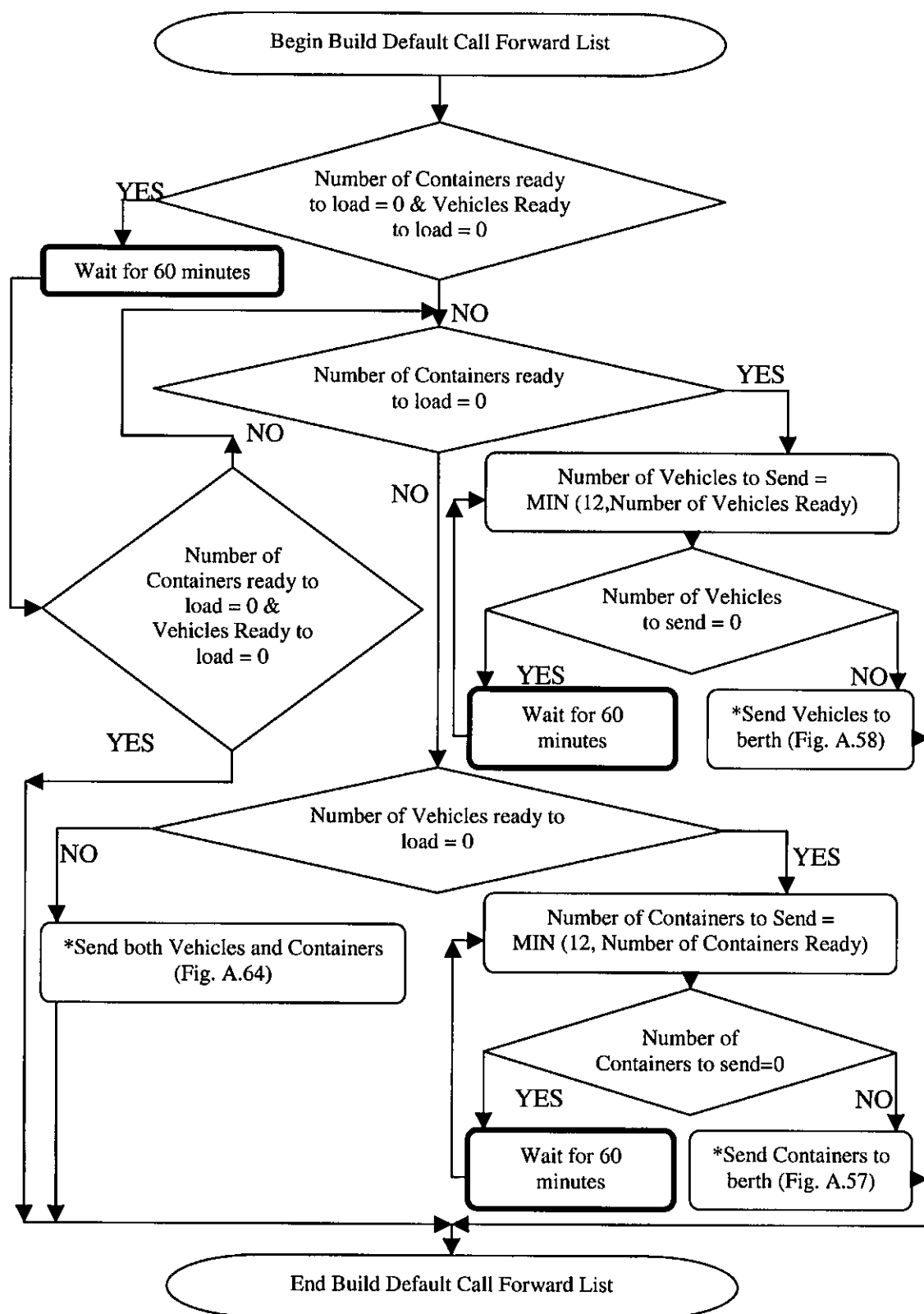


Fig. 26 Modified representation for Fig. 22

### **Step 6 Revisited**

The alternate model developed above also makes the "Call Forward Cargo" process behave as a WAIT statement when there are no more cargo items left to arrive at the staging area, and it also satisfies the minimum waiting period for more cargo items to arrive at the staging area.

### **Step 7**

Now we can convert the flow diagrams shown in Figures 24, 25 and 26 into MODSIM III constructs using the notations developed. The MODSIM III code for these figures is shown in Fig. C.3.

### **Step 8**

The modified code has been executed and we found that the ships were leaving the port when there is no more cargo left to load, irrespective of the remaining cargo space on ships.

It should also be noted that this fix allows the code to operate as originally modeled. However, this is not the correct port operation. A significant effort is required to develop a new berthing/call forward model. This need was not uncovered until this modeling effort allowed insight to the implementation model by the domain experts.

## **4.2 Dialog Boxes Closing Problem**

### **Step 1**

Problem definition: Users are not able to close any of the graphical windows that pop up during the course of creating and running a scenario, by clicking on the window's close button. Close buttons are simply the X buttons on the right hand top corner of the graphical windows. The main graphical windows behave correctly; only the sub-windows require examination.

From the problem statement, we may surmise that there is a difference in the object types of the windows that are responding to the close buttons and the windows that are not responding to the close buttons. The difference in the object types may be due to the difference in the base object types or due to the difference in the way the objects are derived from base object.

### **Step 2**

The flow diagrams representing the functionality of various graphical window objects are shown in Figures B.1 through B.7.

### Step 3

- 1) From step 1, we decide that we need to concentrate on the difference between the graphical object types and the code associated with them.
- 2) From Figures B.1 through B.7 we can see that
  - There is a difference in the base object types from which the graphical objects shown in Fig. B.1 and Figures B.2 to B.7 are derived. One is of WindowObj type and another is of DialogBoxObj type.
  - The “cancel” button, which is being used to stop/cancel any process of modifying the parameters of an object that the graphical window is representing, closes the graphical window.
- 3) From Fig. B.1, which represents the functionality of a graphical object derived from WindowObj, we can see that there is no explicit code written to respond to the close button. Same is the case with Figures B.2 through B.7, which represent the functionality of graphical objects derived from DialogBoxObj.

From Section 3.1.1.2, we know that whenever a user clicks on the close button of a graphical object, the corresponding BeClosed method is called automatically. The default action of the BeClosed method for the WindowObj object is to terminate the application. Also, from problem definition, we infer that the graphical objects that are responding to close buttons are of type WindowObj. Thus, we can infer that the default action of the BeClosed method for the DialogBoxObj object is to do nothing.

#### **Step 4**

We need to overwrite the `BeClosed` method for the objects of type `DialogBoxObj` in such a way that it actually helps to close the graphical window to which it belongs.

#### **Step 5**

The important decision has to be made about the complete functionality of the close button before solving this problem. We should keep in mind that by achieving the condition mentioned in step 4, we are actually changing the behavior of the dialog boxes. Therefore, we need to concentrate on the consistent behavior of the dialog boxes. As closing generally signifies that the graphical object and hence its services are not needed anymore, it is a good idea to make the complete functionality of the close button same as that of “cancel” button of that dialog box.

From Figures B.2 through B.7, we can see that it is impossible to have a unique solution for all types of dialog boxes and yet have their previous functionality unchanged. Thus, the best way to solve this problem is by finding a solution for each type of dialog box.

##### **1. Solution for instances of `DialogBoxObj` object:**

In order to overwrite the default behavior of the `BeClosed` method of the `DialogBoxObj` object, an additional object of type `wrapDialogBoxObj` is derived from `DialogBoxObj` object with only its `BeClosed` method overwritten. The overwritten method should contain code to make the `AcceptInput` return when a user clicks on the close button. This can be achieved by either disposing the instance or erasing the instance within the



BeClosed method. However, from Figures B.2 through B.6 we can see that the instance cannot be disposed as it is referred after AcceptInput returns, so the best way to make the AcceptInput return is by erasing the instance in the Beclosed method. Fig. 27 shows the representation of BeClosed method.

Solution for Fig. B.2:

From Fig. B.2, we can see that the above described modification alone is sufficient to solve the problem for this type of instances of DialogBoxObj object. The solution is achieved by simply changing the type of instances with behavior shown in Fig. B.2 from DialogBoxObj to wrapDialogBoxobj. Fig. 28 shows the modified representation of Fig. B.2.

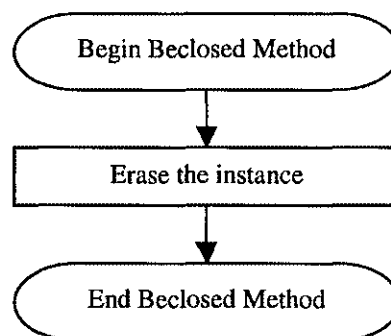


Fig. 27 Representation of BeClosed method using "Erase"

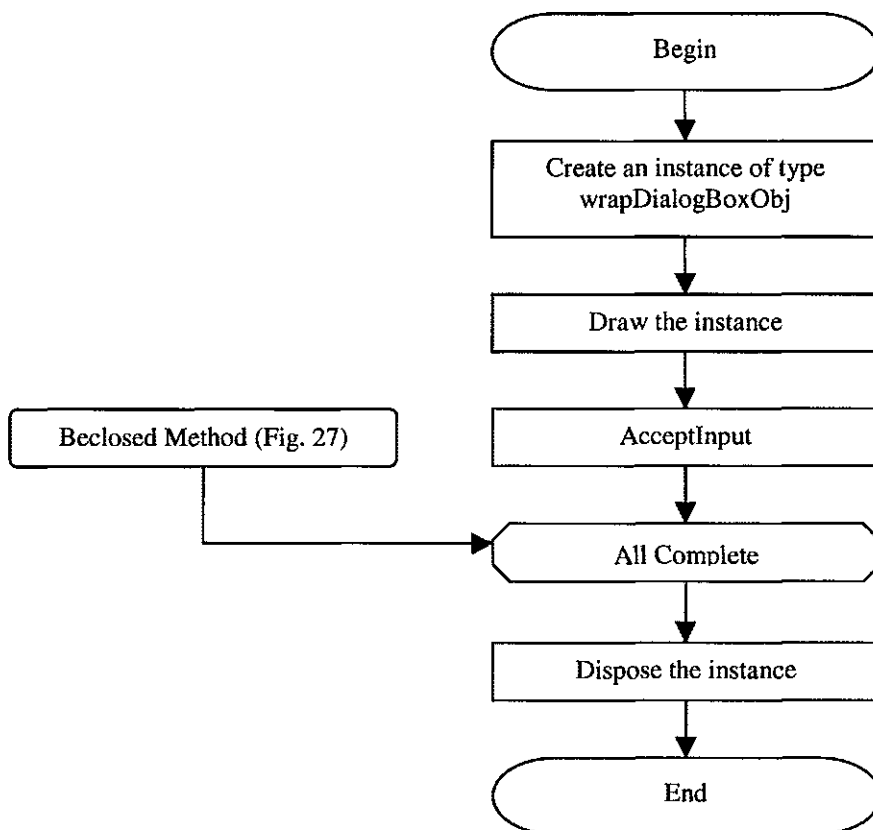


Fig. 29 Modified representation of Fig. B.2

Solution for Fig. B.3:

From Fig. B.3, we can see that the loop has to be broken in addition to changing the type of instance from DialogBoxObj to wrapDialogBoxObj to solve the problem for this type of instances. The loop can be broken by checking the returned value from the AcceptInput method. If the returned object is of type NILOBJ, we exit from the loop

[3.1.1.2]. The returned object type has to be checked again outside the loop to make the behavior of the close button the same as that of the cancel button. Fig. 29 shows the modified representation of Fig. B.3.

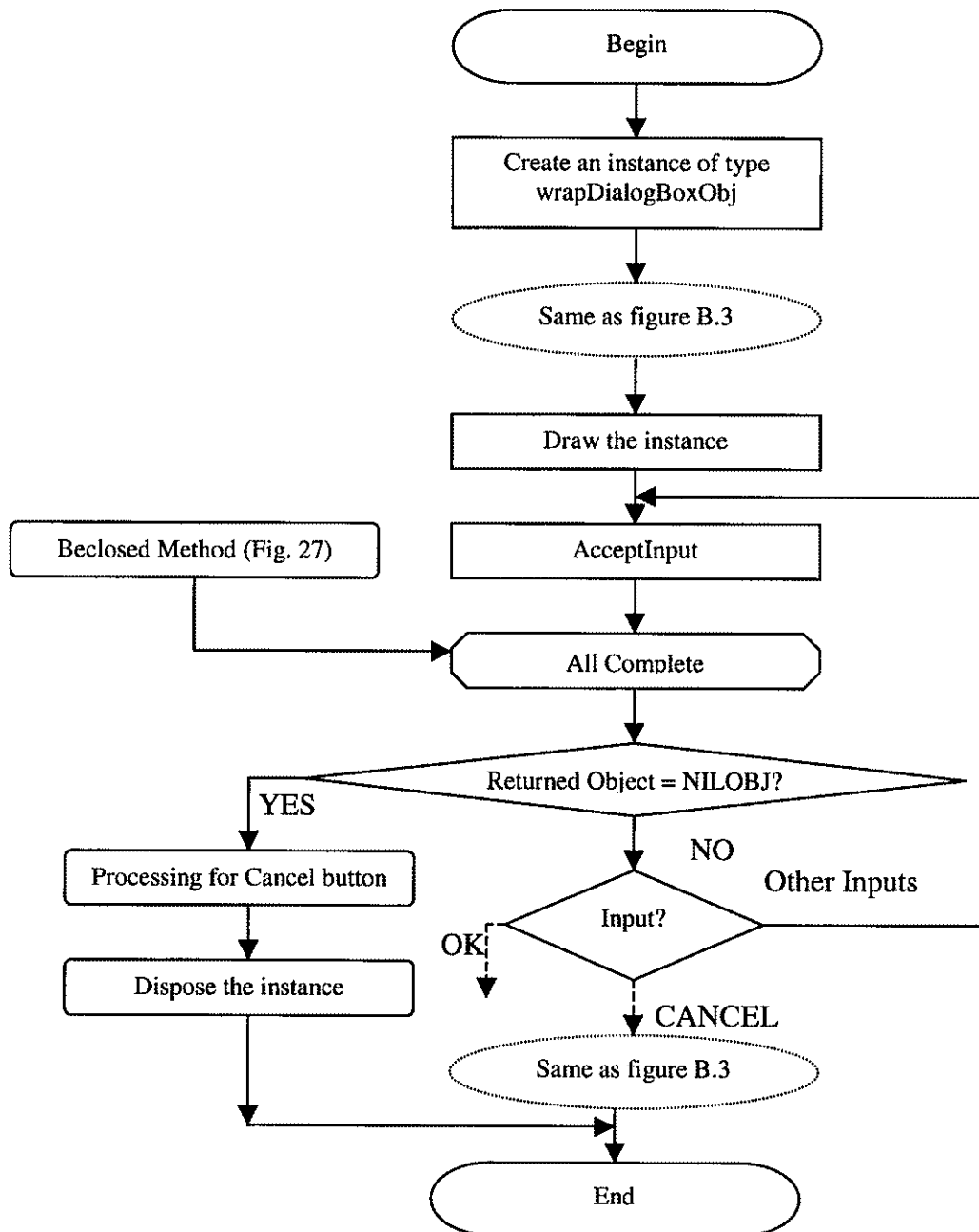


Fig. 29 Modified representation of Fig. B.3

Solution for Fig. B.4:

From Fig. B.4, we can see that for this type of instances we also need to check the object type of the returned object from AcceptInput method immediately after it returns in order to make the behavior of the close button same as that of the cancel button. Fig. 30 shows the modified representation of Fig. B.4.

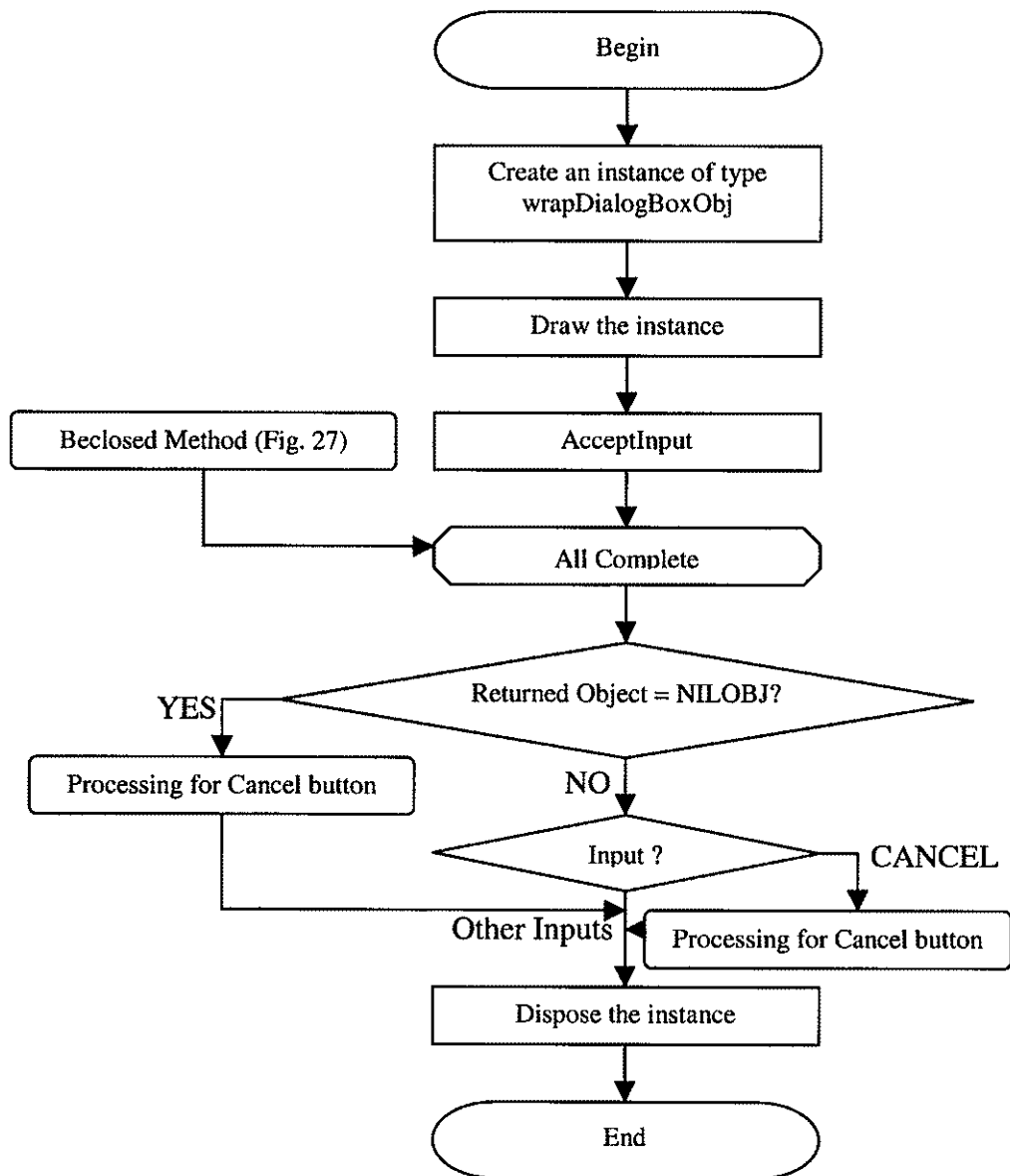


Fig. 30 Modified representation of Fig. B.4

Solution for Fig. B.5:

From Fig. B.5 we can see that these types of instances are not accepting any input from a user. So this kind of dialog boxes are to be left alone.

2. Solution for instances of objects derived from DialogBoxObj object:

As dialog boxes have an object type that is inherited from DialogBoxObj, we can simply overwrite the BeClosed method in the inherited objects itself.

Solution for Fig. B.6:

From Fig. B.6a, we can see that this type of dialog box is not referenced once after the cancel button is pressed. Therefore, just copying the code for the cancel button into the body of BeClosed method is sufficient to solve the problem. The code for the cancel button can be achieved from the dialog boxes BeSelected method under the “cancel” option. From Fig. B.6b, we can see that the functionality of cancel button is to dispose the instance, so we dispose the instance in BeClosed method as shown in Fig. 31. Fig. 32 shows the modified representation of Fig. B.6a.

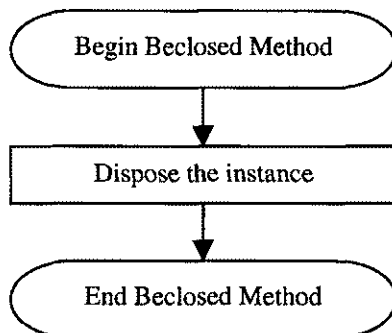


Fig. 31 Representation of BeClosed method using "Dispose"

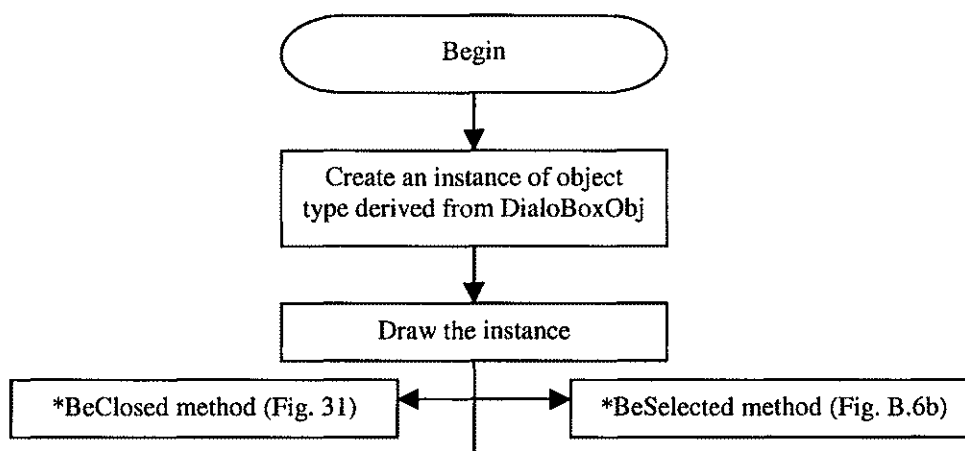


Fig. 32 Modified representation of Fig. B.6a

Solution for Fig. B.7:

From Fig. B.7, we can see that these dialog boxes have two different types of behavior. The two different behaviors are shown in Figures B.8a and B.9a, respectively. The similarity between these figures is that in both cases the instances are referenced after the AcceptInput method returns, so we need to:

- 1) Erase the instance in BeClosed method, as shown in Fig. 27.
- 2) Check the returned object type from AcceptInput method for NILOBJ type.

The difference between Figures B.8a and B.9a is that the instances are disposed in Fig. B.8a but are not in Fig. B.9a. We can see that this difference between Figures B.8a and B.9a is due to the difference in the behavior of the cancel button in the corresponding BeSelected methods, shown in Figures B.8b and B.9b respectively. From Fig. B.8b, we can see that the function of the cancel button is to do nothing. From Fig. B.9b, we can see that the functionality of the cancel button is to check for a closing condition and if the condition is satisfied the instance is removed or else an error message is displayed. So, in order to make the functionality of X button the same as that of cancel button, in the case of Fig. B.8a, it is sufficient to erase the instance in BeClosed method and check the returned object type from Acceptinput method for NILOBJ type in the loop of Fig. B.8a. However, in the case of Fig. B.9a, we need to place the functionality of cancel button within the loop if the returned object type from Acceptinput method is of type NILOBJ. Figures 33 and 34 show the modified representations of Figures B.8a and B.9a, respectively.

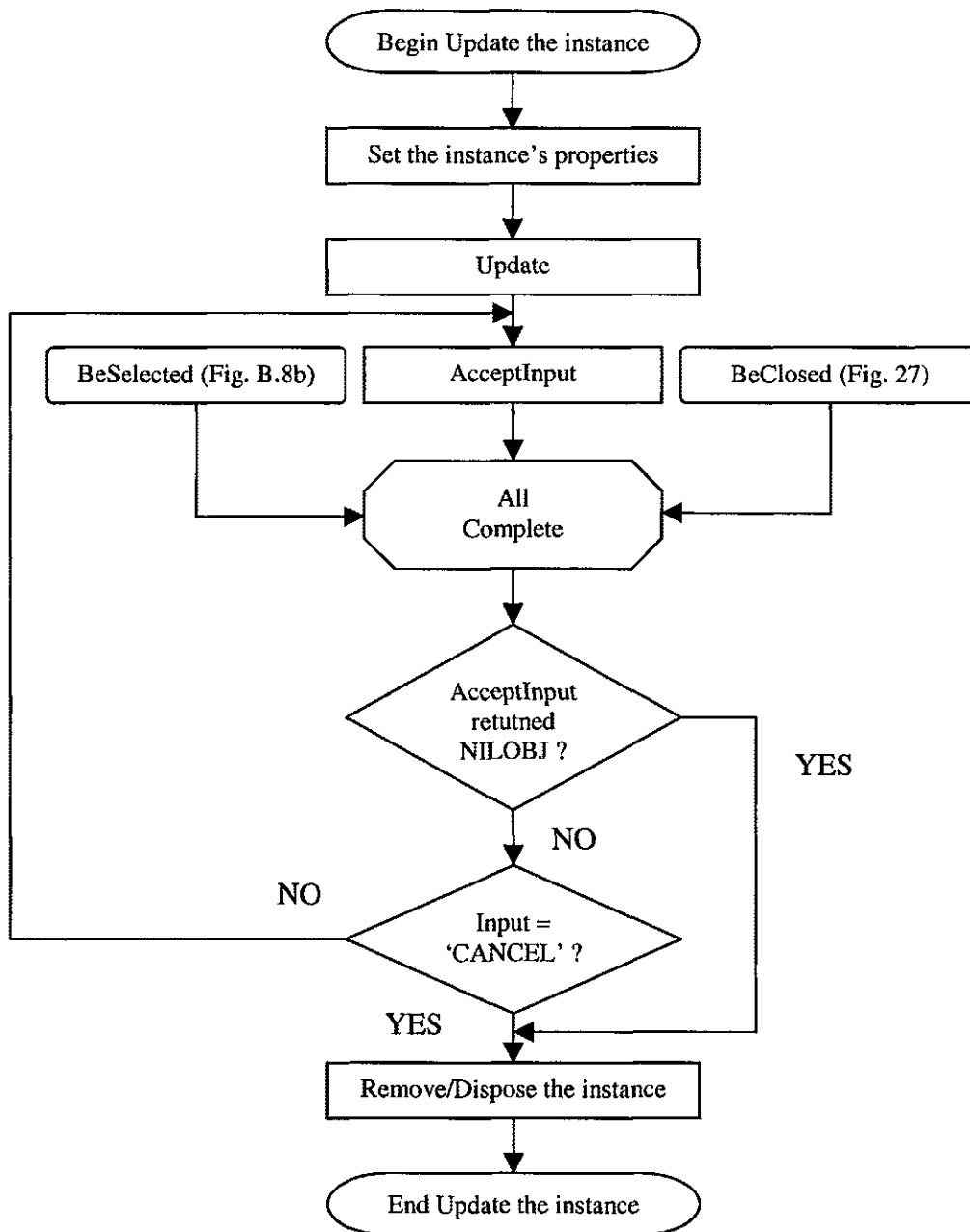


Fig. 33 Modified representation of Fig. B.8a



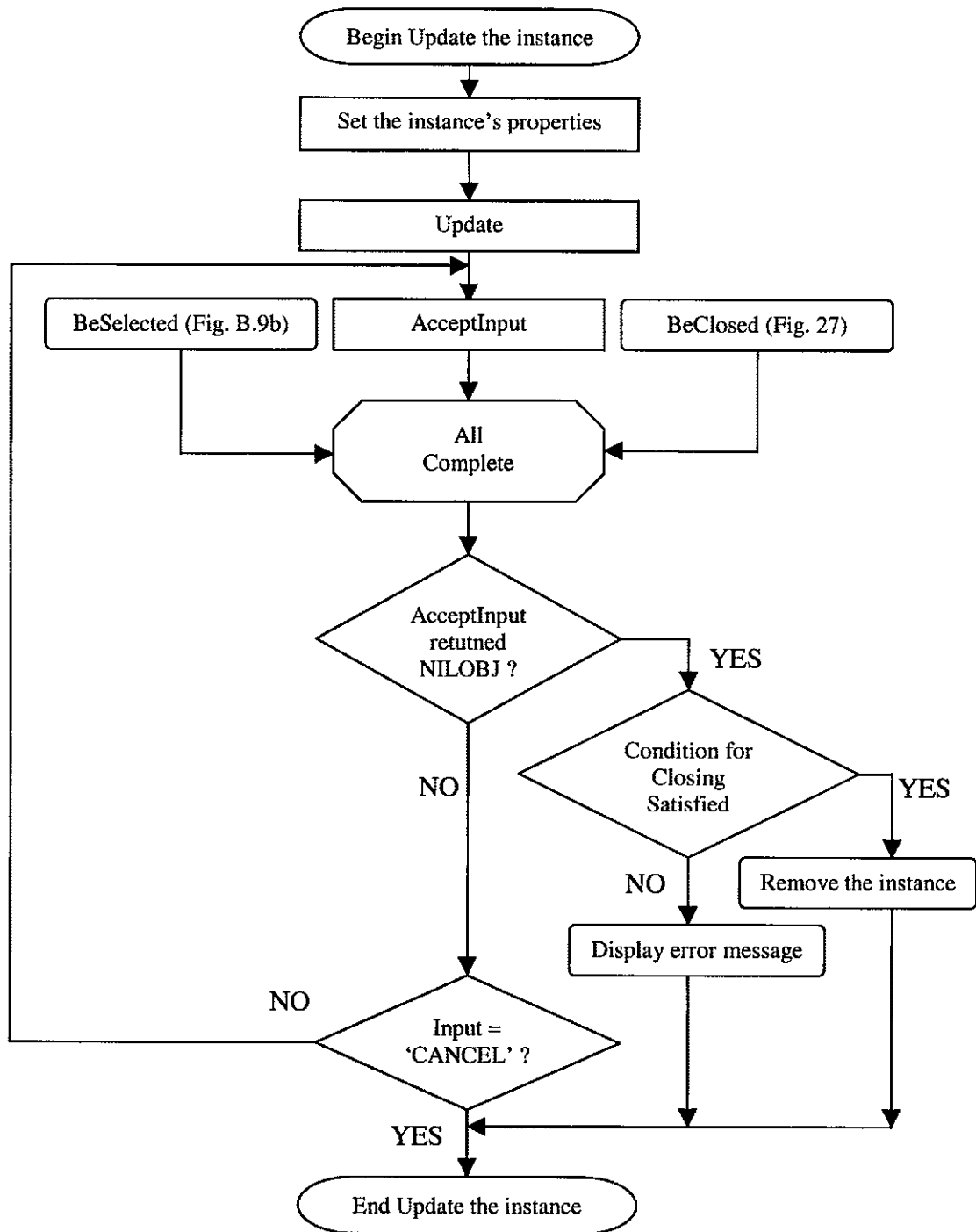


Fig. 34 Modified representation of Fig. B.9a

**Step 6**

The models developed above satisfy the condition mentioned in step 4 by making the dialog boxes disappear whenever a user clicks on the close button of a dialog box. The models developed above also makes sure that the functionality of the close button and the corresponding cancel button is same.

**Step 7**

Now we can convert the flow diagrams shown in Figures 27 to 34 into MODSIM III constructs using the notations developed. The codes for Figures 27 to 34 are shown in Figures C.4 through C.9.

**Step 8**

The modified code has been executed, and we found that all dialog boxes are closing when clicked on respective close buttons. Also, dialog boxes whose life cycles are controlled by the underlying process are not closing by their close buttons.

## **SECTION 5**

### **CONCLUSIONS**

This section summarizes the important achievements of the developed flowcharts model for pre-existing MODSIM software systems and later discusses drawbacks and possible enhancements.

#### **5.1 Achievements**

This section discusses the benefits of the flowchart model developed in Section 3.

- 1) Flowcharts help to understand the underlying processes that are responsible for accomplishing various tasks by the software.
- 2) The logical separation of processes helps to find the part of code that is most likely to have a bug.
- 3) The logical processes help to understand the side effects, if any, of the proposed solution to any given problem.
- 4) As whole processes are represented using flowcharts, it is quiet easy to understand the underlying model of any process.
- 5) As needed, the flowcharts do carry the information about the programming language environment and thus enable designers to model their solution accordingly.
- 6) The application of the developed method for modification consistently makes the flowcharts represent the recent models of the processes in the software.
- 7) If the initial flowcharts are developed properly, incorporation of any changes in the flowcharts into the software processes is mostly straightforward.

- 8) The developed method for modification does include all the three key phases of software development, namely analysis, design and implementation. The method is also iterative and incremental.

Previously when there was no model for debugging, the conceptual starting point to look for defective code was the first executable statement of the software. As one keeps working on the software, a good guess as to which module of the software may contain the bug can be made. To identify the problem causing code segment, one has to check all the possible defective code segments; this is mainly due to the message passing ability between various objects. Even once one finds the potential problem causing code segment, the probability that the first fix is the best and most effective fix is low, as the proposed fix may affect the underlying model as one does not know the entire model the software follows for its processes. However, once the model is established there is no need to go through the code first. Instead, one goes through the developed flowchart structure looking for the possible process model that may contain the bug. The process model can be easily guessed from the definition of the problem. Thus, the developed model is providing a relationship between the bugs and the PORTSIM process models. This relationship saves lot of time in finding the problem causing code. Once the process model is found, it is easy to pin point the problem causing structure and thus the code. Once the problem causing structure is found, one can propose the necessary changes to fix the bug without affecting the underlying model. It is the total picture of the underlying related process models that enables the bugs to be fixed effectively.

## 5.2 Drawbacks and Possible Enhancements

This section discusses drawbacks in the method and possible solutions.

- 1) The effectiveness of the method depends on how well the initial flowcharts were developed for each process. Conceptually, one could develop an automated method for abstracting the flowchart structure from the code and then allow the developer to fill in the resulting flowchart constructs according to the code.
- 2) Process models were developed when the requirement to work on that process came, so complete understanding of all the processes in the software is not available. At present, one does not know the complete Graphic User Interface (GUI) model and the complete memory semantics within the software. To solve this problem, one can develop models for all the processes within the software.
- 3) The flowchart constructs developed for MODSIM III language constructs do not include all the available MODSIM III constructs, like monitors, thus restricting the designers to the subset of available MODSIM III constructs. Development of flowchart constructs after a thorough study of the MODSIM III language manual can eliminate these restrictions.
- 4) There is currently no direct connection between the code and the flowcharts to further facilitate the code identification process. Inserting comments around code segments would provide a link between the code and the flowchart constructs. The comments should refer to the particular flowchart construct that is representing the code segment.

The developed model was useful in making modifications to the PORTSIM code by abstracting documentation of the software in the form of process flow. This helps to locate bugs in the code easily. The documentation captures the language constructs inherent to MODSIM, thus acting as a common platform between designers and programmers. The documentation helps to understand the underlying models of various processes and also the relation between these processes. This helps to make modifications in the code to fix bugs effectively. The model is currently being used by the Virginia Modeling, Analysis and Simulation Center (VMASC) and Argonne National Labs to continue PORTSIM development.

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## APPENDIX A

### PORTSIM PROCESS FLOW

As a whole PORTSIM can be divided into three phases:

- 1) An initialization phase.
- 2) A simulation phase.
- 3) An output phase.

Fig. A.1 represents the program flow of PORTSIM at macro level.

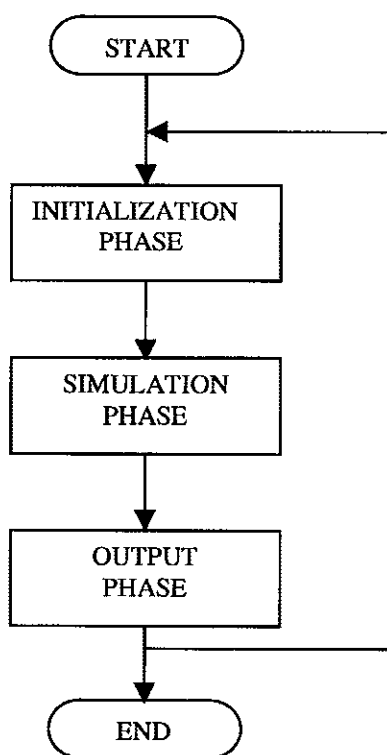


Fig. A.1

### **A.1 Initialization phase**

It is in this phase that we prepare PORTSIM to run the scenario we want. For simplicity and ease of understanding, this section deals only with POE-Standalone type. We can either create a new scenario in which PORTSIM relies heavily on the underlying database to get required data or load an existing scenario in which PORTSIM reads data from the scenario file. To clearly understand all parts of initialization phase, we will look into the process of creating a new scenario.

As soon as we start PORTSIM, it is connected to the underlying database. When we decide to create a new scenario, PORTSIM provides a dialog box, which guides us in creating a new scenario. The various parameters that we have to decide upon are a scenario name, a port and a terminal name, a force name and a list of ships upon which the selected force has to embark. Fig. A.2 represents the flow of program control within PORTSIM during the creation of new scenario.

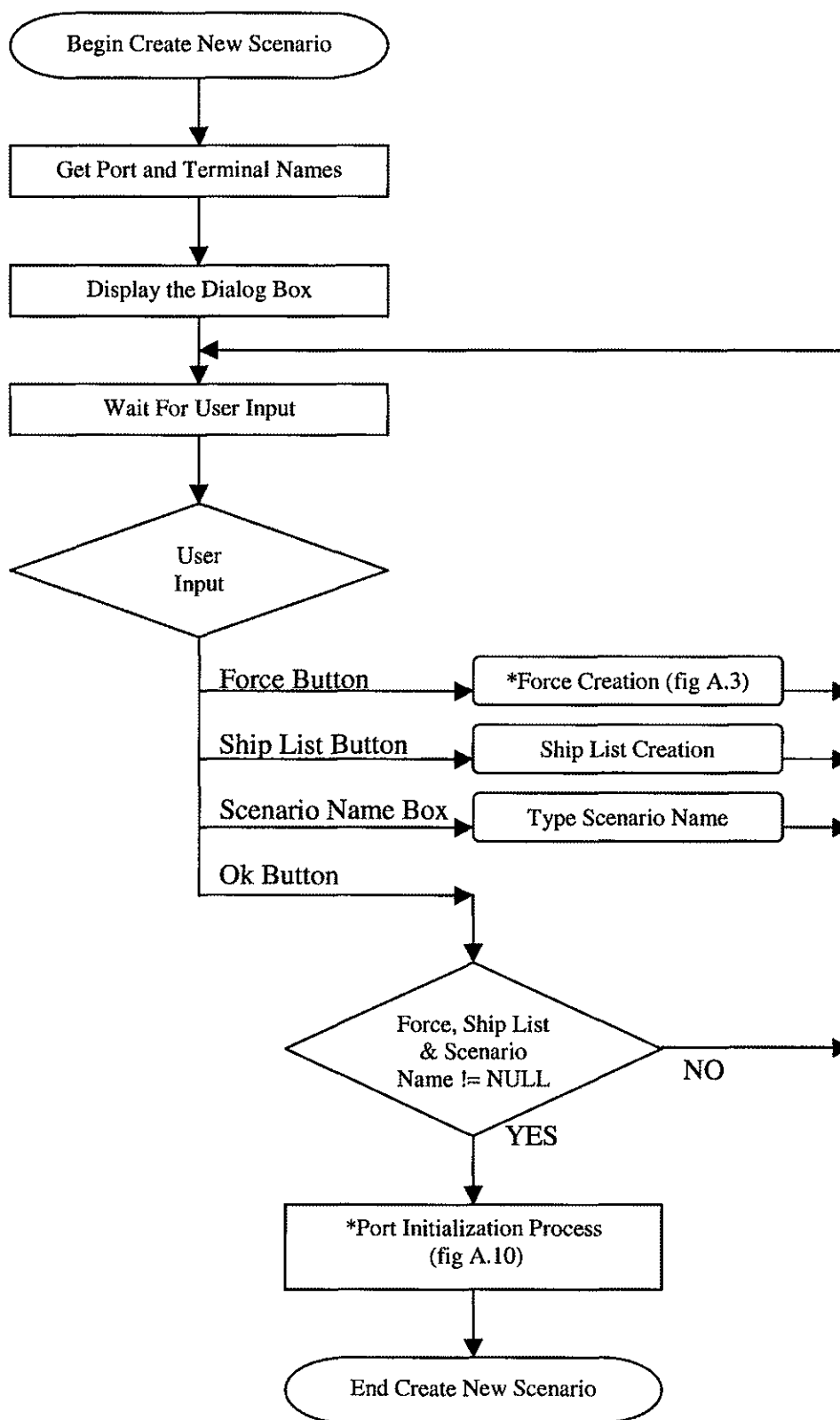


Fig. A.2

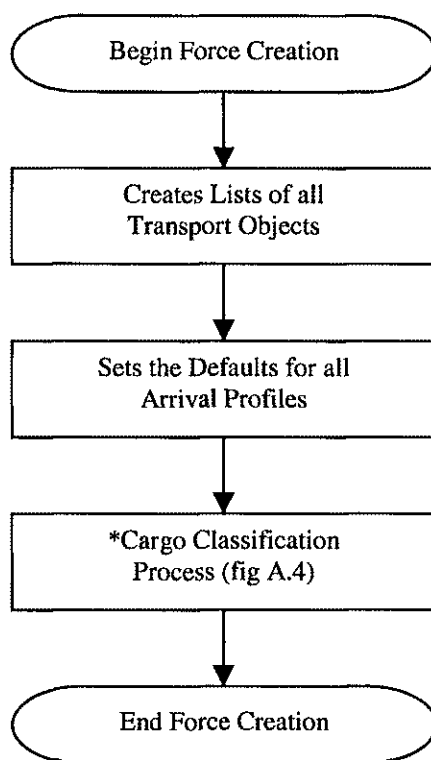


Fig. A.3

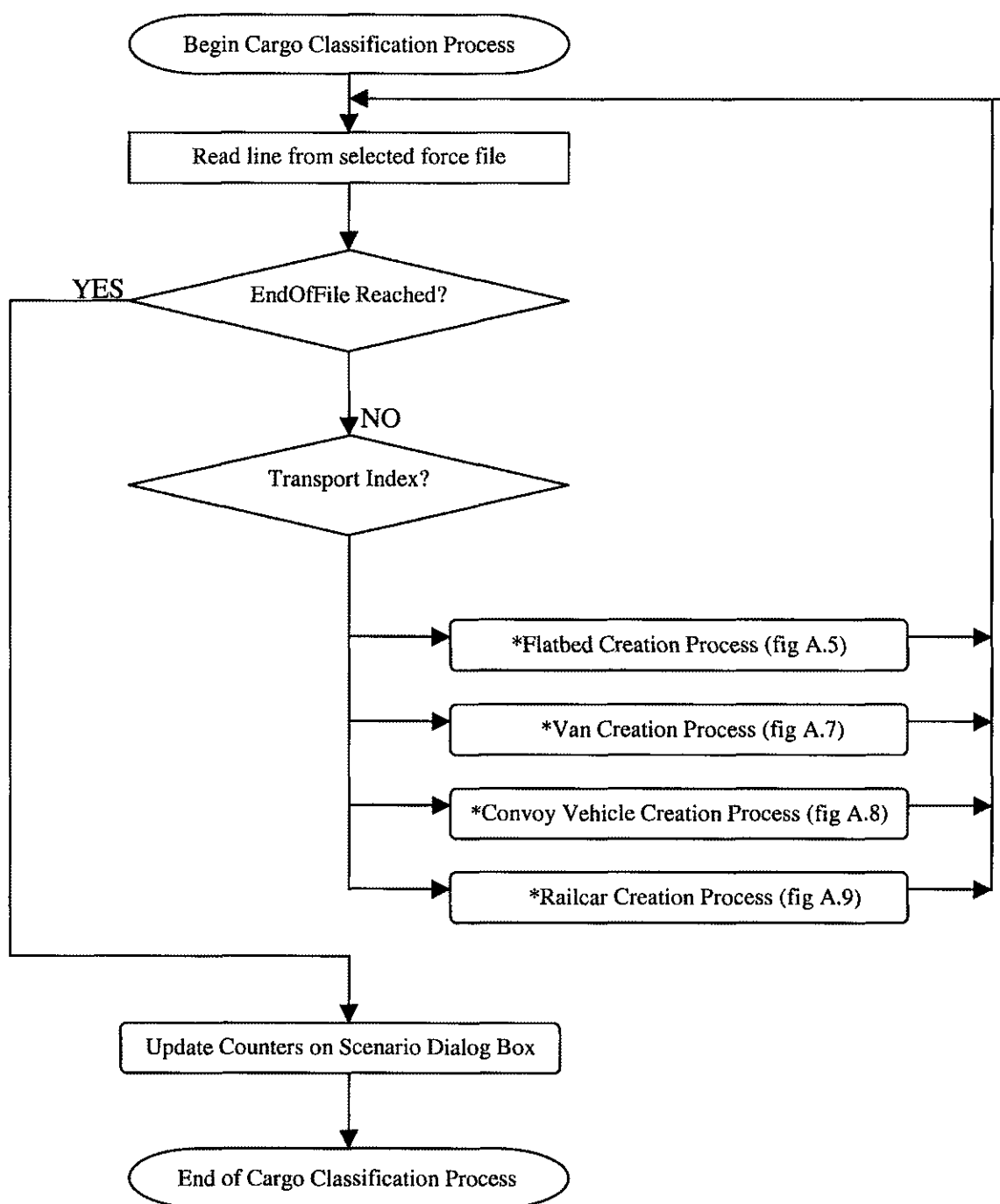


Fig. A.4

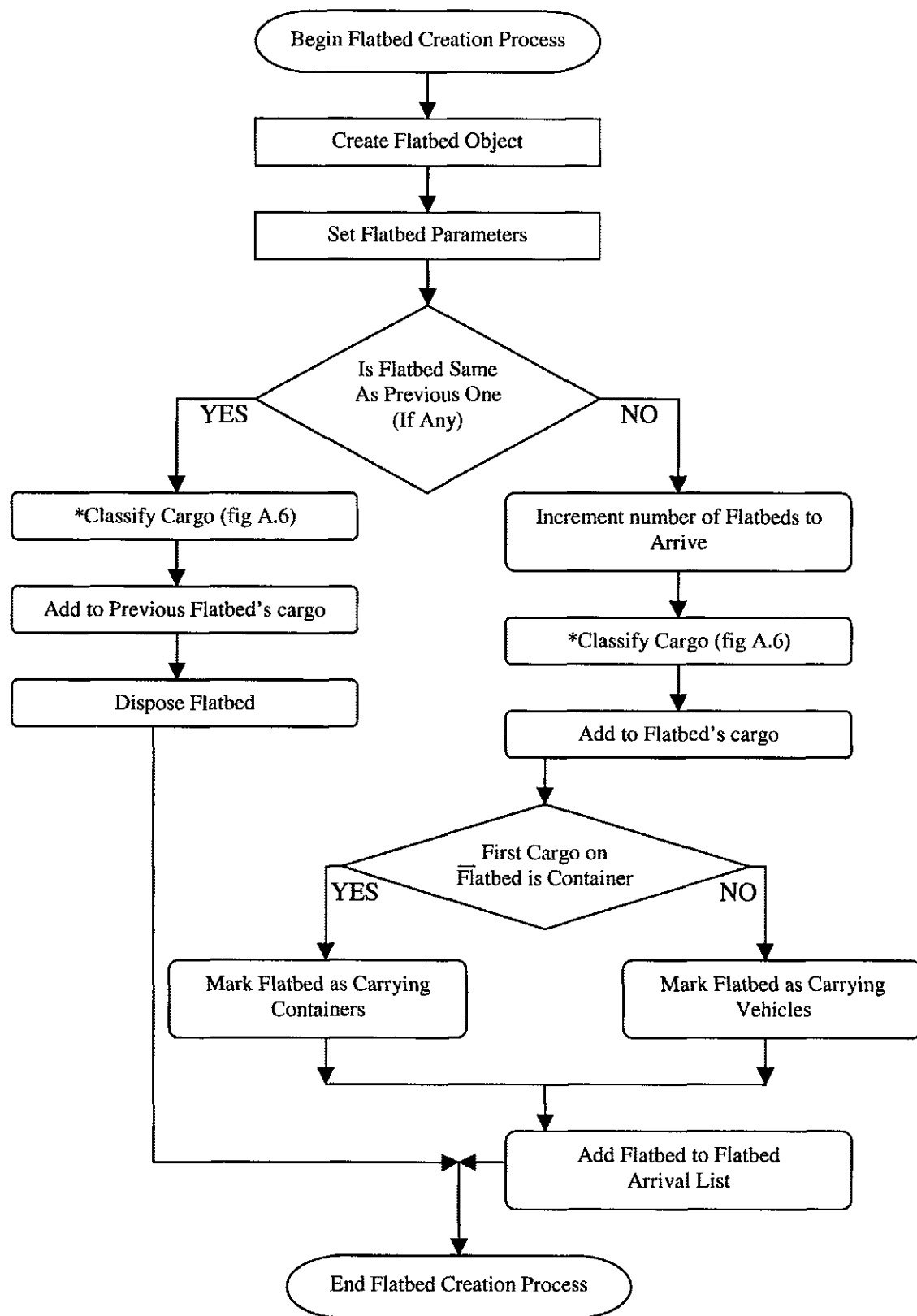


Fig. A.5

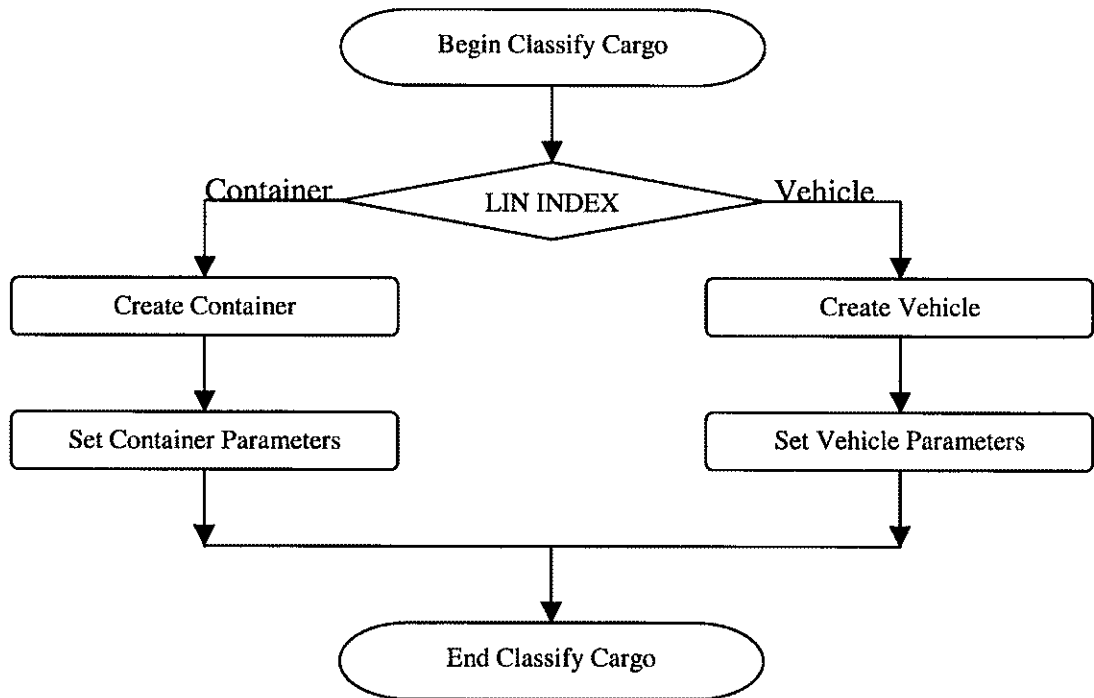


Fig. A.6

Fig. A.6 shows how containers and vehicles are differentiated in PORTSIM. This type of index-based classification is used for vehicles and containers only and not for pallets, as both flatbeds and railcars are capable of carrying both vehicles and containers. Pallets are carried by vans only.

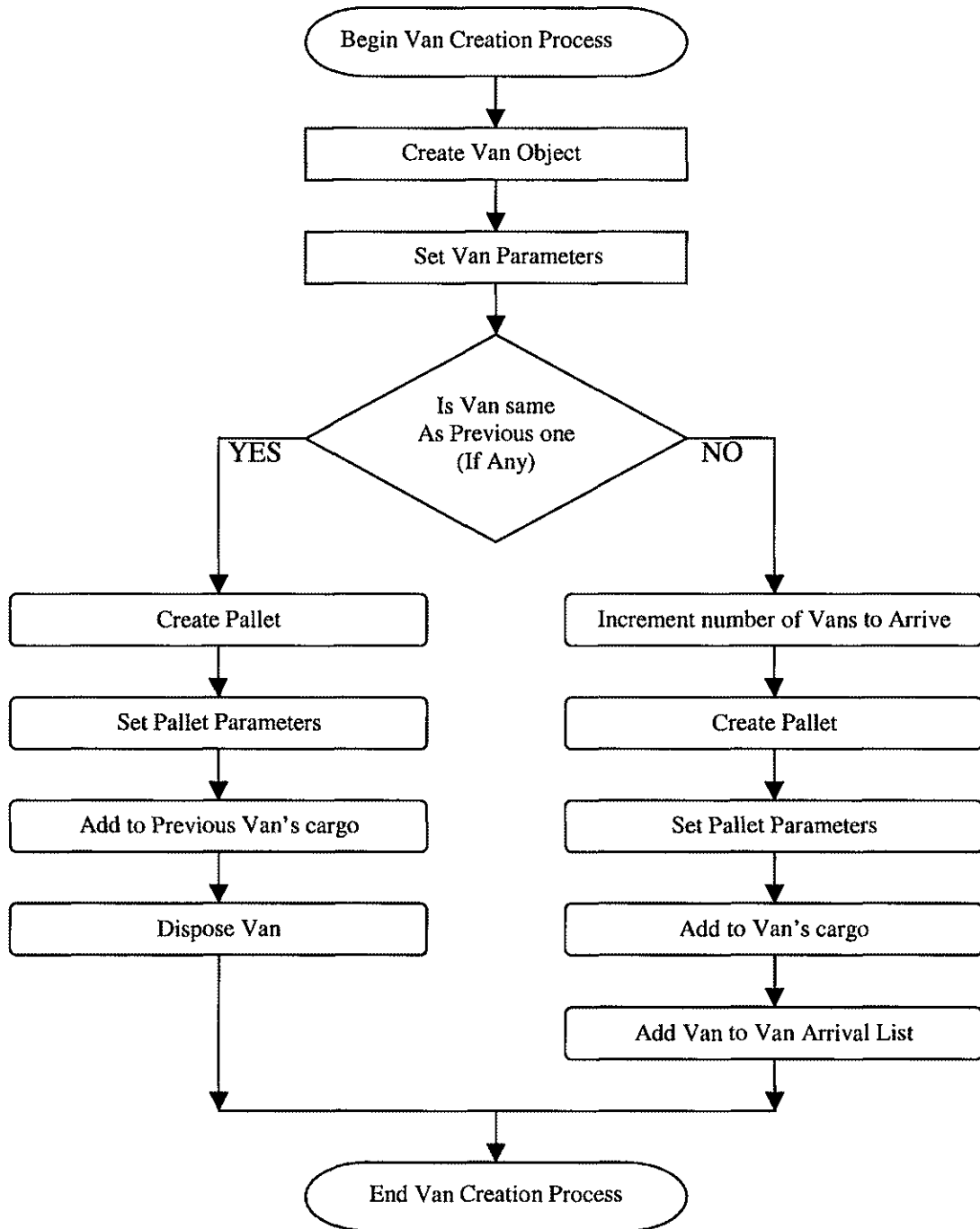


Fig. A.7



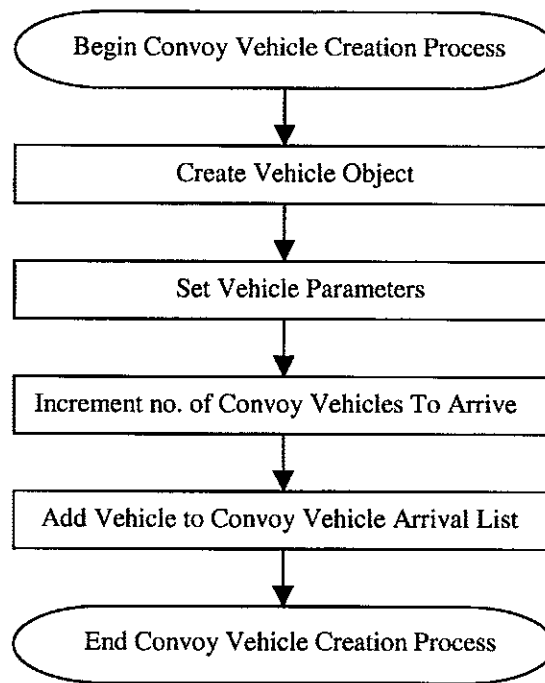


Fig. A.8

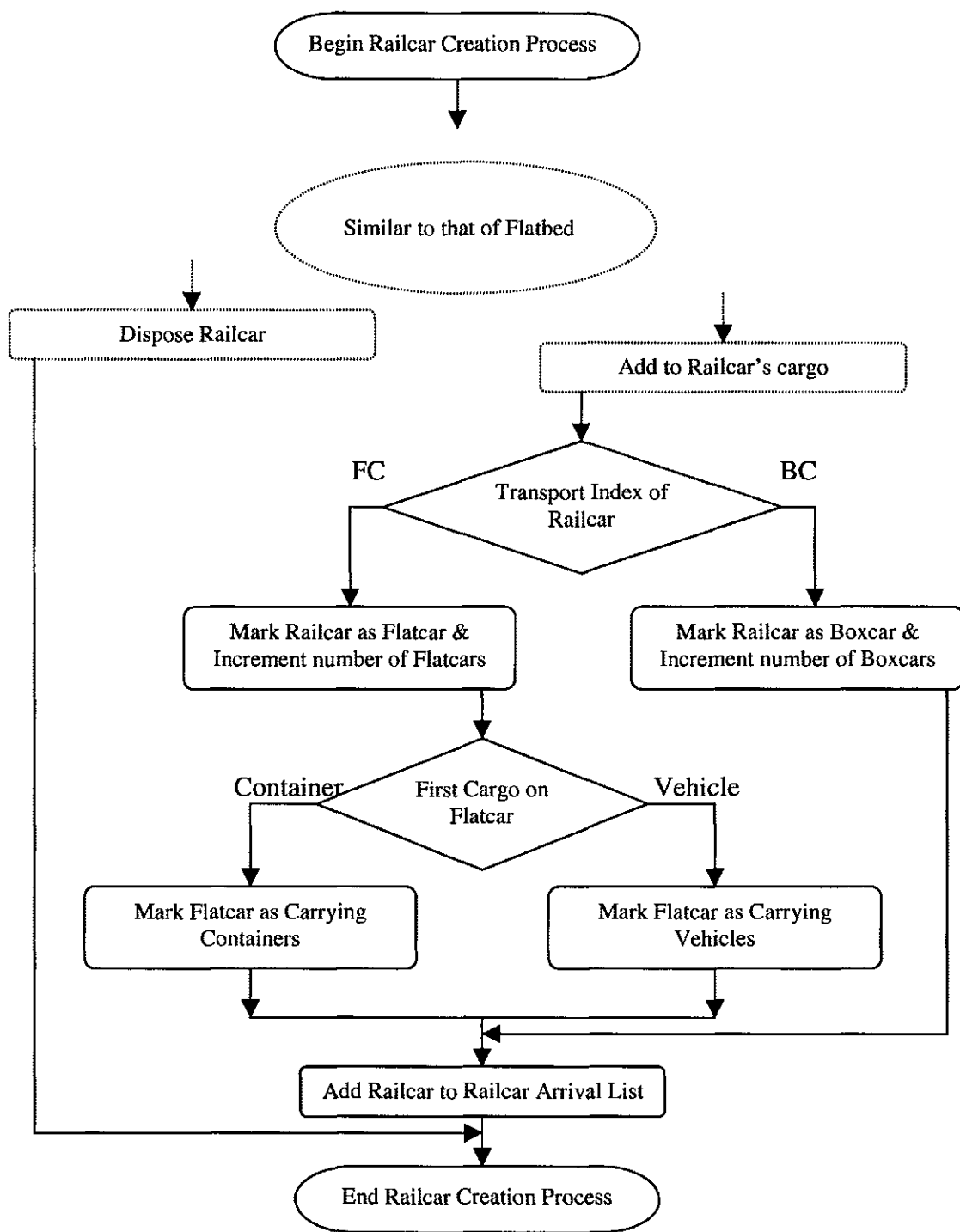


Fig. A.9

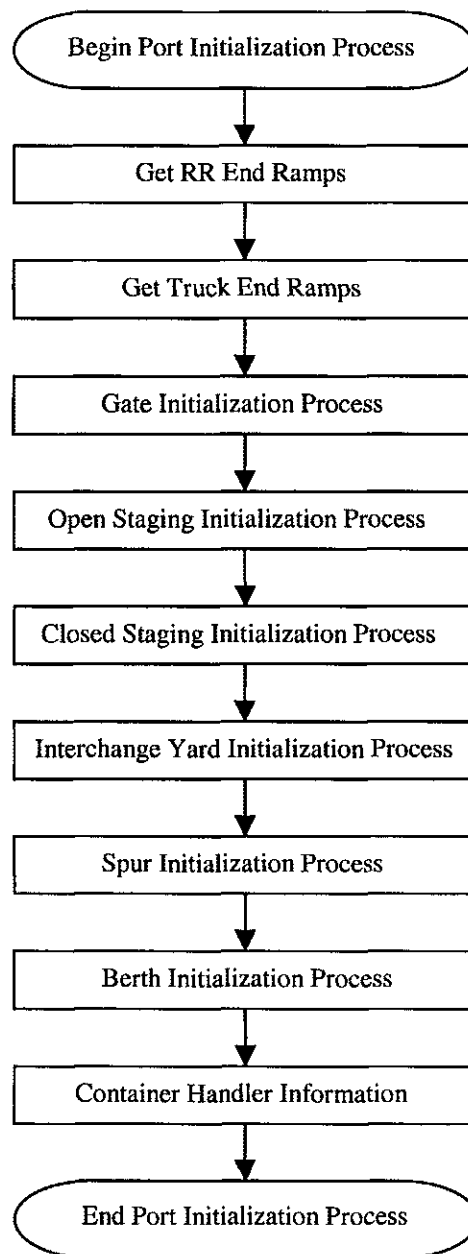


Fig. A.10

Each process mentioned above can be generalized as shown in Fig. A.11.

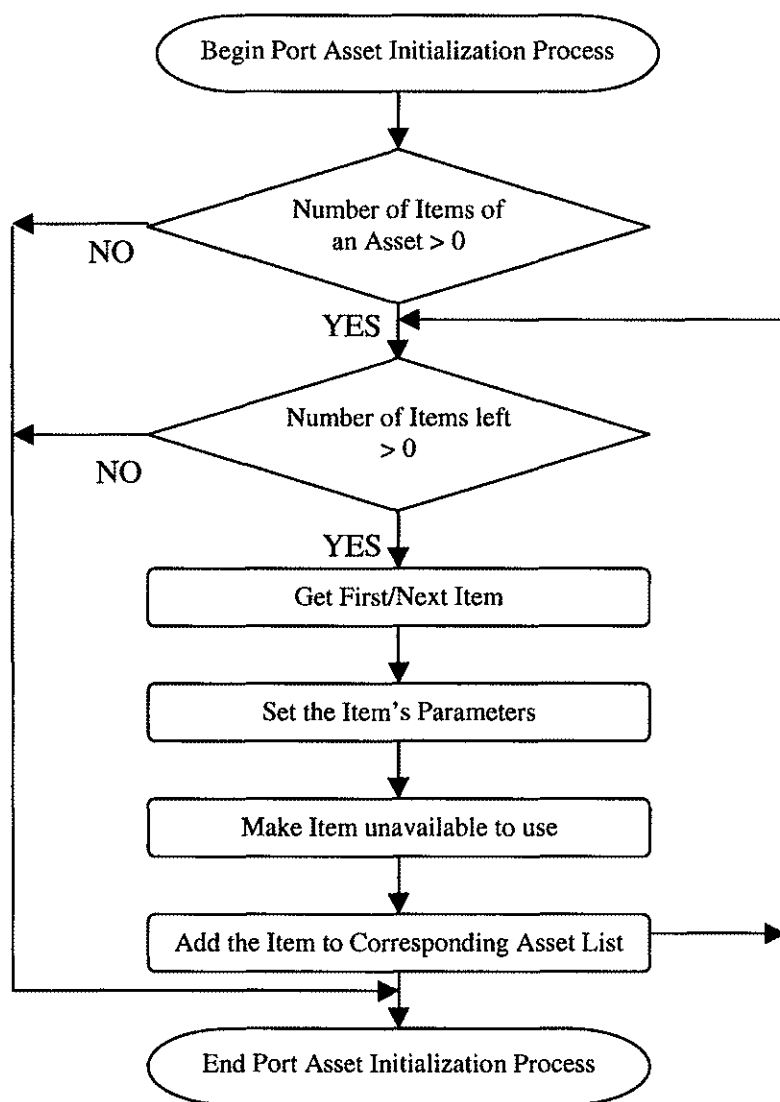


Fig. A.11

NOTE: After all of this initialization is done, PORTSIM checks if there is any other scenario with same name. If there is, it asks us whether to overwrite the previous scenario or not. Now comes the important part of this phase. As we can see from Fig. A.11, all the port assets are made unavailable for military use. It is our duty to decide which assets are

to be made available for military use. After this, PORTSIM is ready to simulate the created scenario.

## A.2 Simulation Phase

It is in this phase that we watch the simulation run to completion. In most cases, there is little interaction with the software. Fig. A.12 shows the sequence of processes that PORTSIM goes through during this phase.

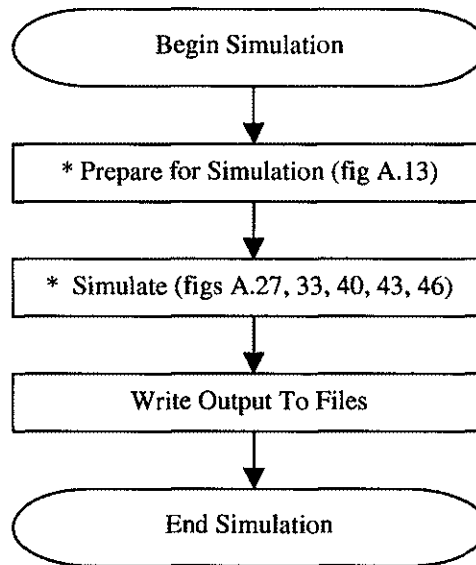


Fig. A.12

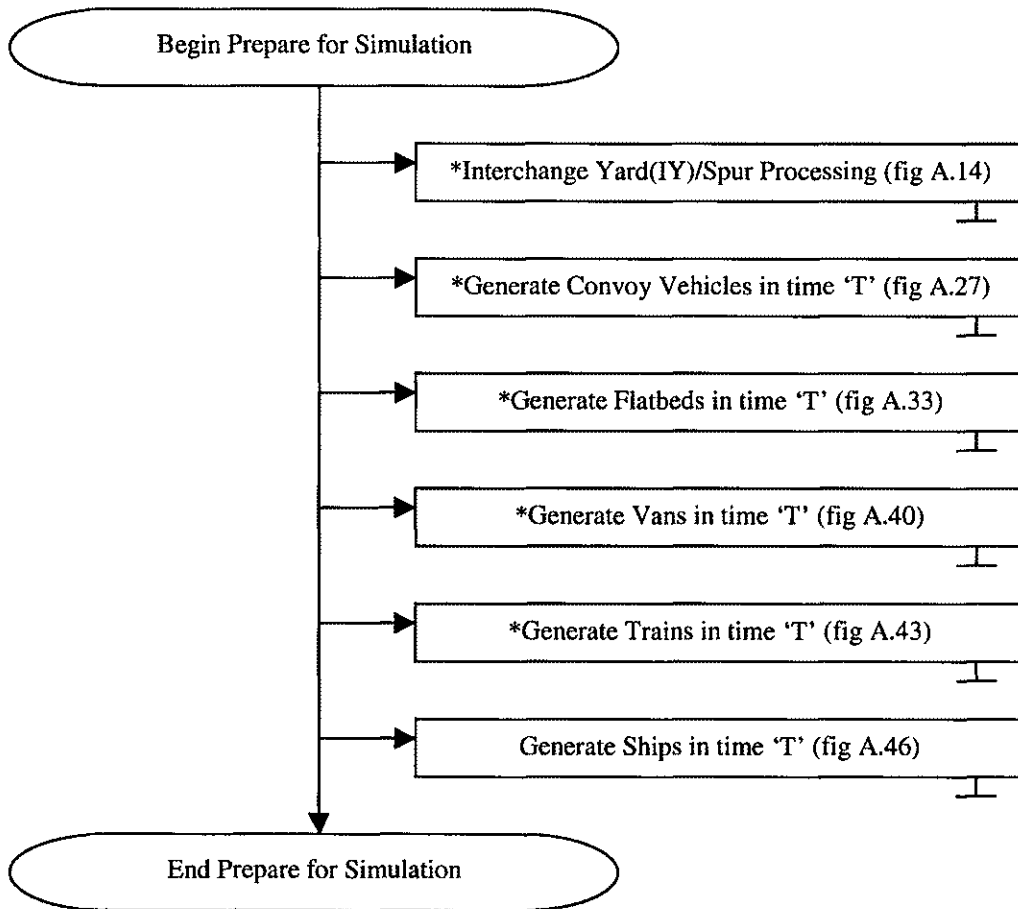


Fig. A.13

The time 'T' in the above figure is with respect to the start of simulation time. If the value of 'T' is one, that method is executed one minute after the start of simulation.

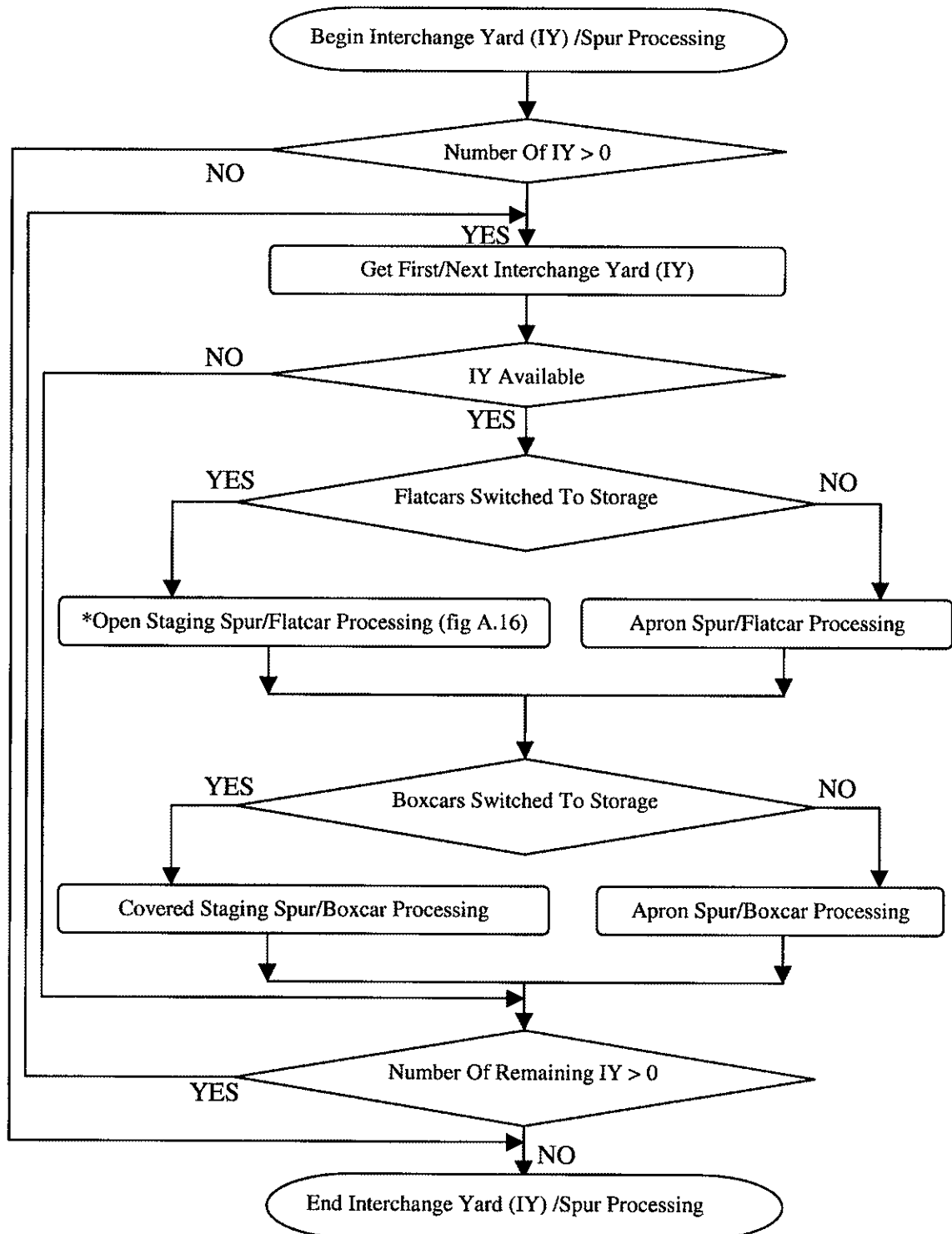


Fig. A.14



All the above mentioned 'spur type/railcar type processing' can be generalized as shown in Fig. A.15.

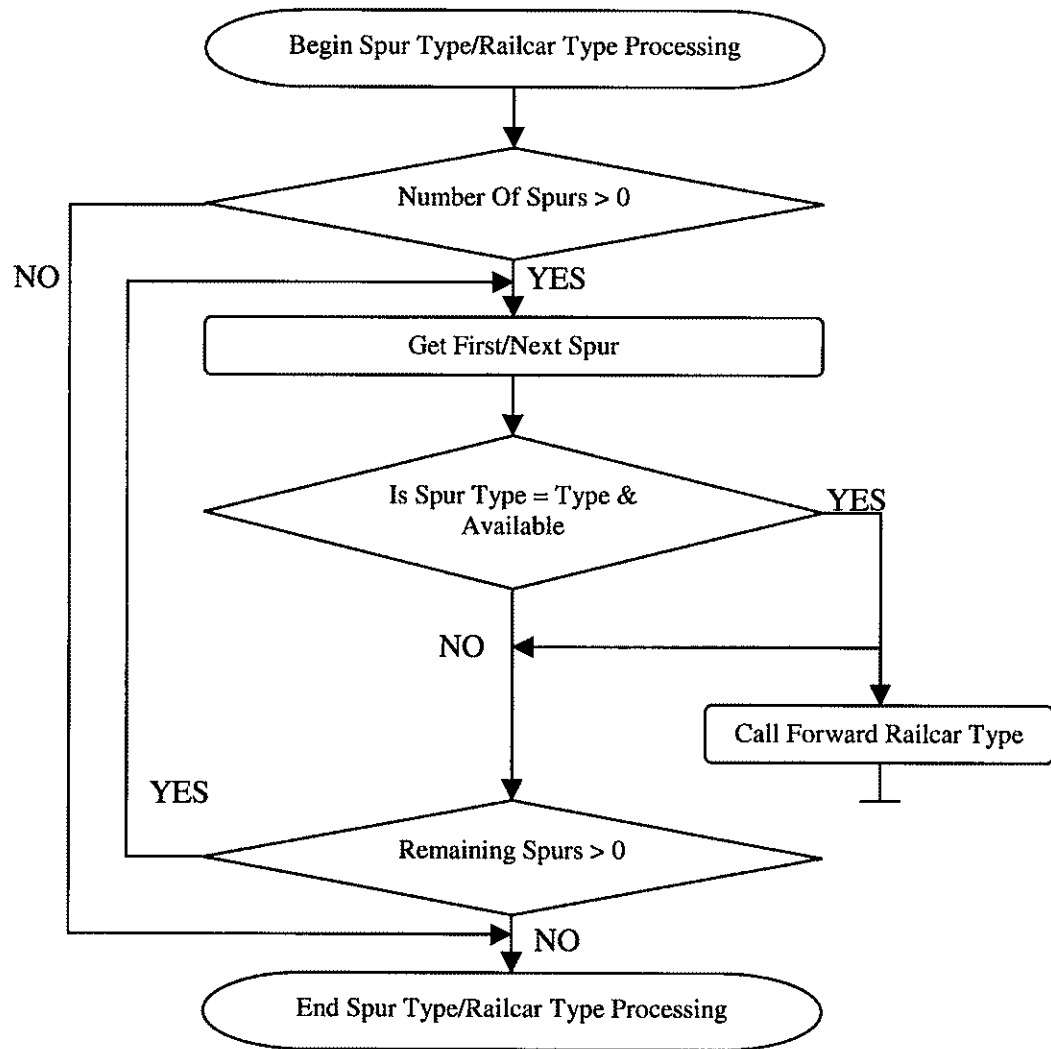


Fig. A.15

As an example, 'Open Staging Spur/Flatcar Processing' can be represented as shown in the following figure.

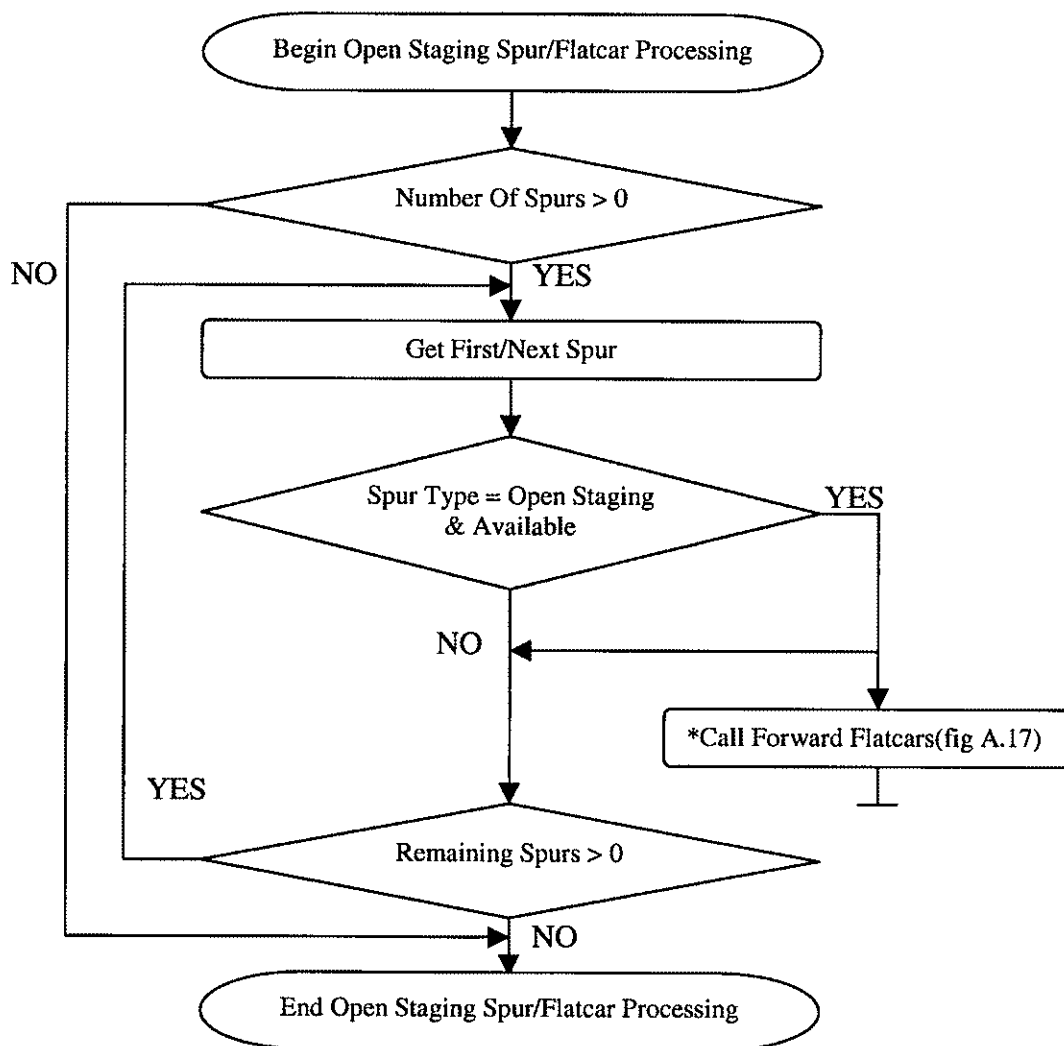


Fig. A.16

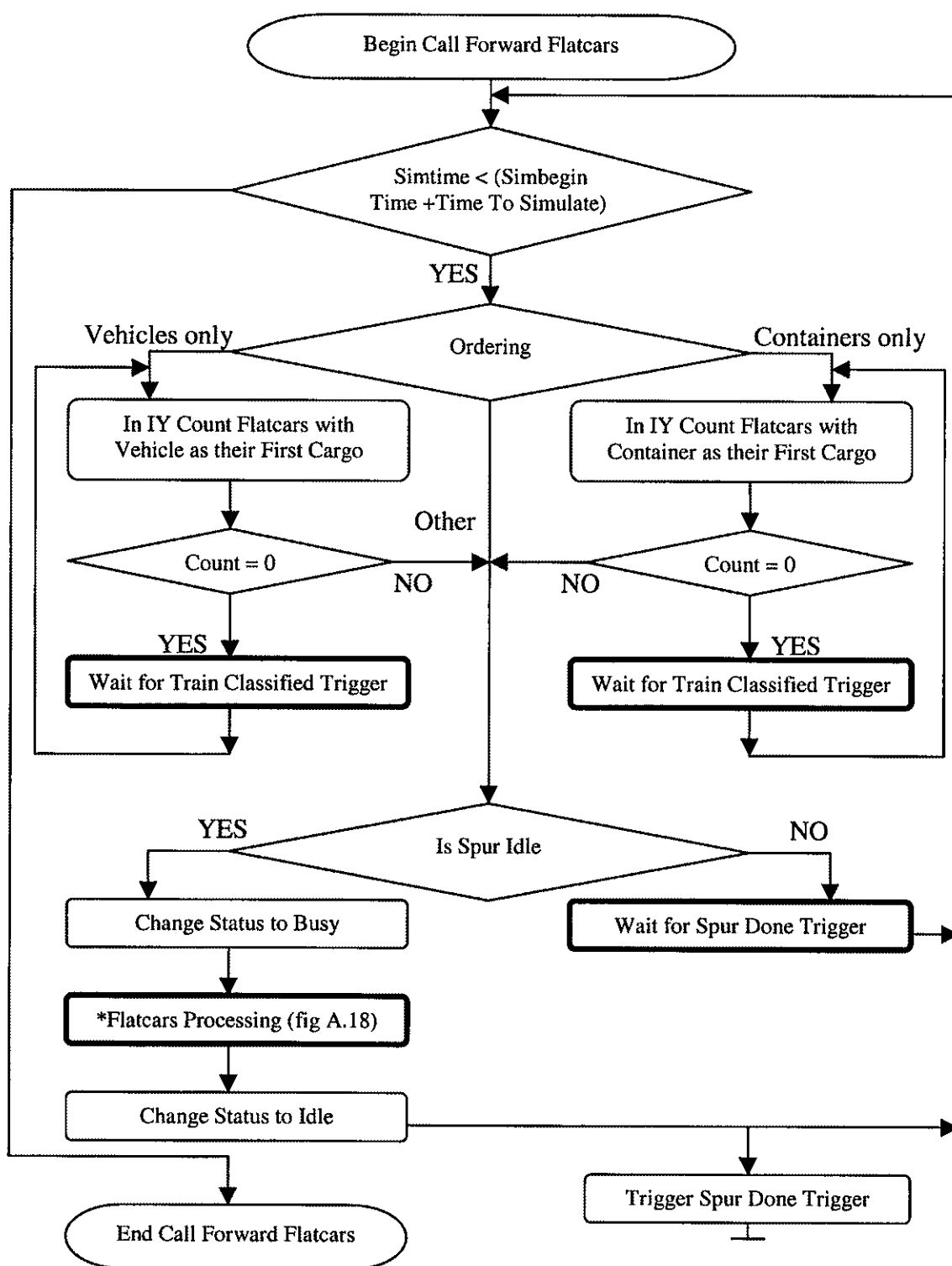


Fig. A.17

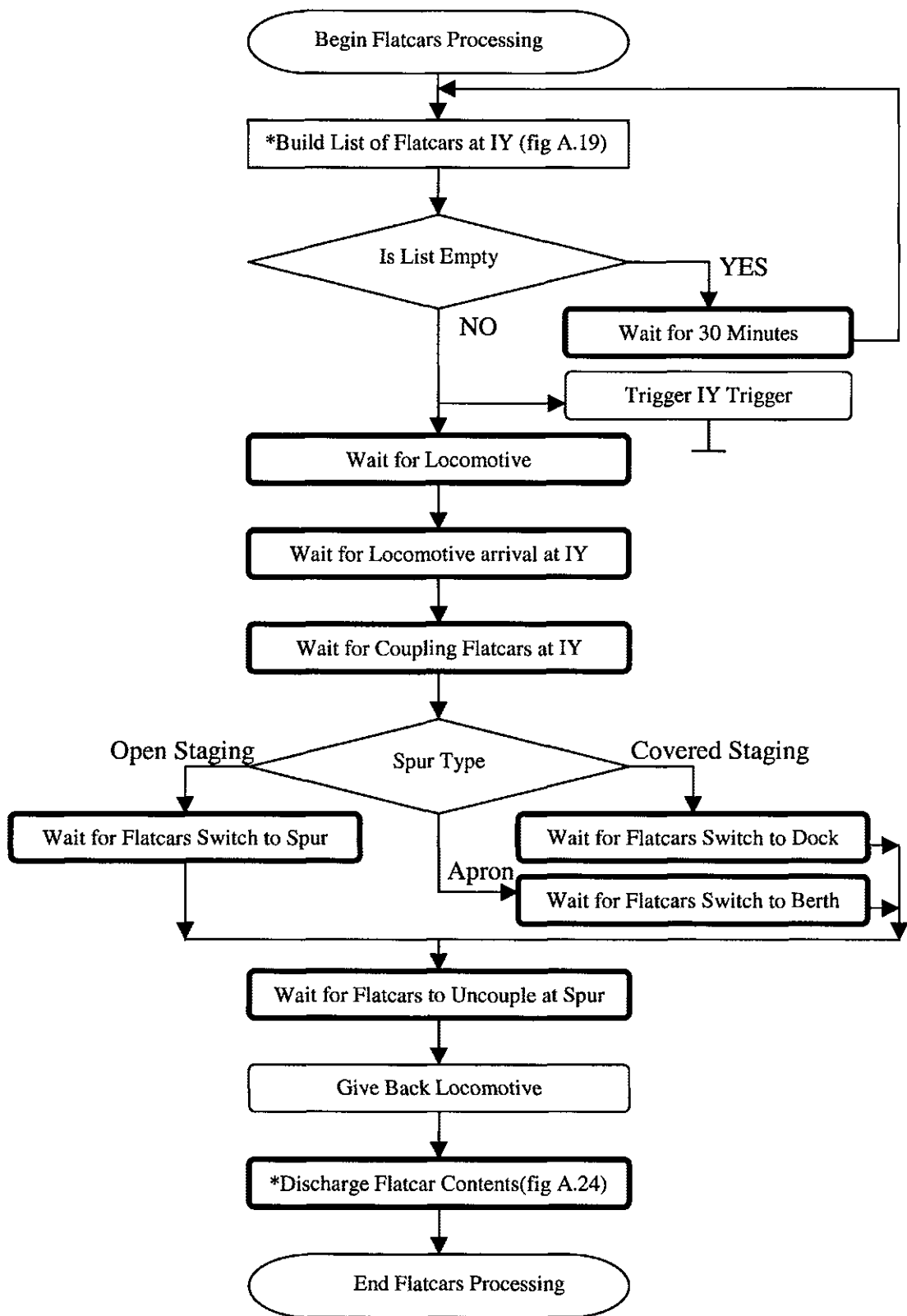


Fig. A.18

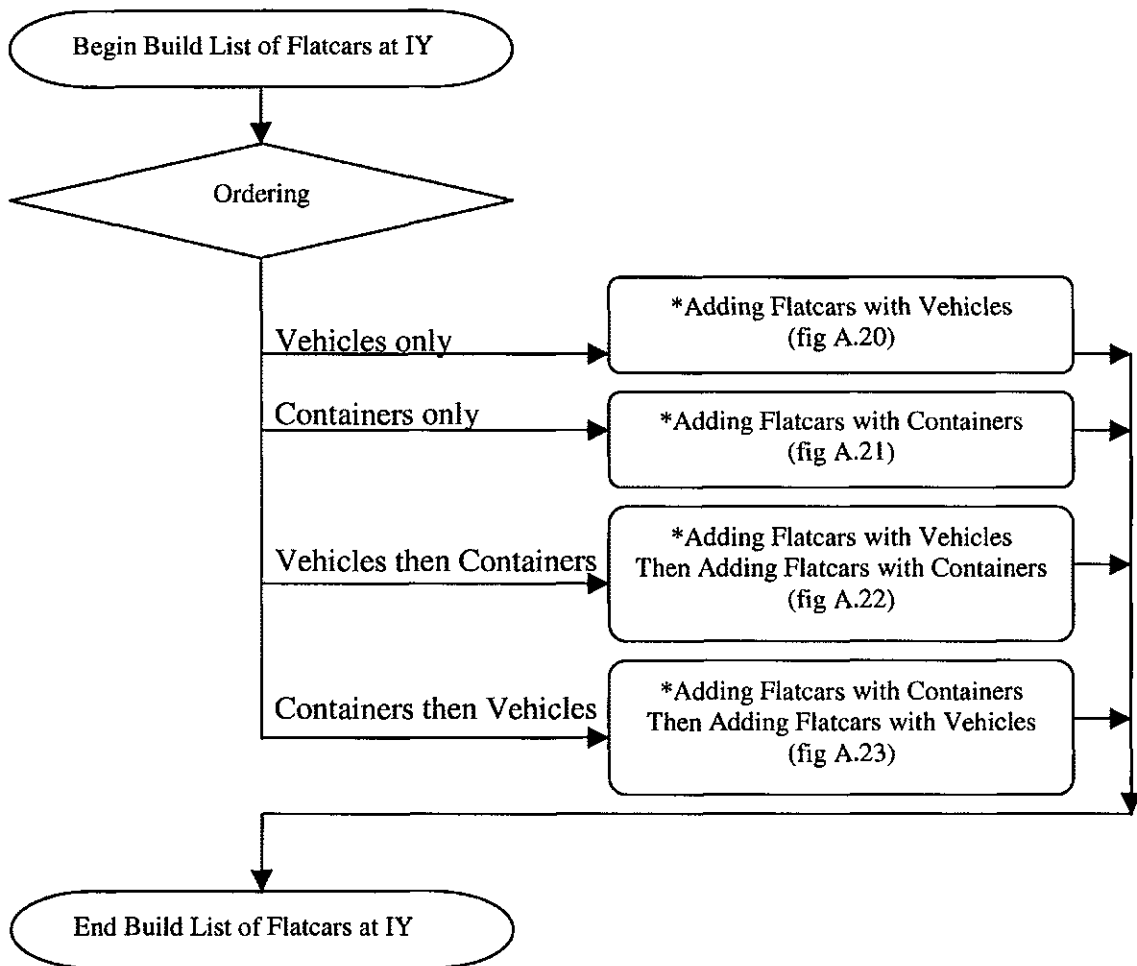


Fig. A.19

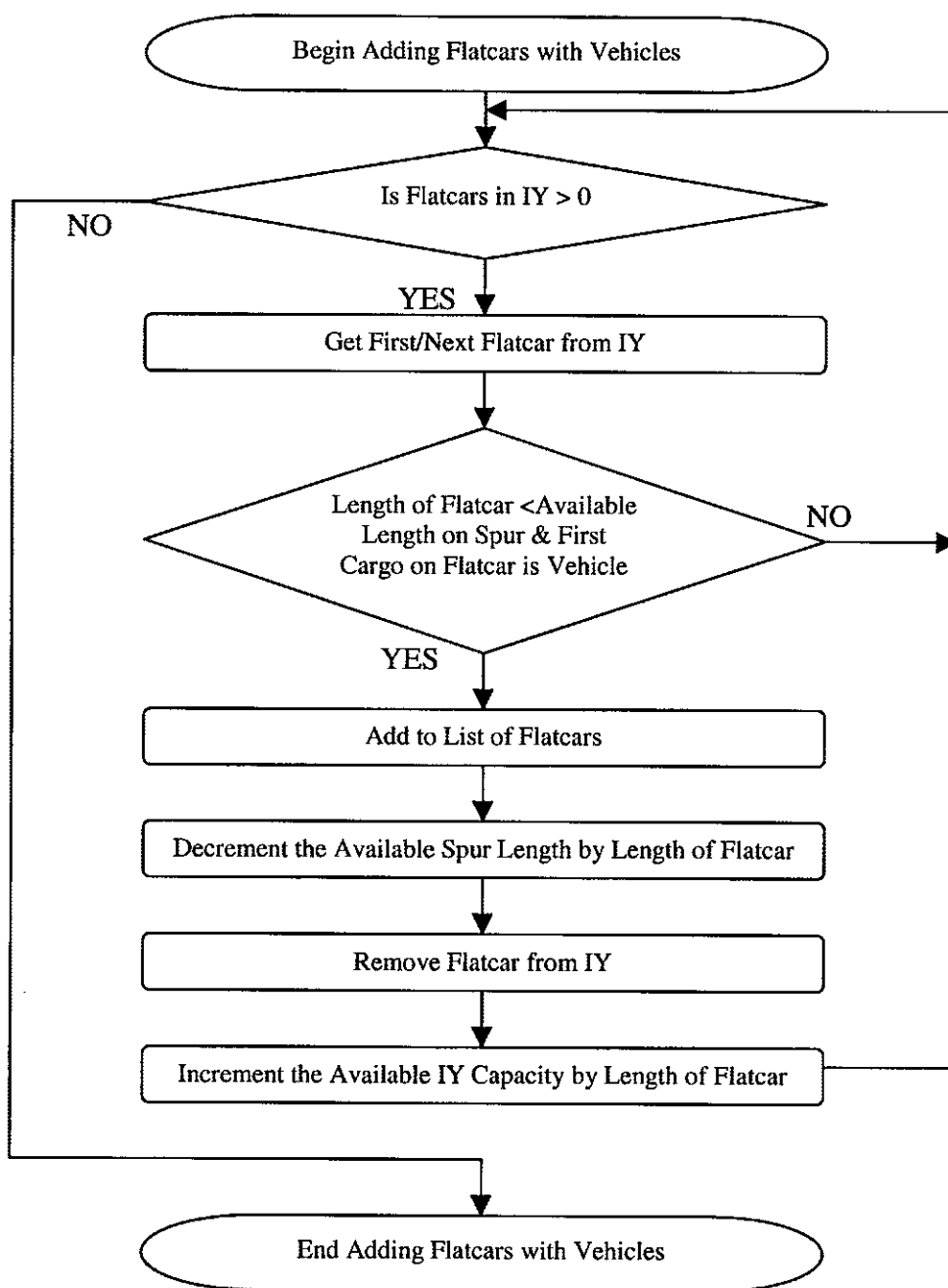


Fig. A.20

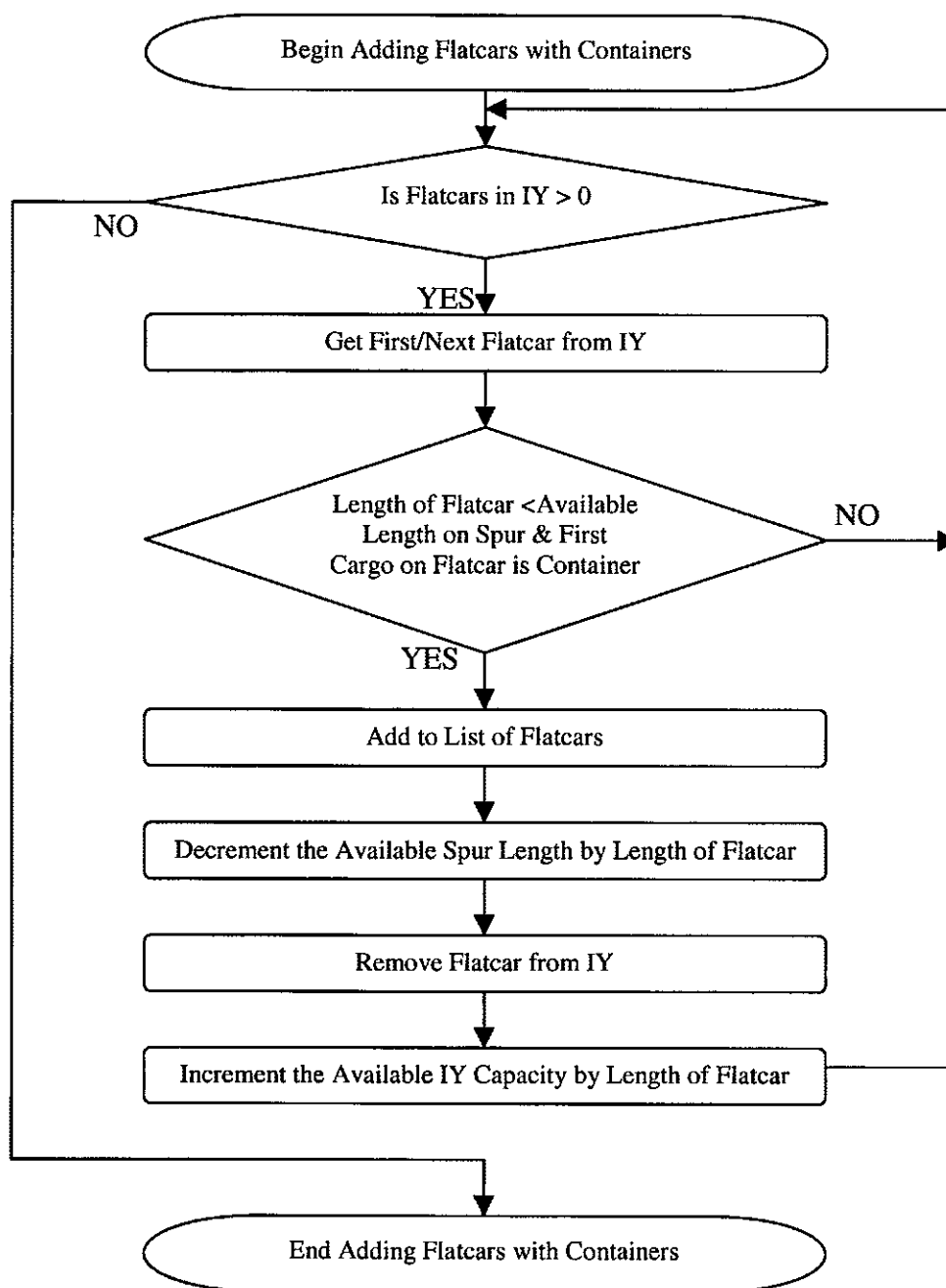


Fig. A.21

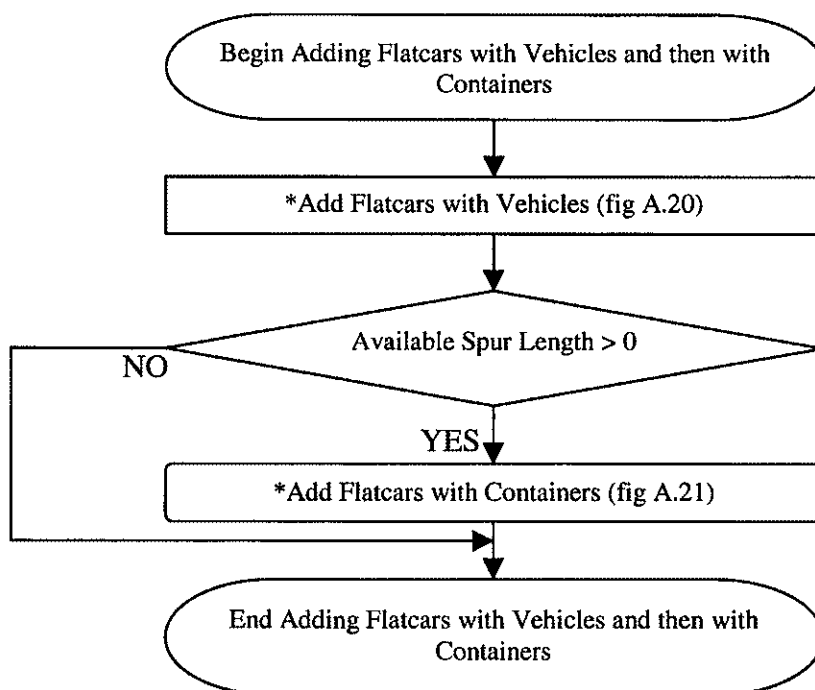


Fig. A.22

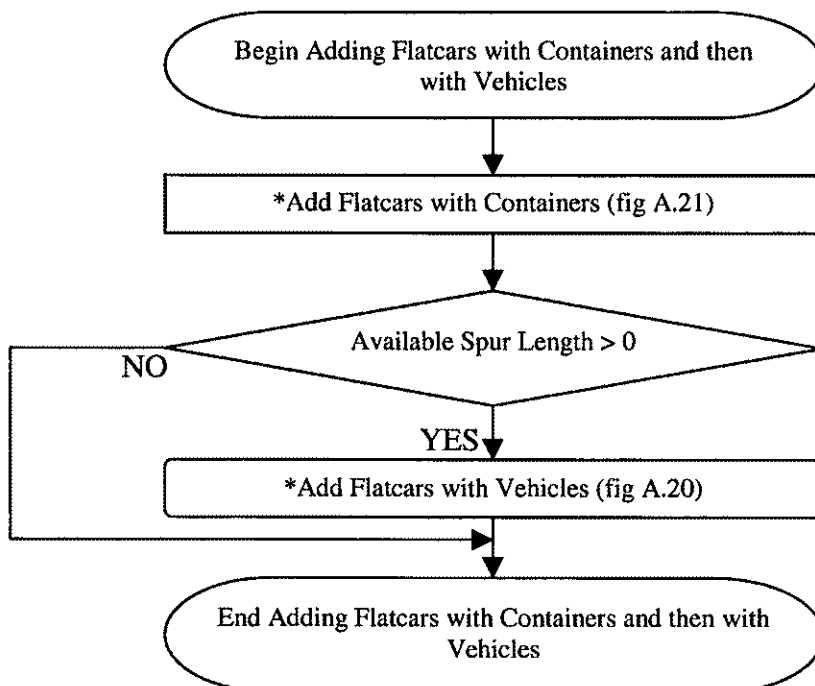


Fig. A.23



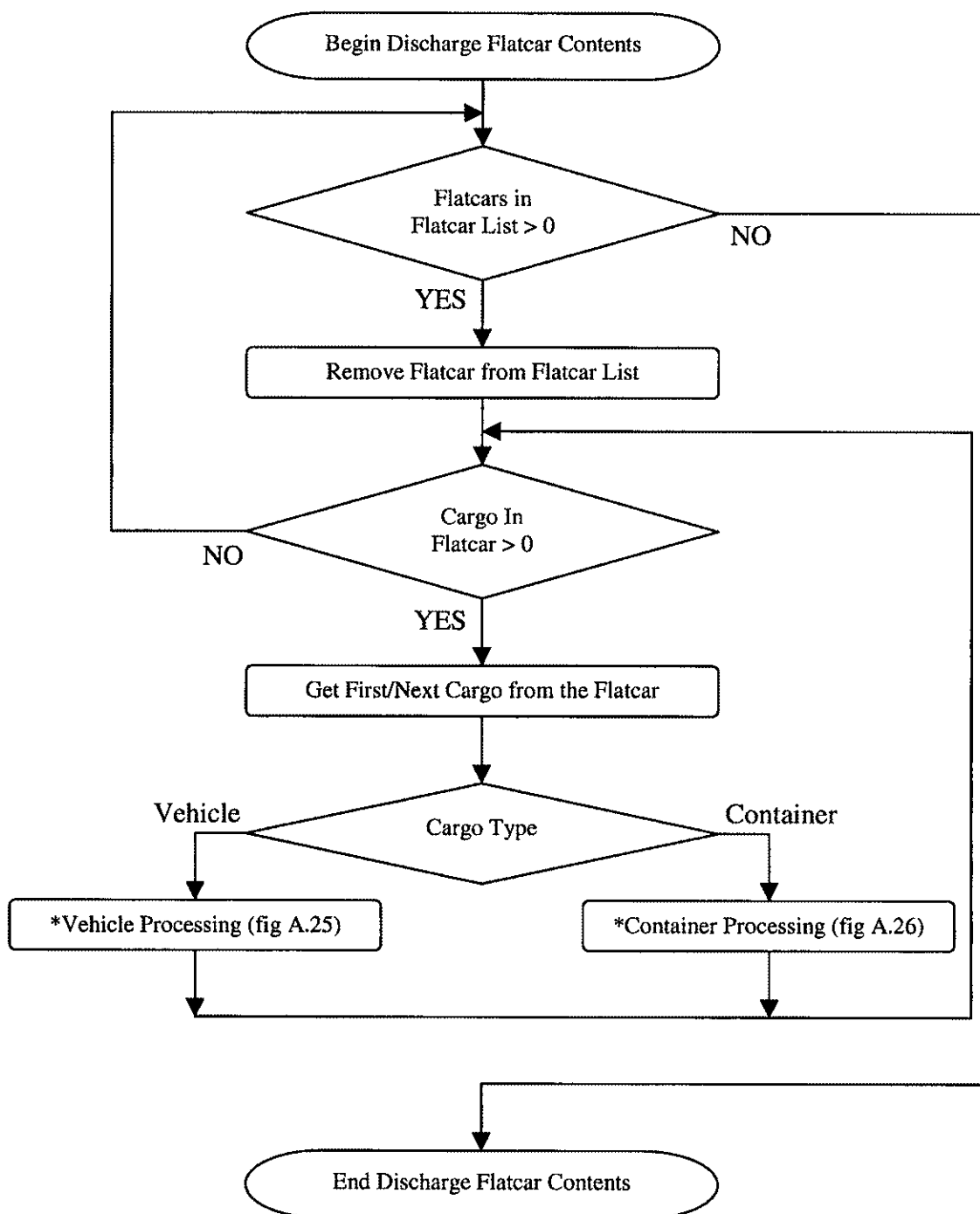


Fig. A.24

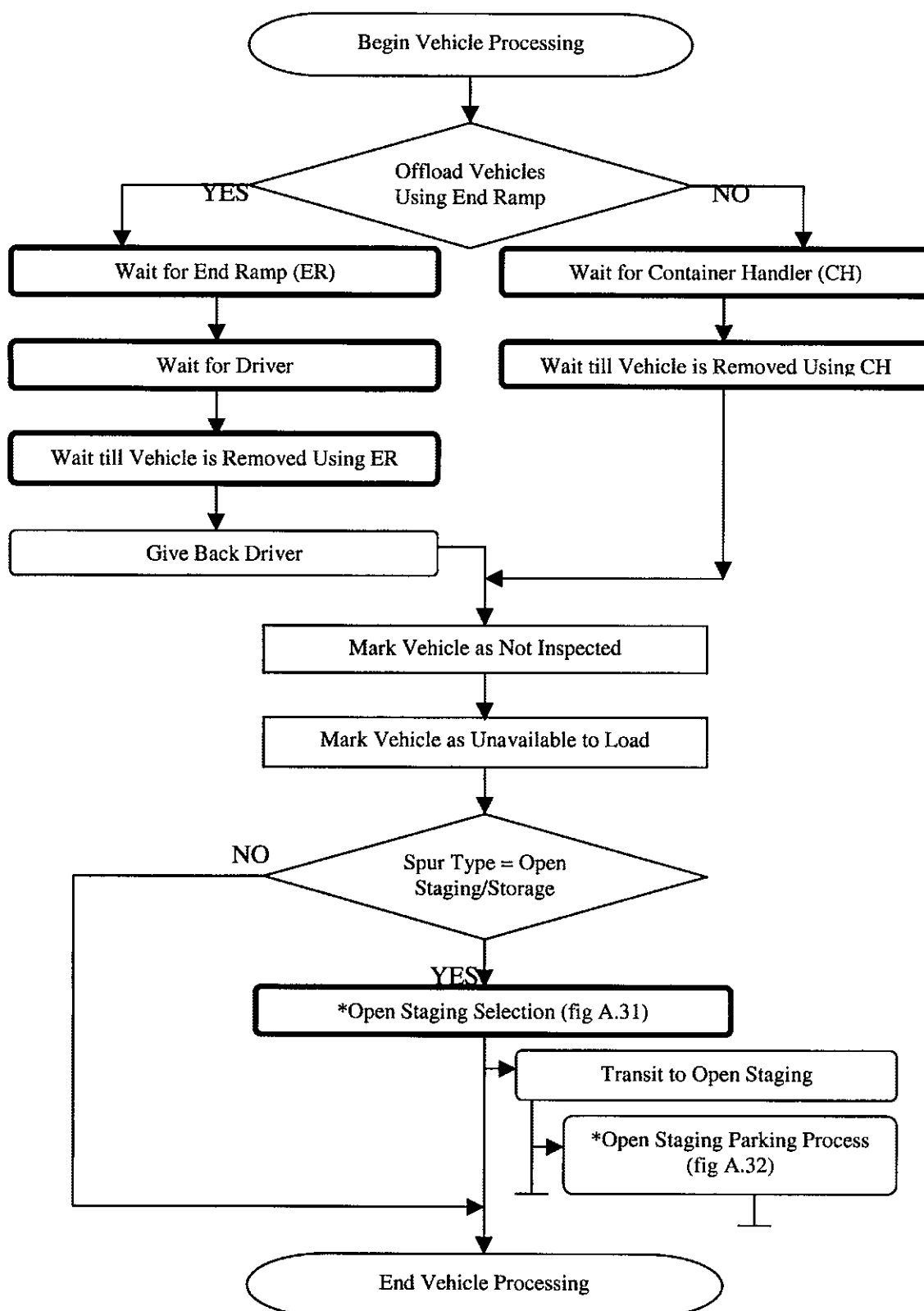


Fig. A.25

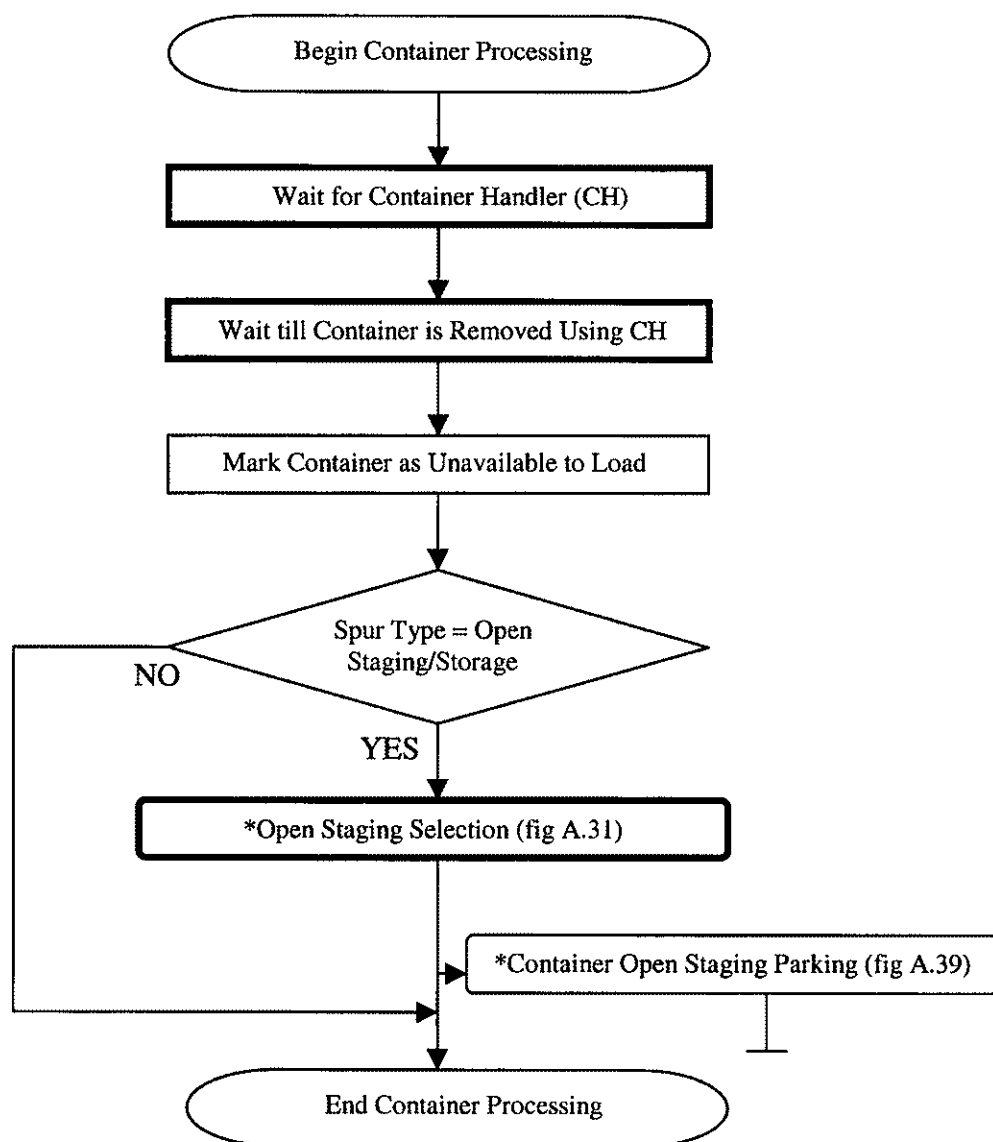


Fig. A.26

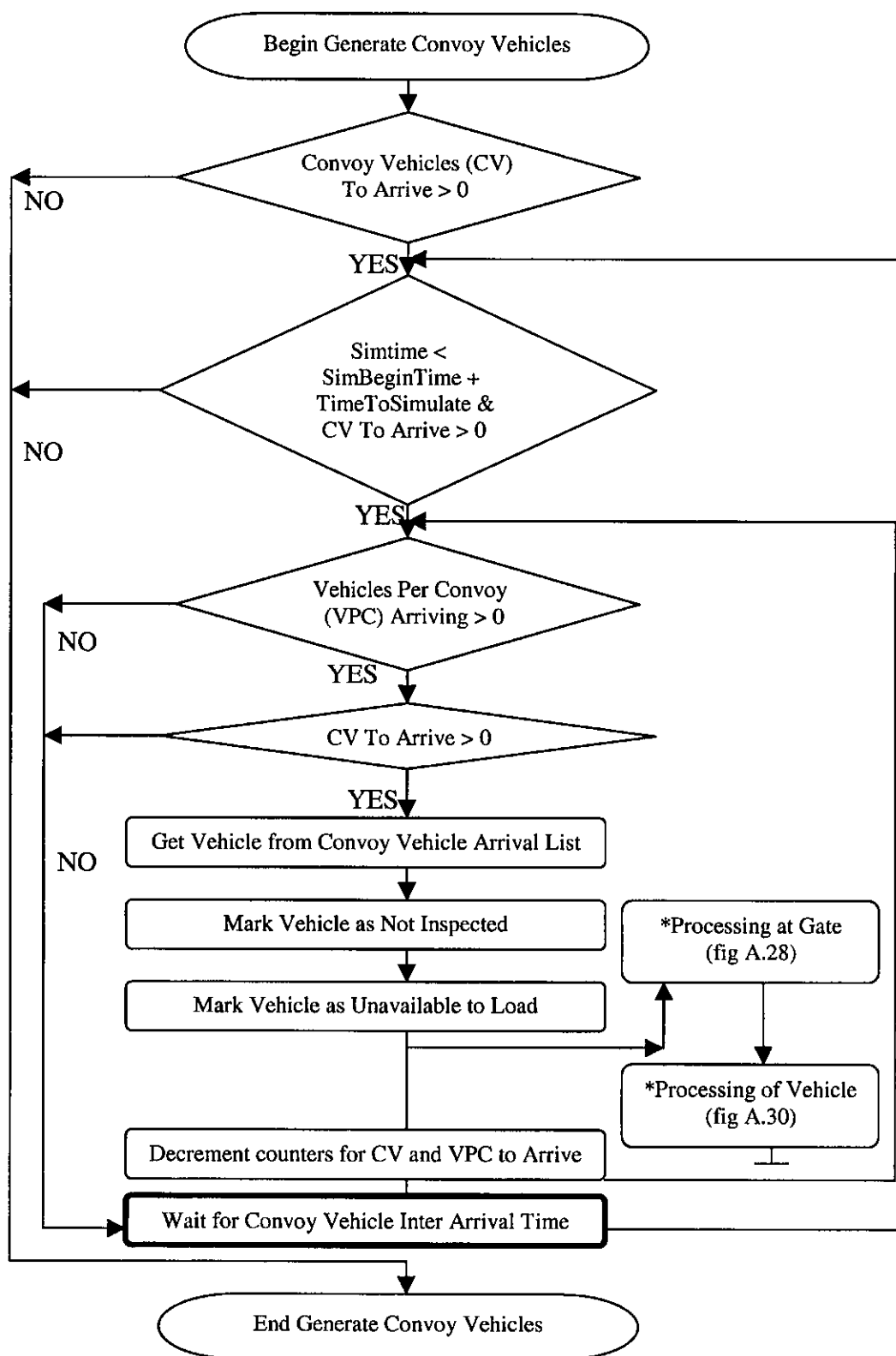


Fig. A.27

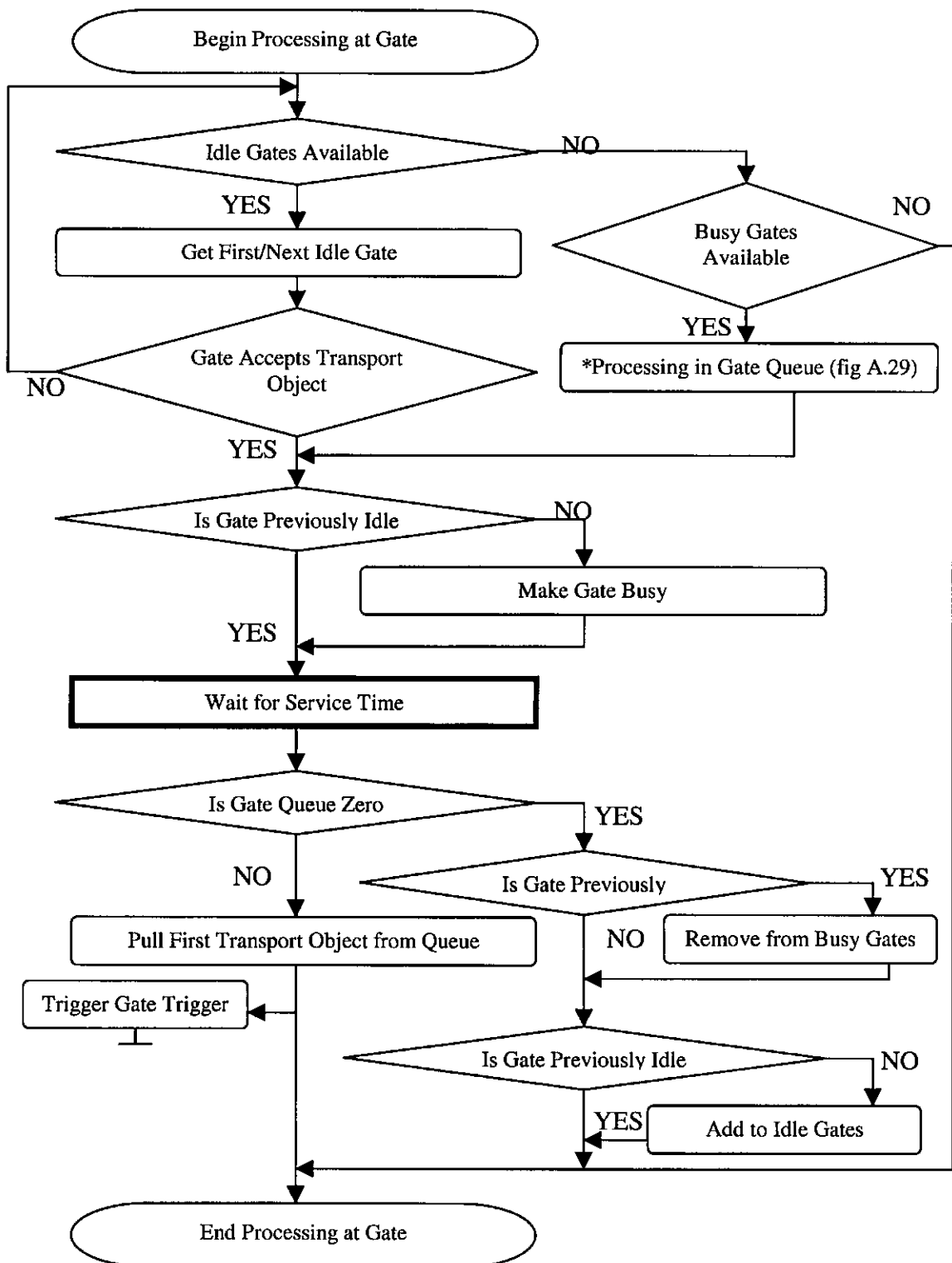


Fig. A.28

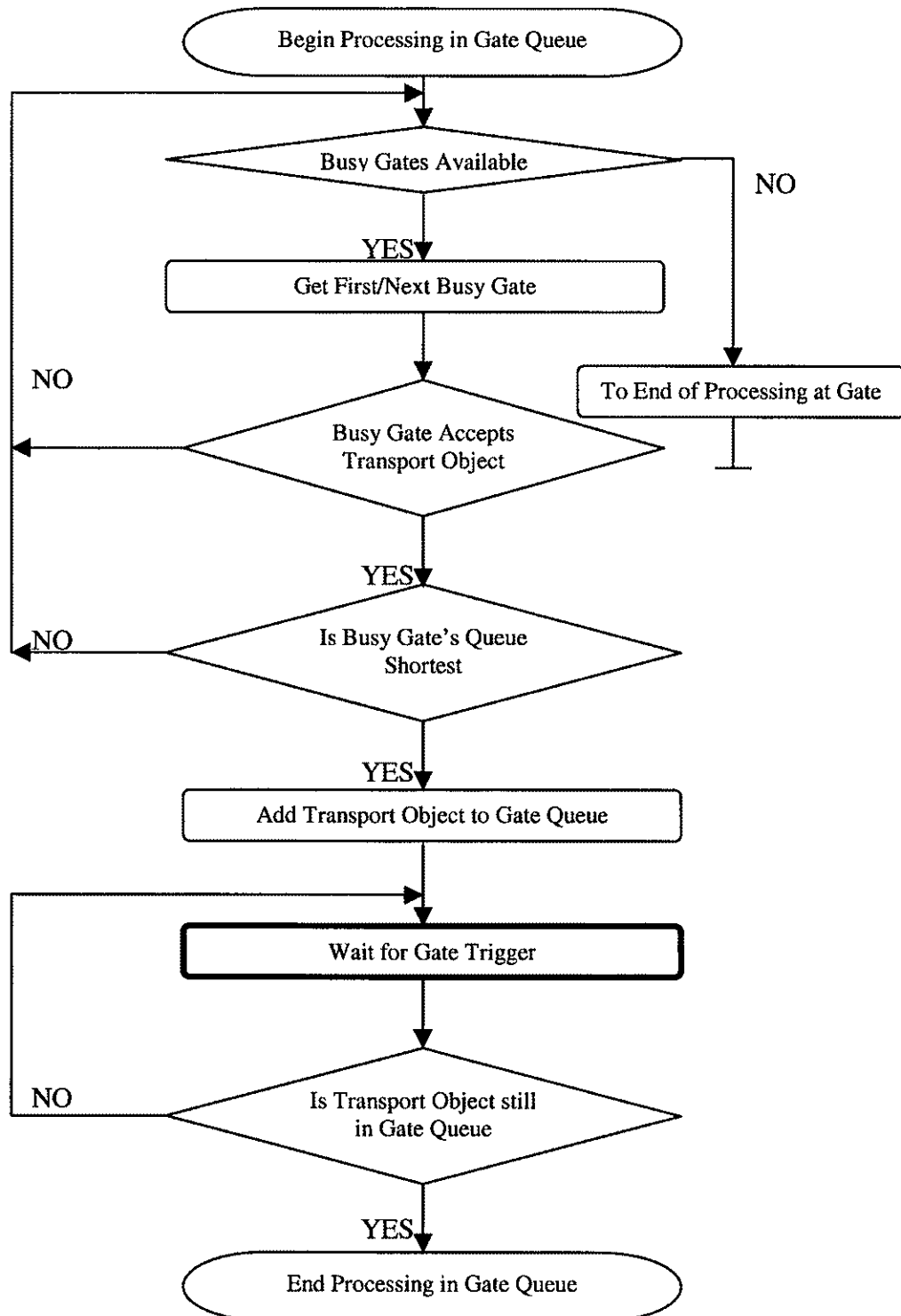


Fig. A.29

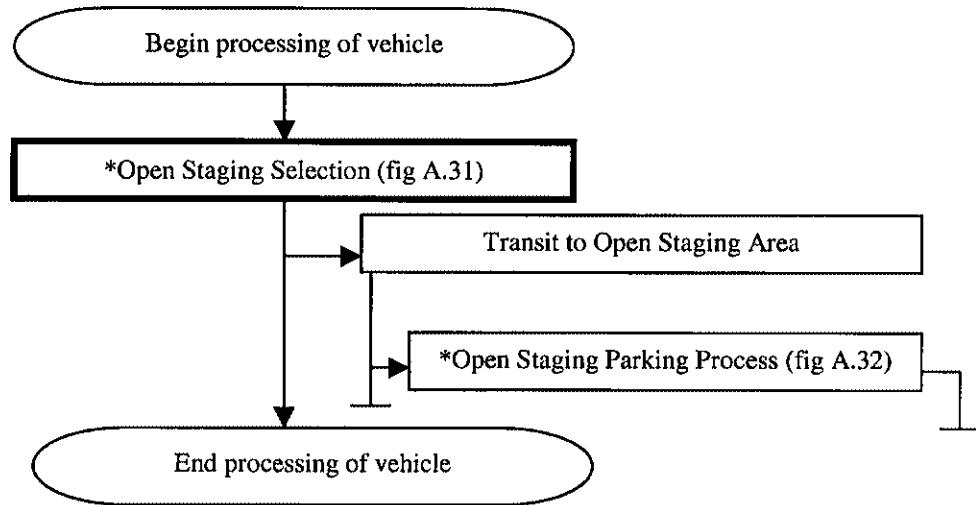


Fig. A.30

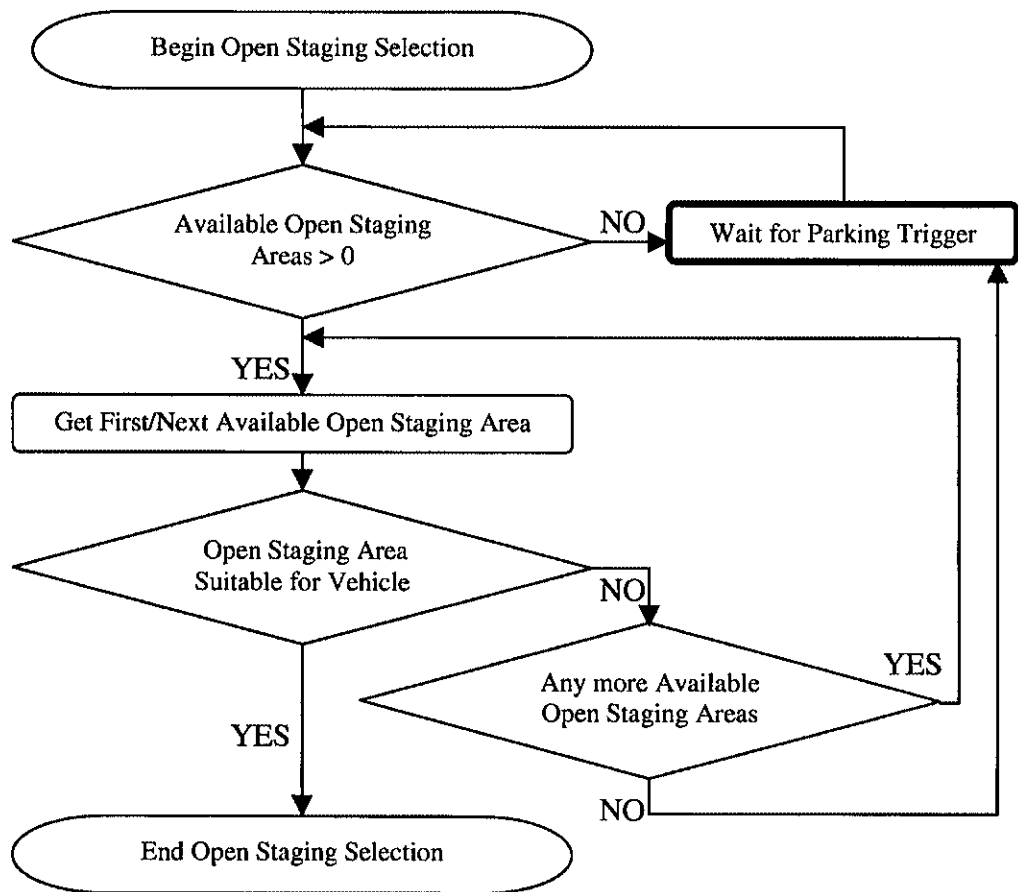


Fig. A.31

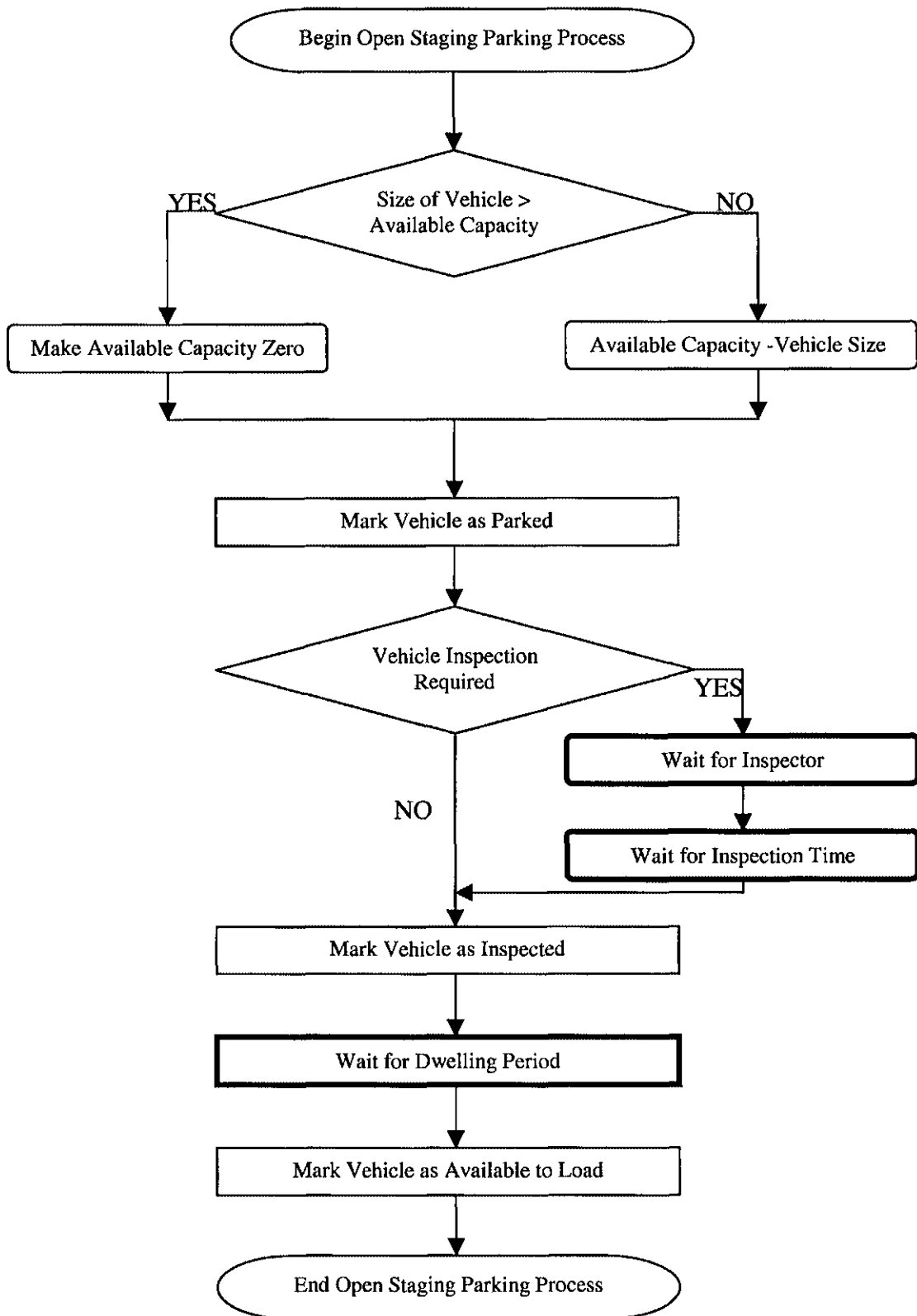


Fig. A.32



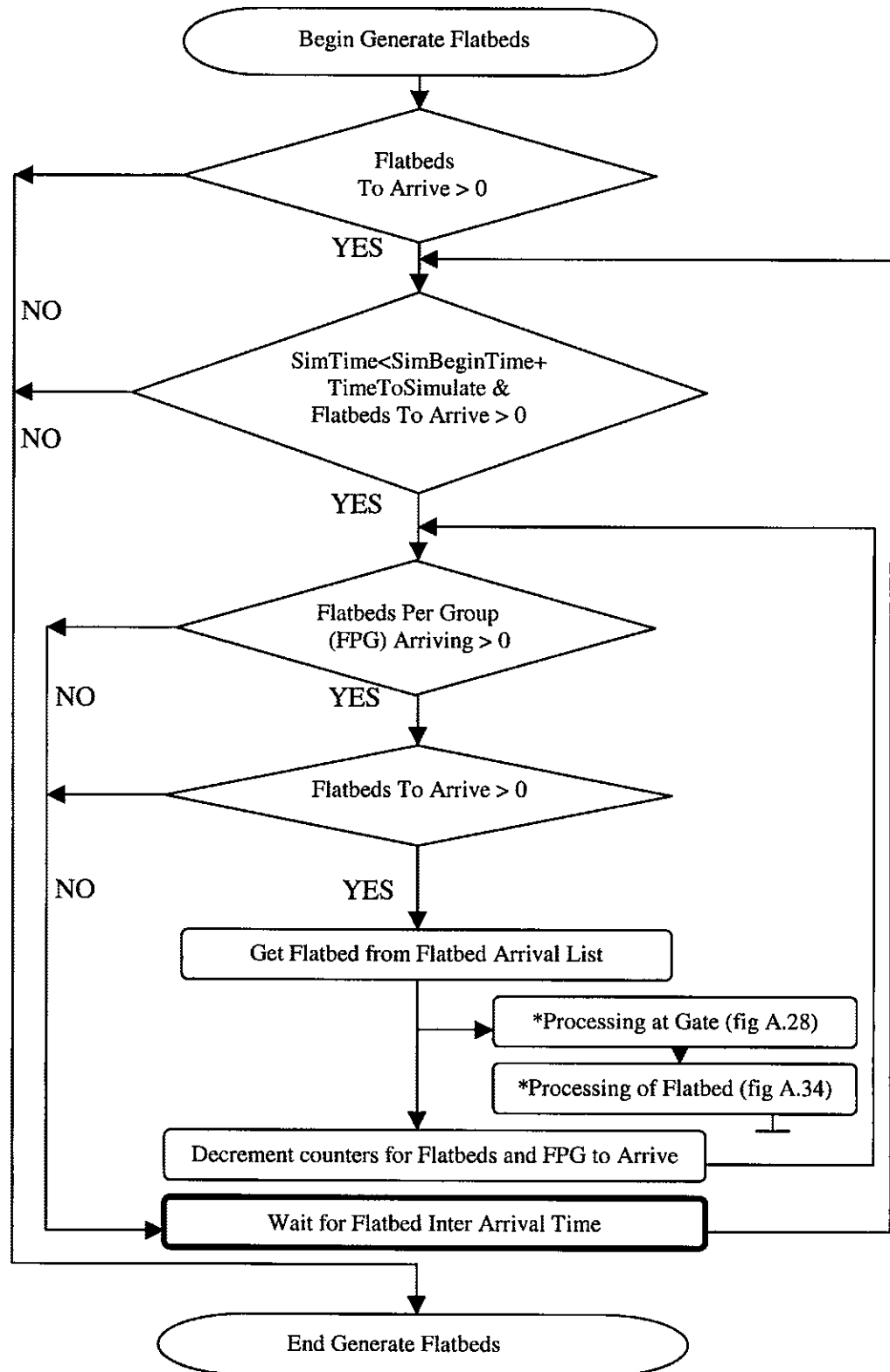


Fig. A.33

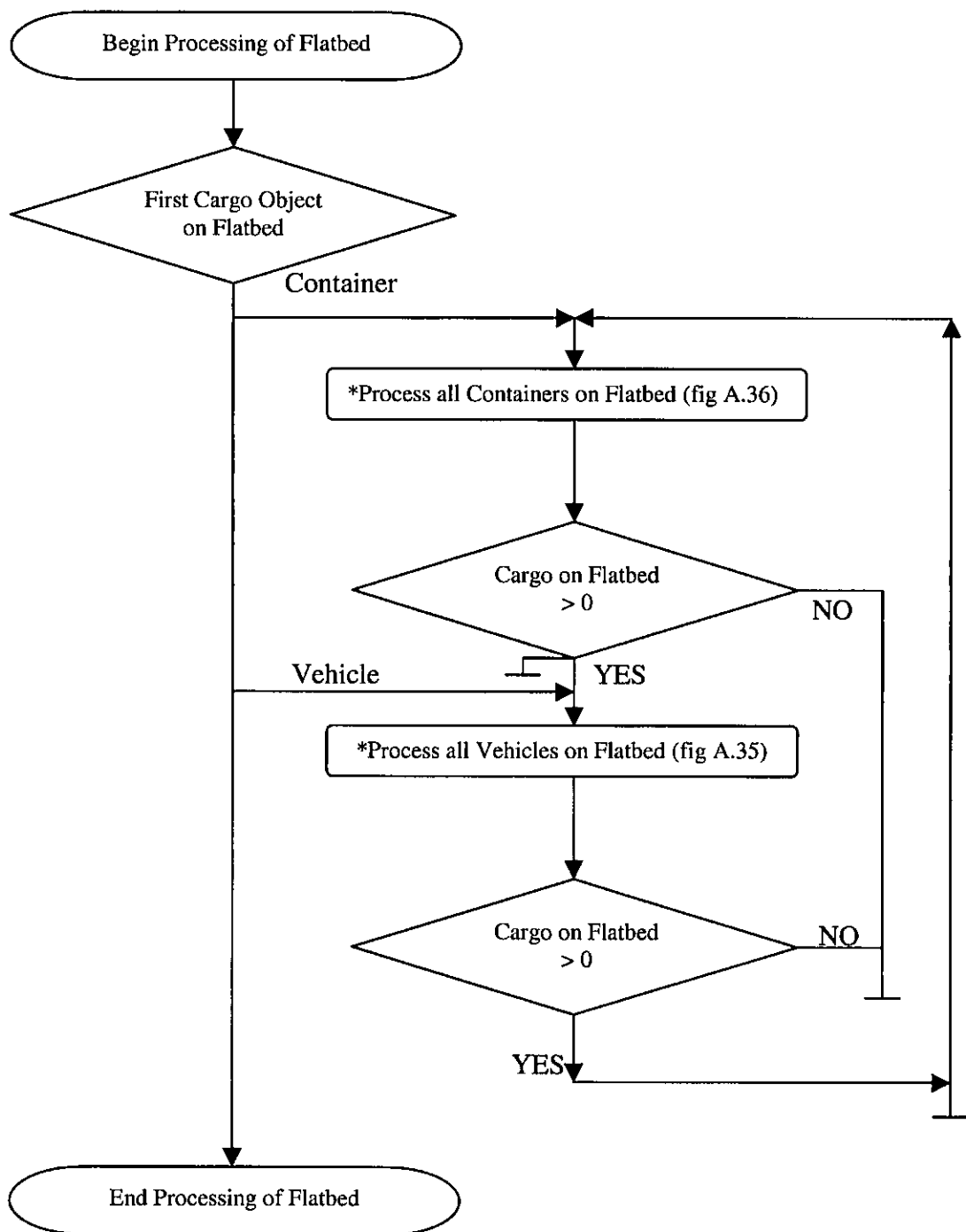


Fig. A.34

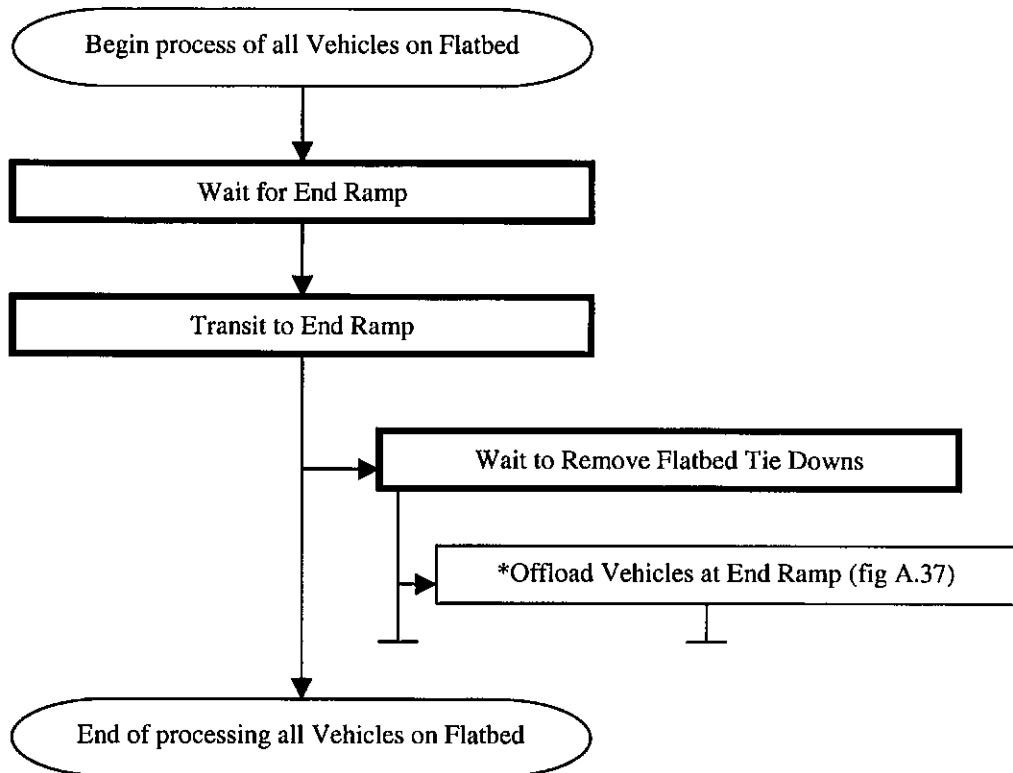


Fig. A.35

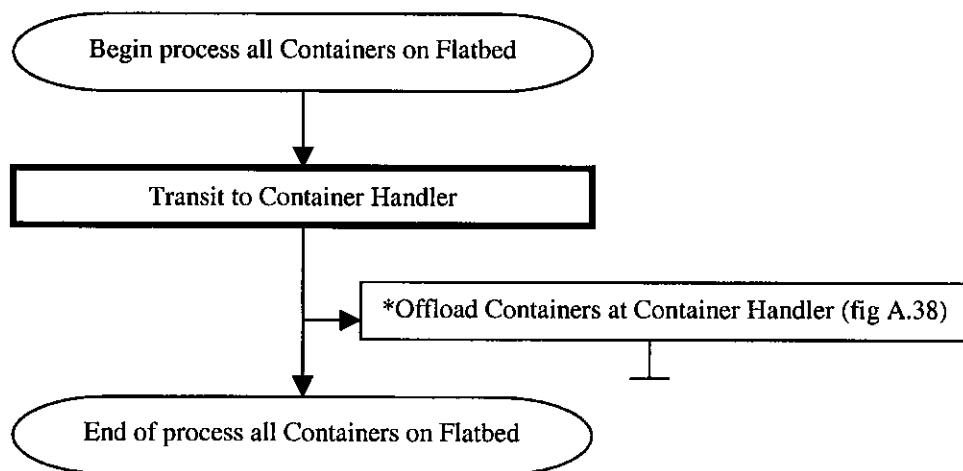


Fig. A.36

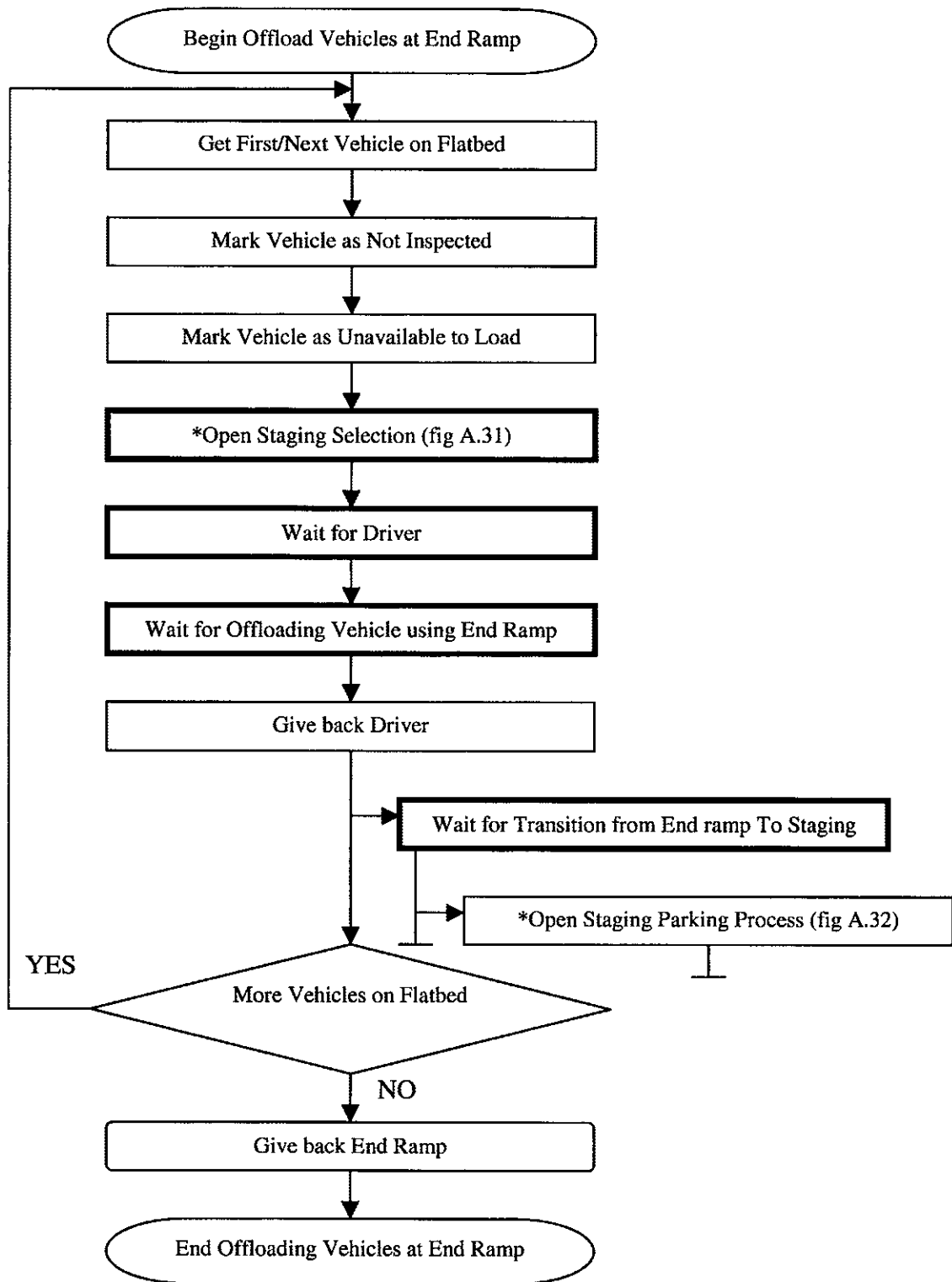


Fig. A.37

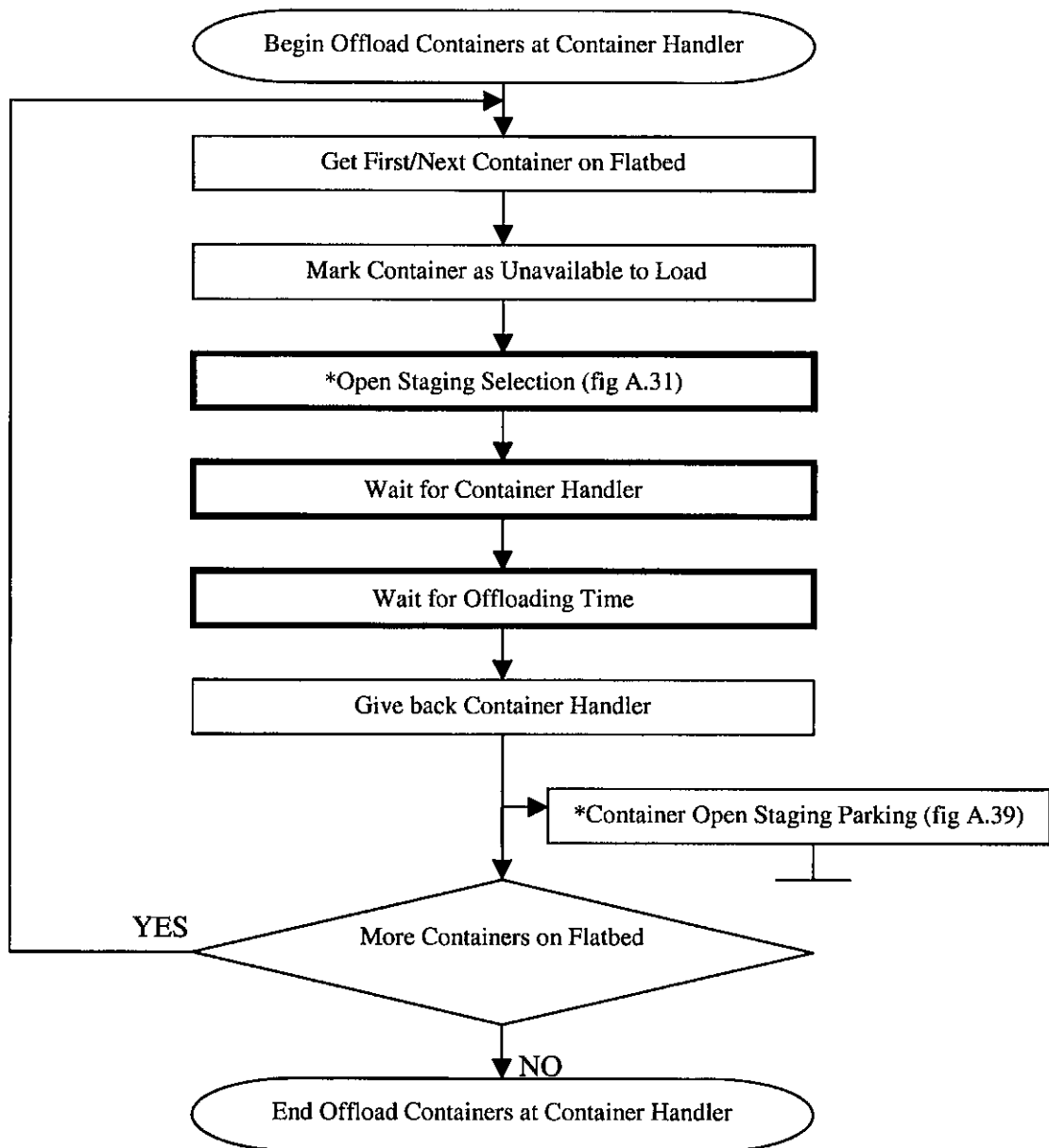


Fig. A.38

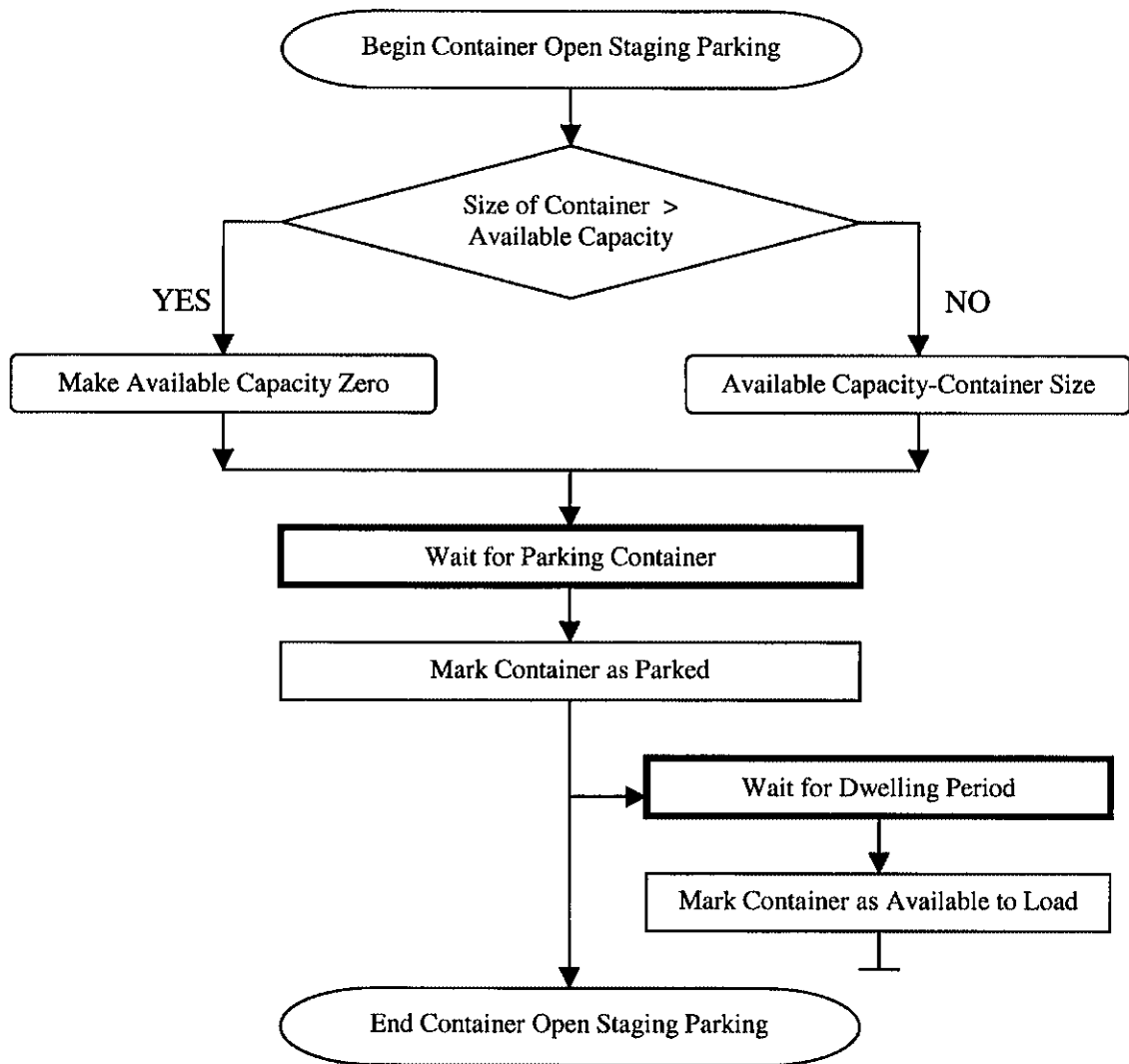


Fig. A.39

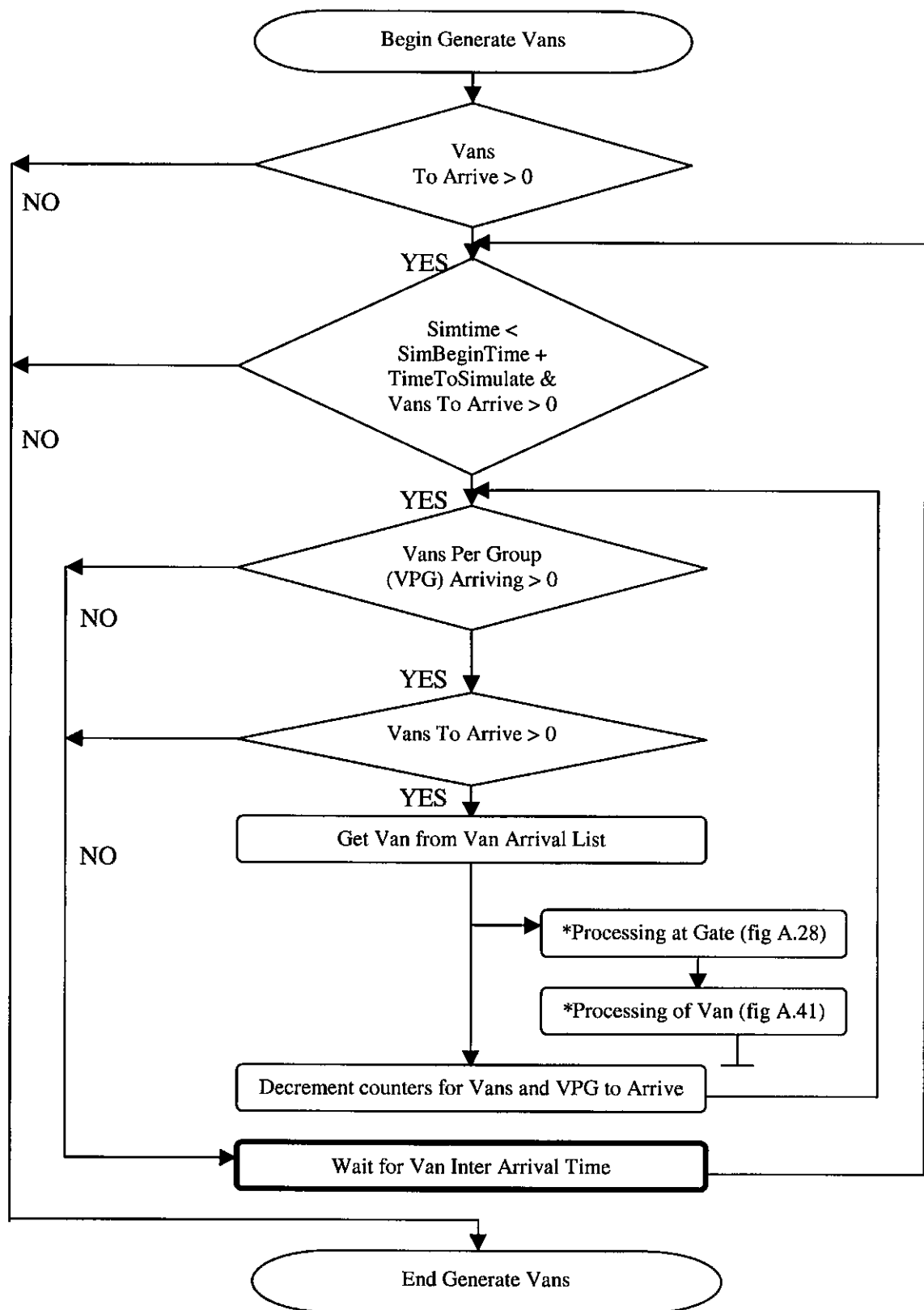


Fig. A.40

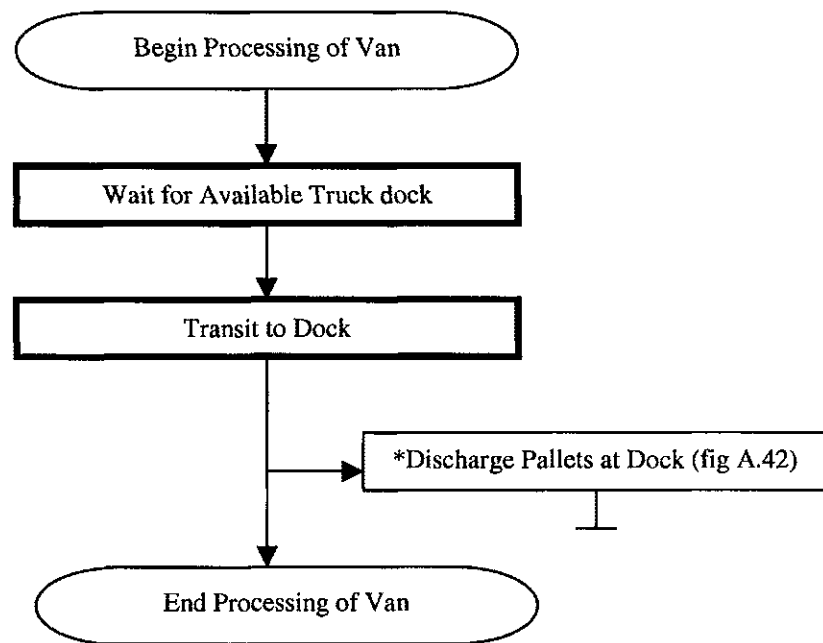


Fig. A.41

In Fig. 42, the program flow for "Covered staging area selection" process is similar to that of "Open staging selection," and the program flow for "Covered staging area parking" process is similar to that of "Container Open staging parking".



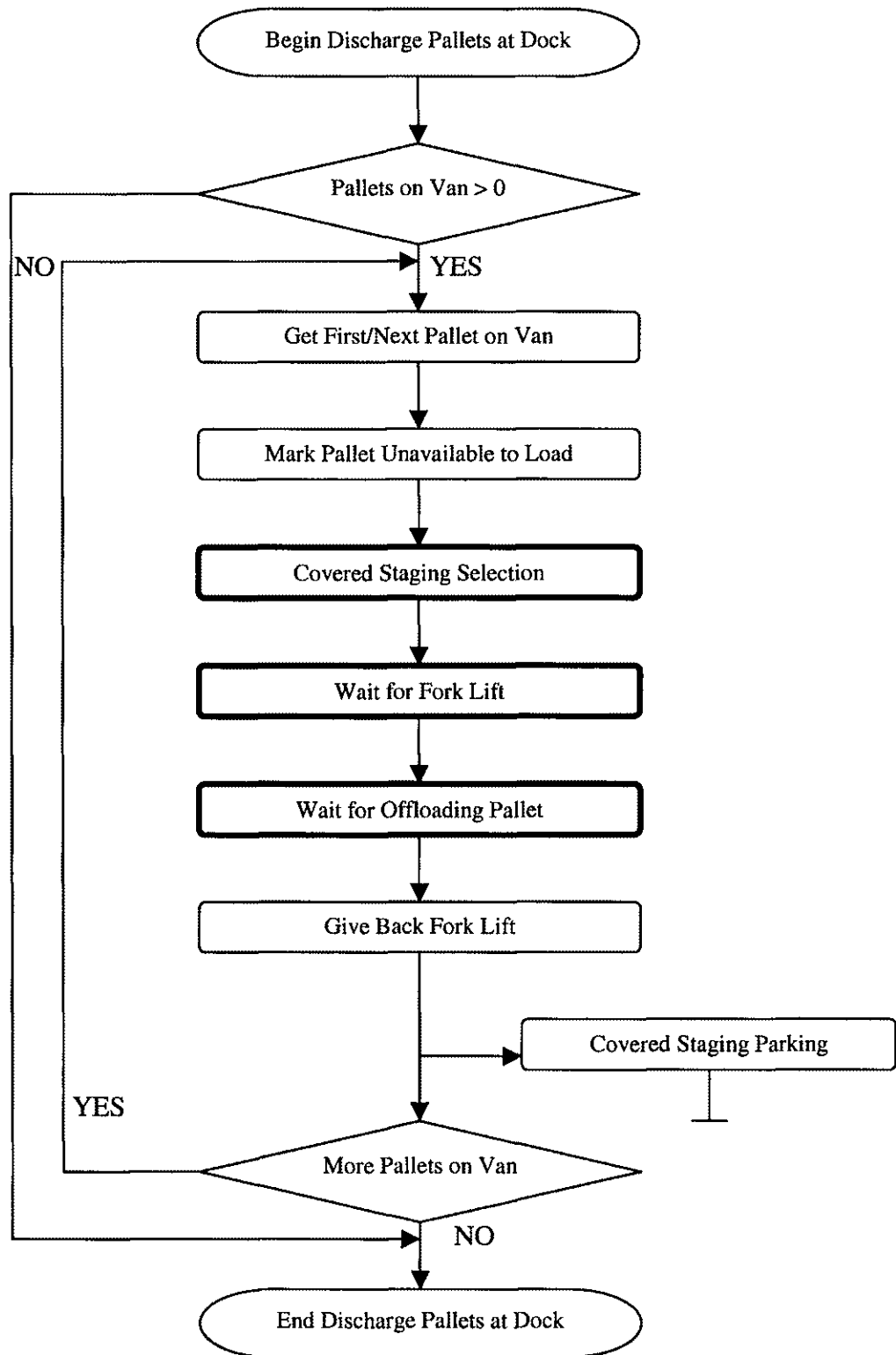


Fig. A.42

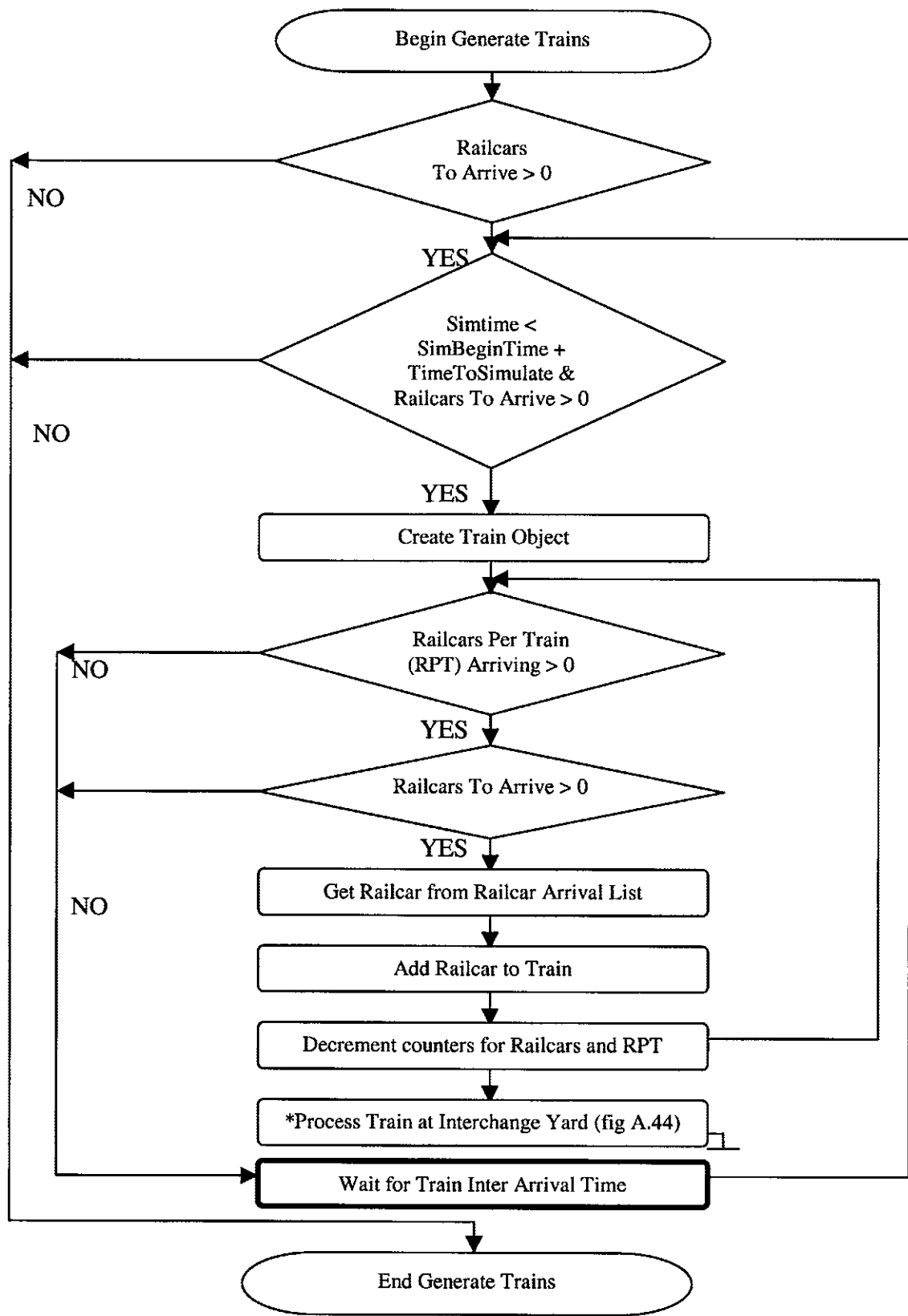


Fig. A.43

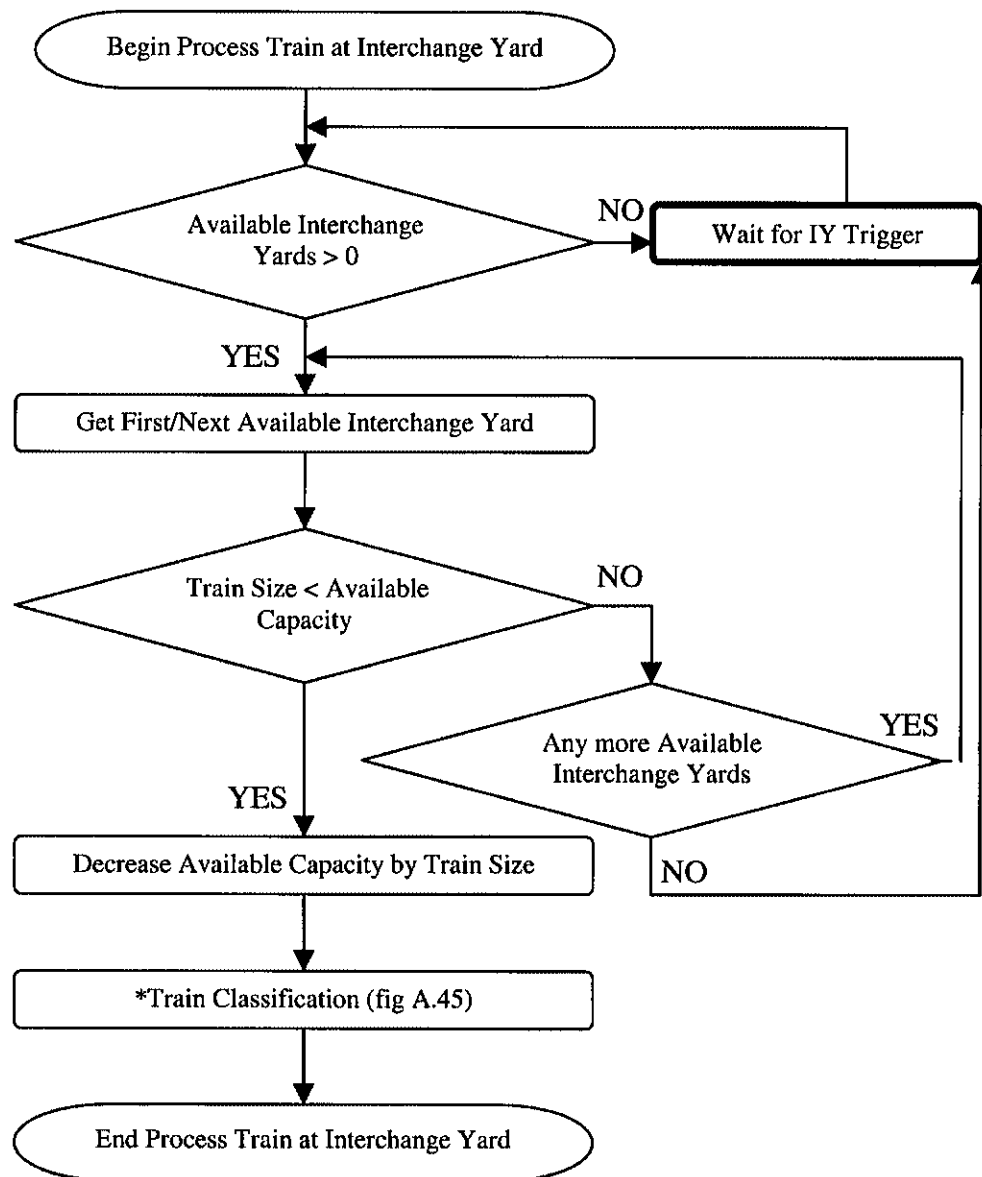


Fig. A.44

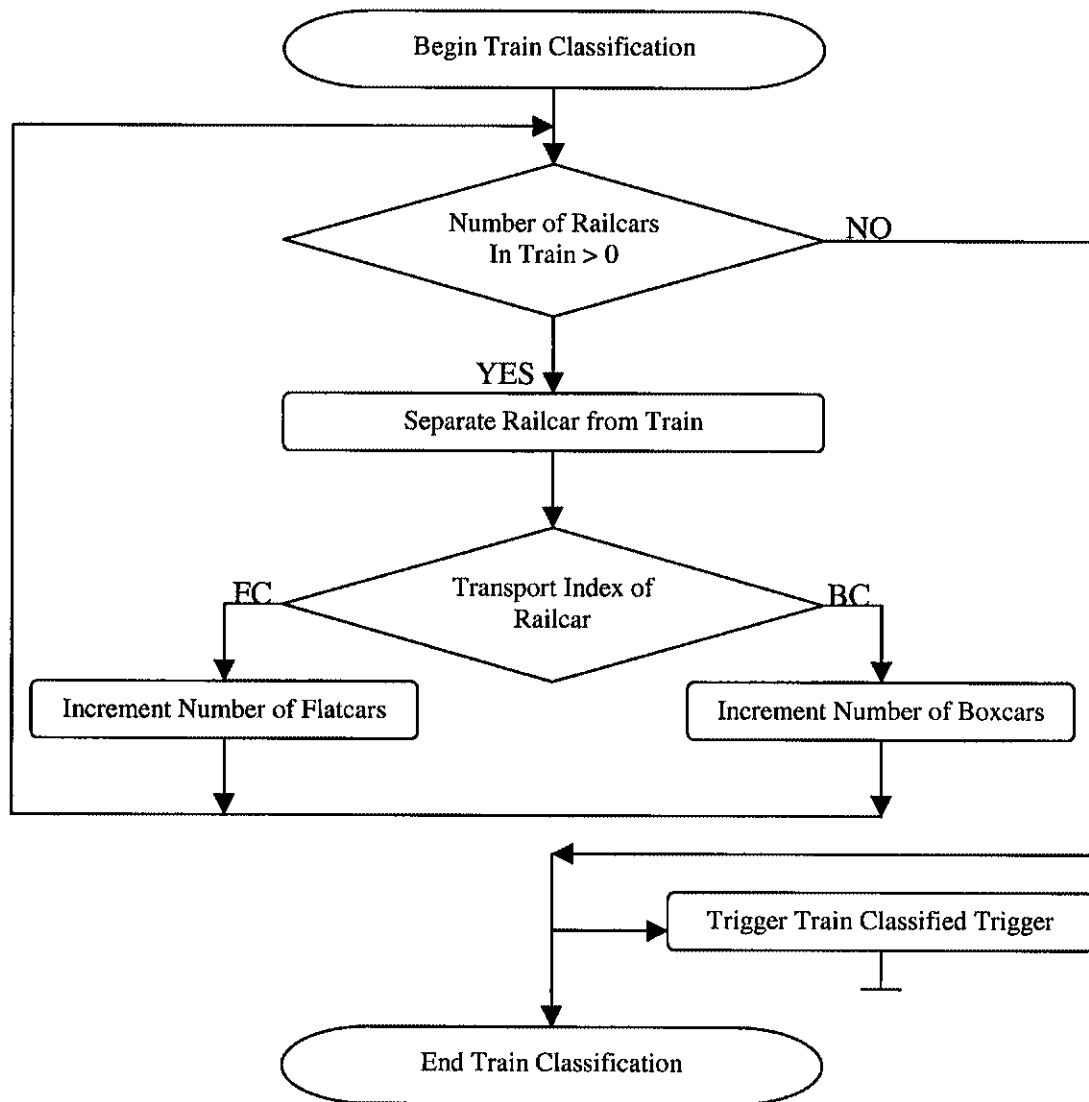


Fig. A.45

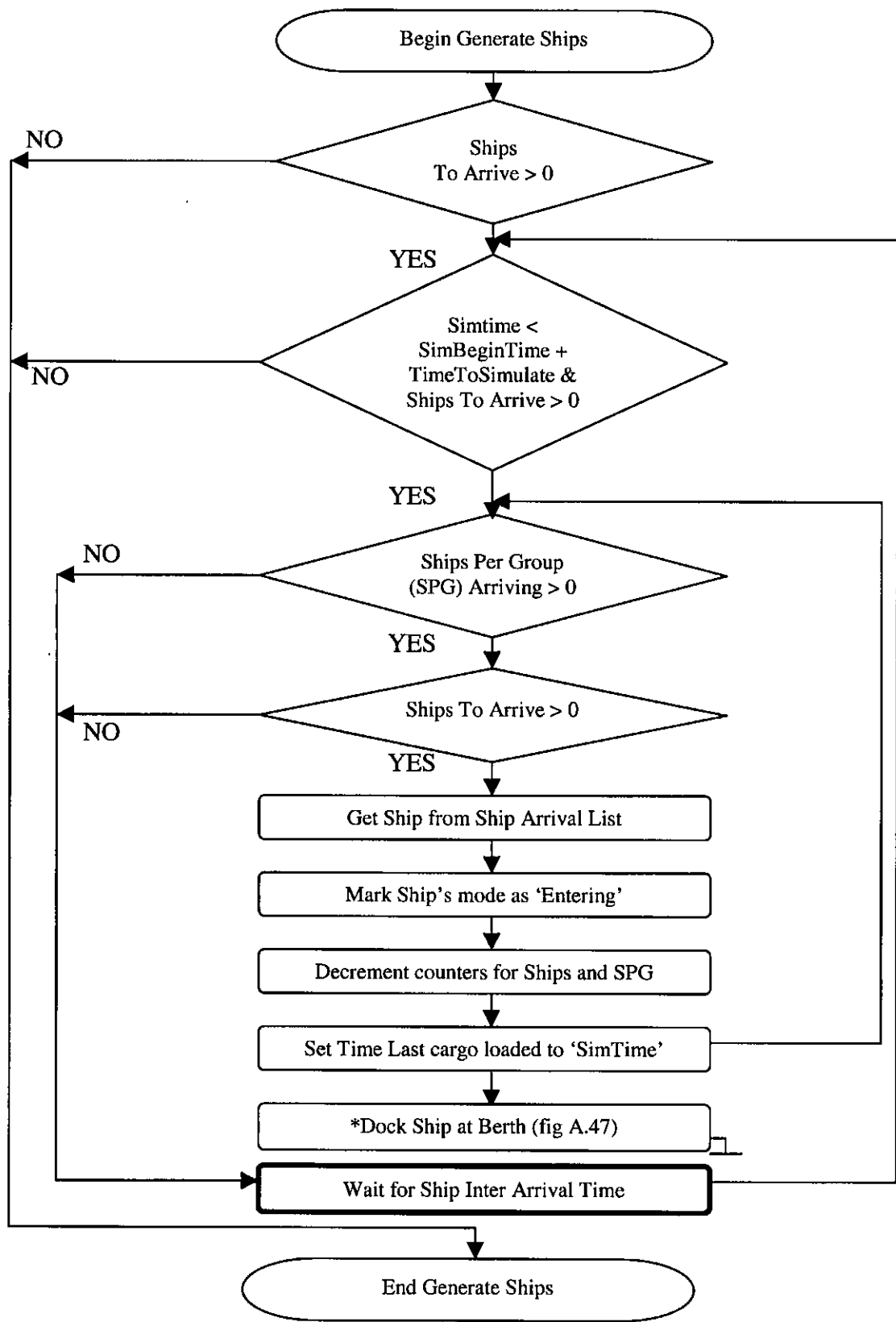


Fig. A.46

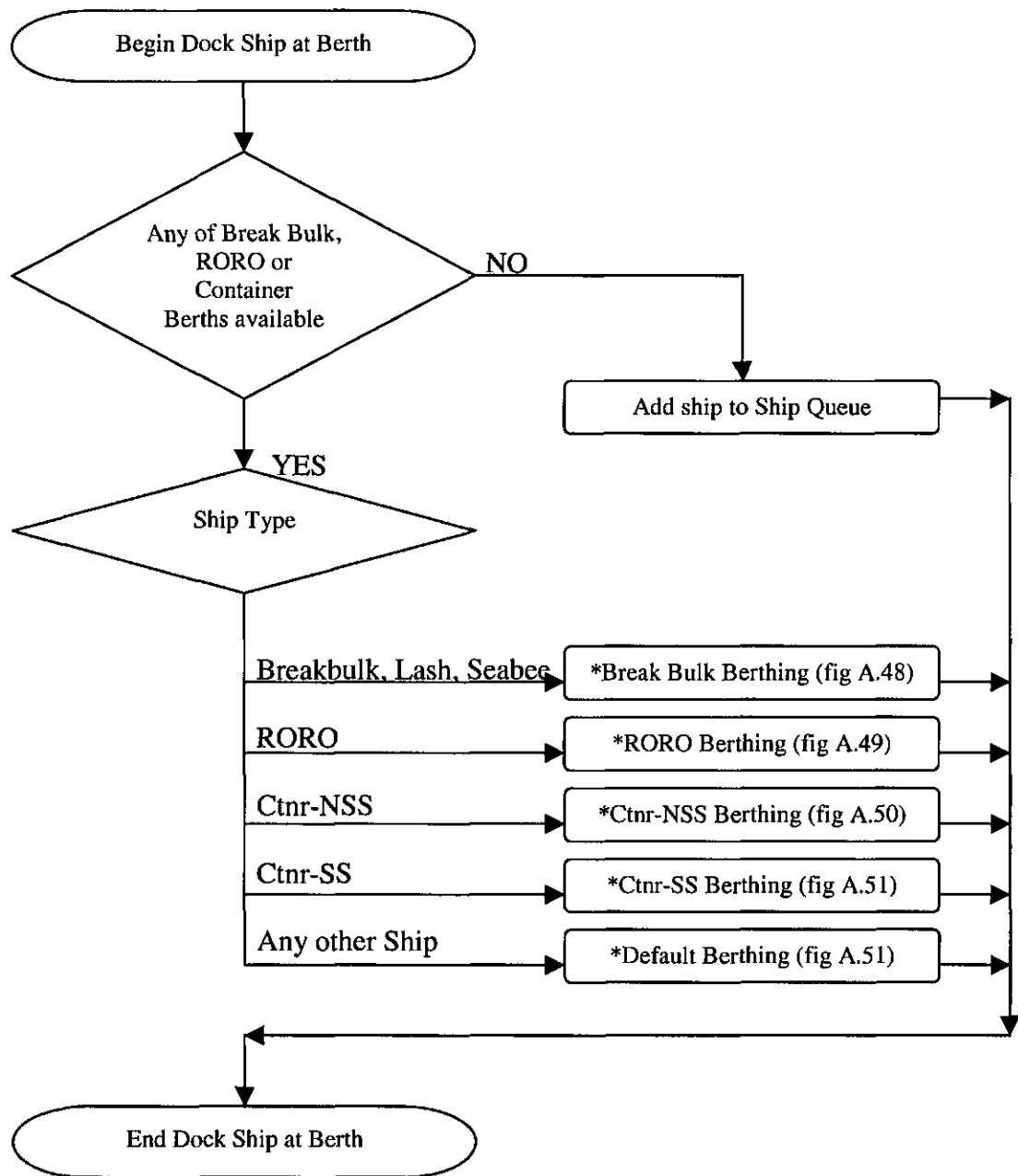


Fig. A.47

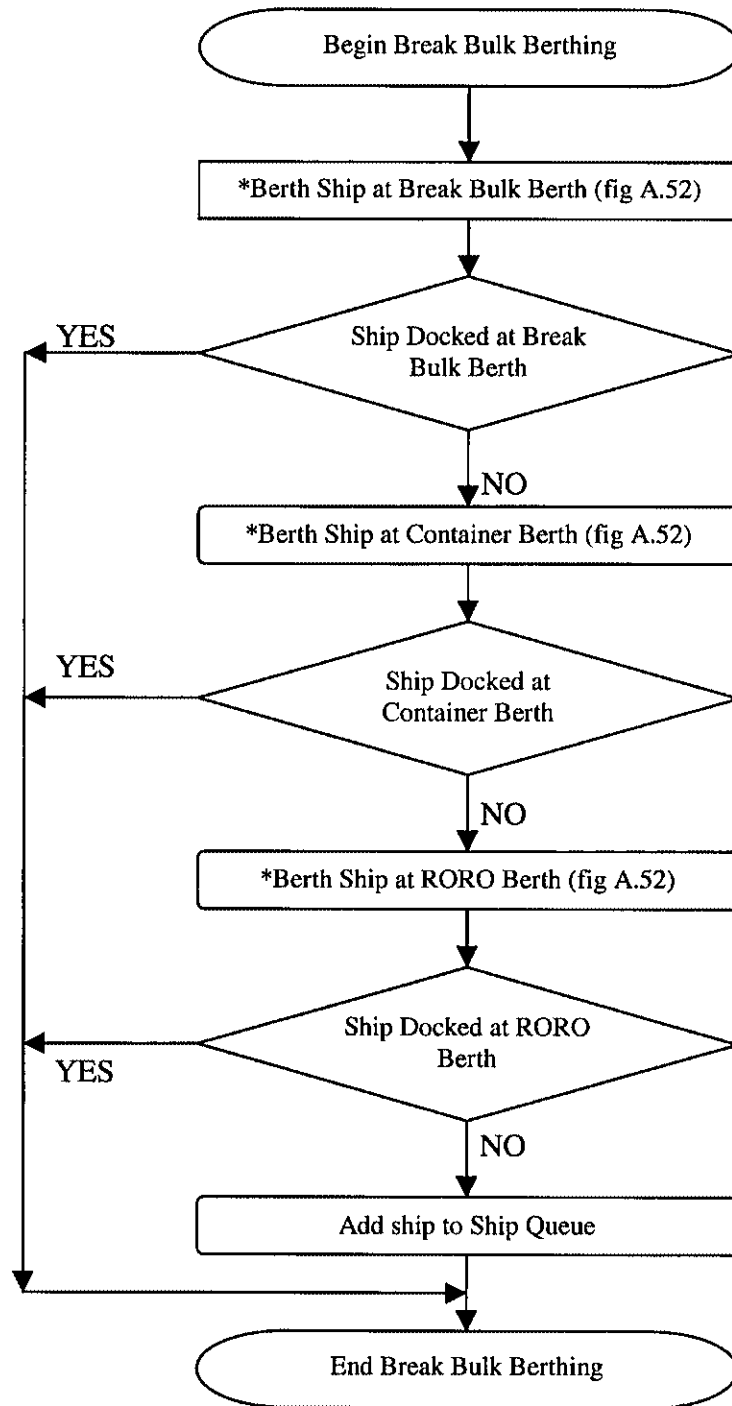


Fig. A.48

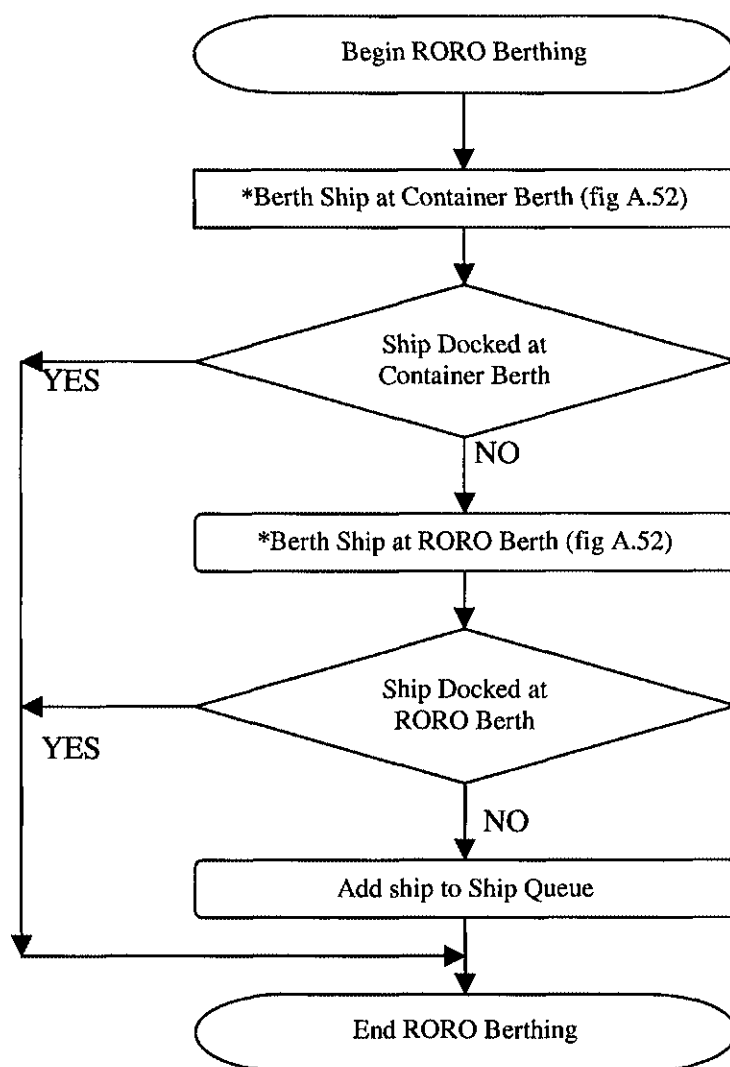


Fig. A.49



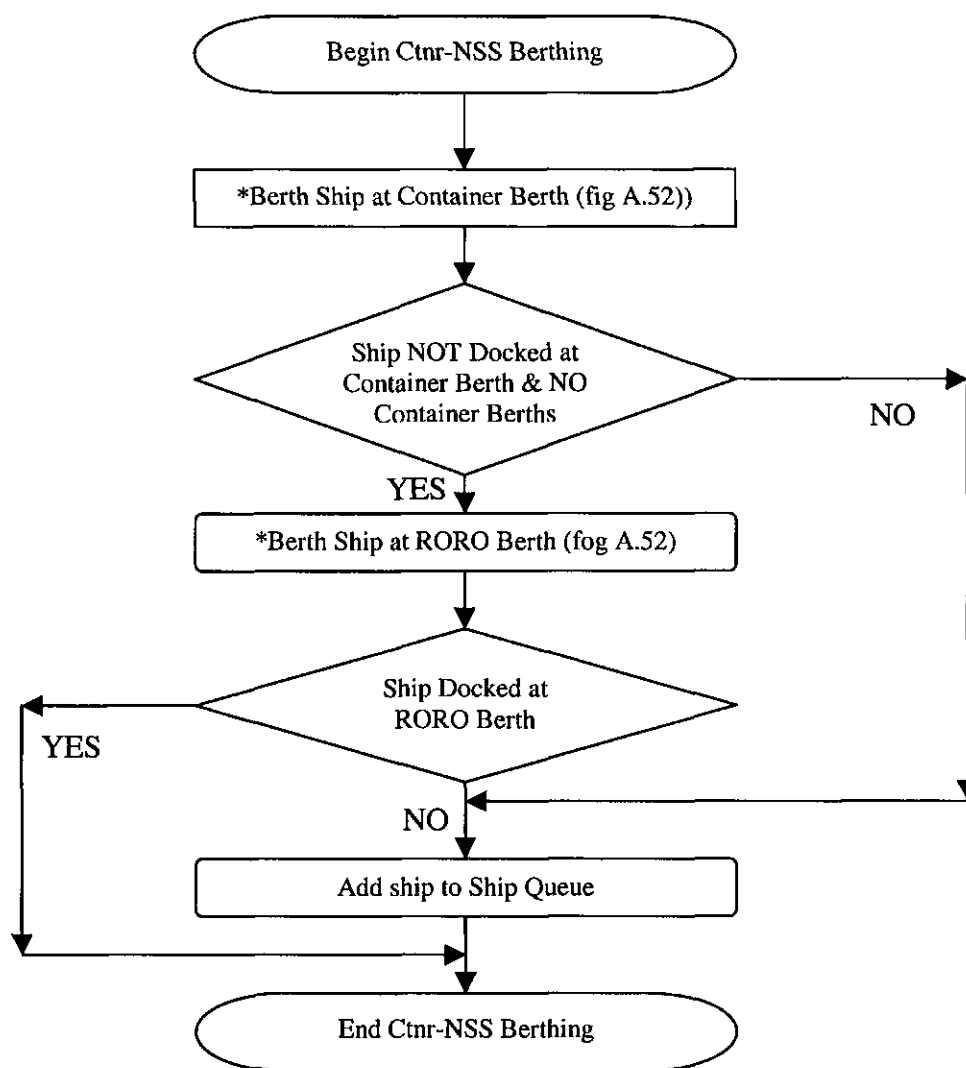


Fig. A.50

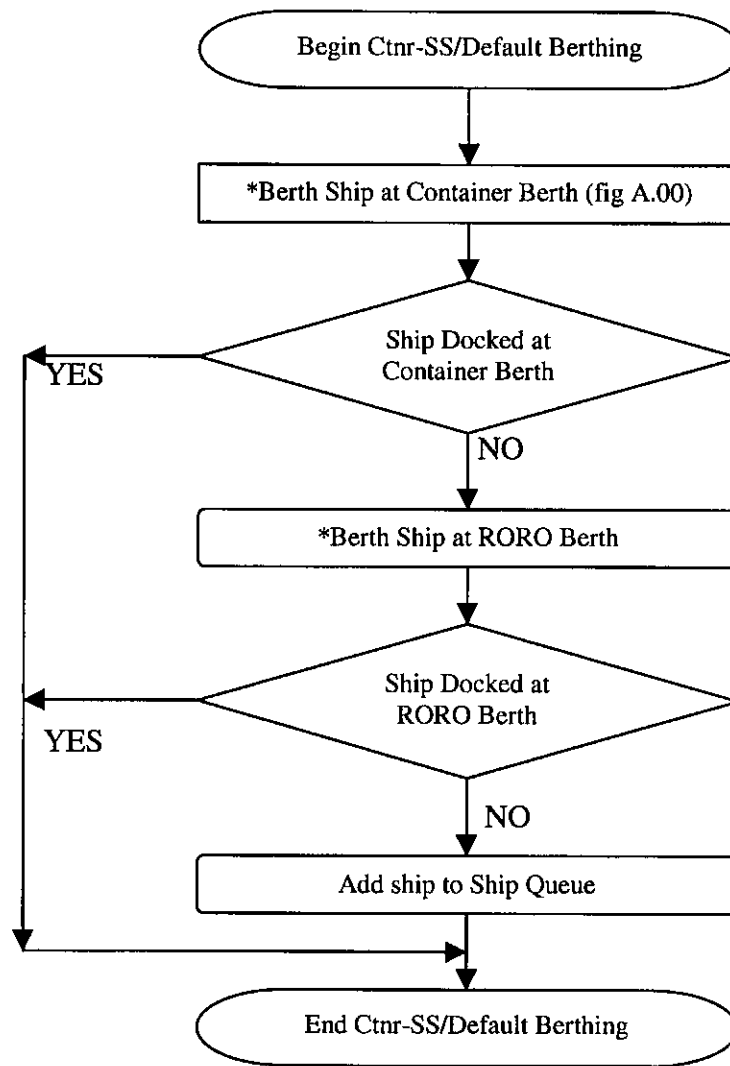


Fig. A.51

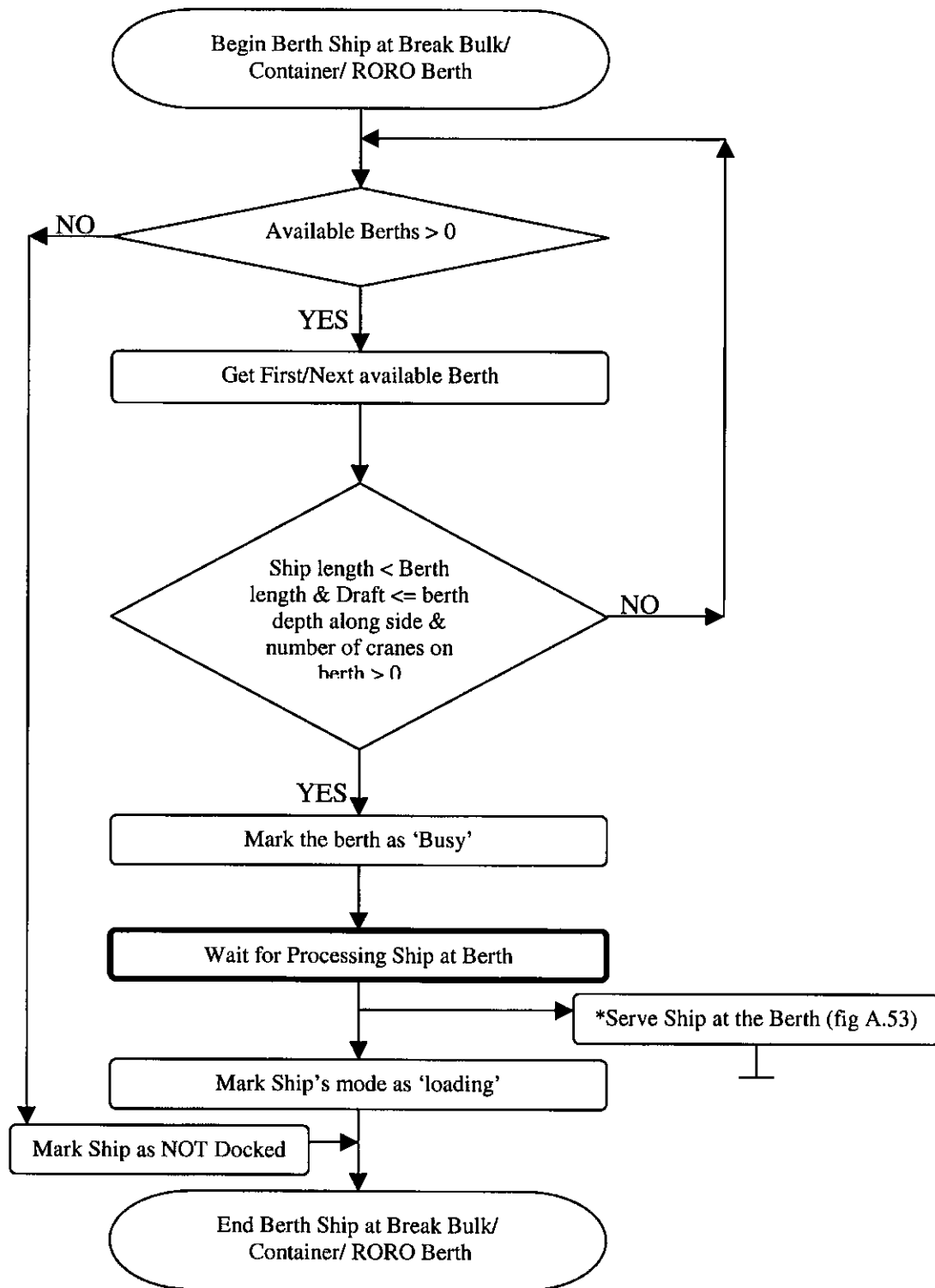


Fig. A.52

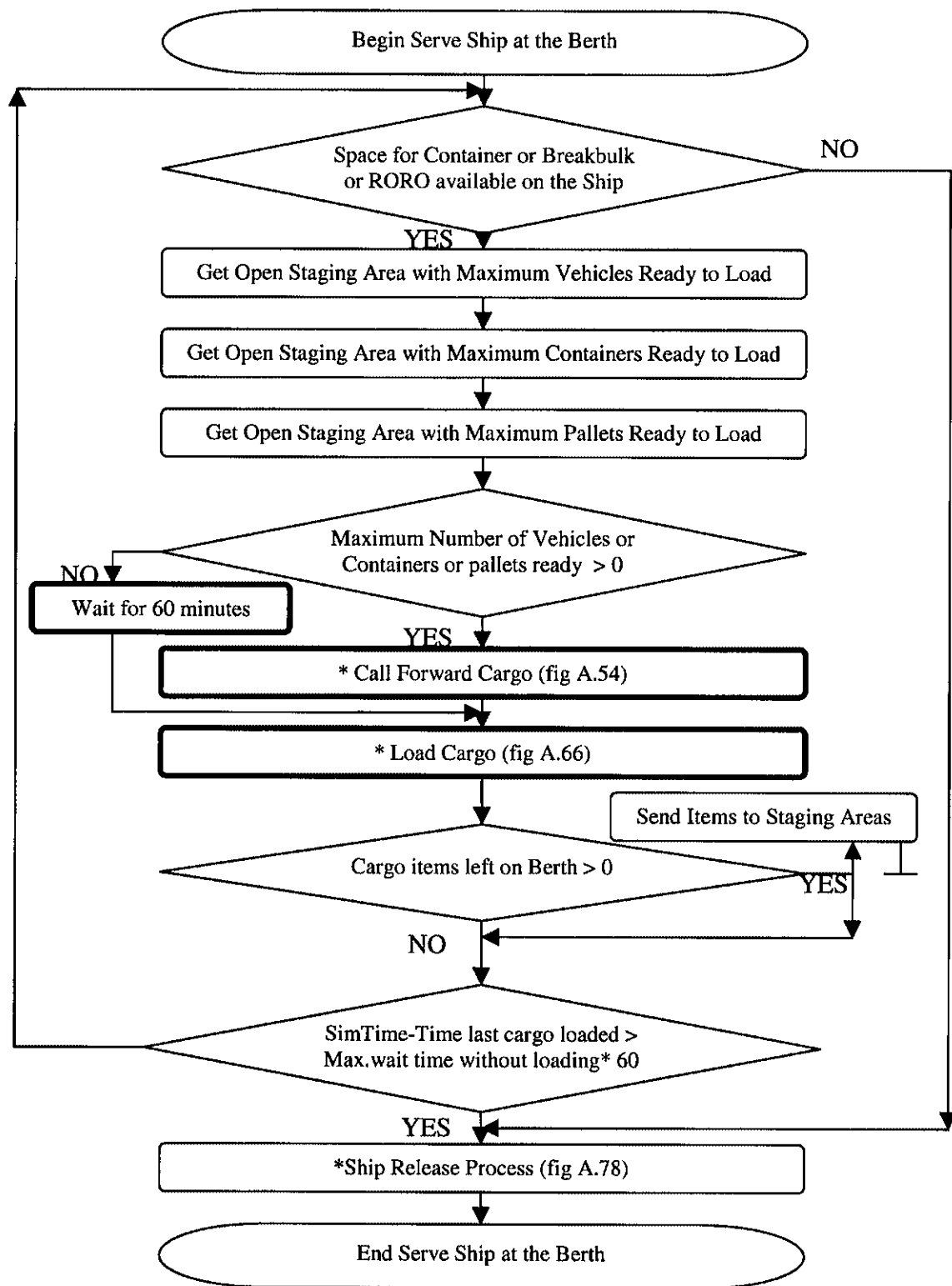


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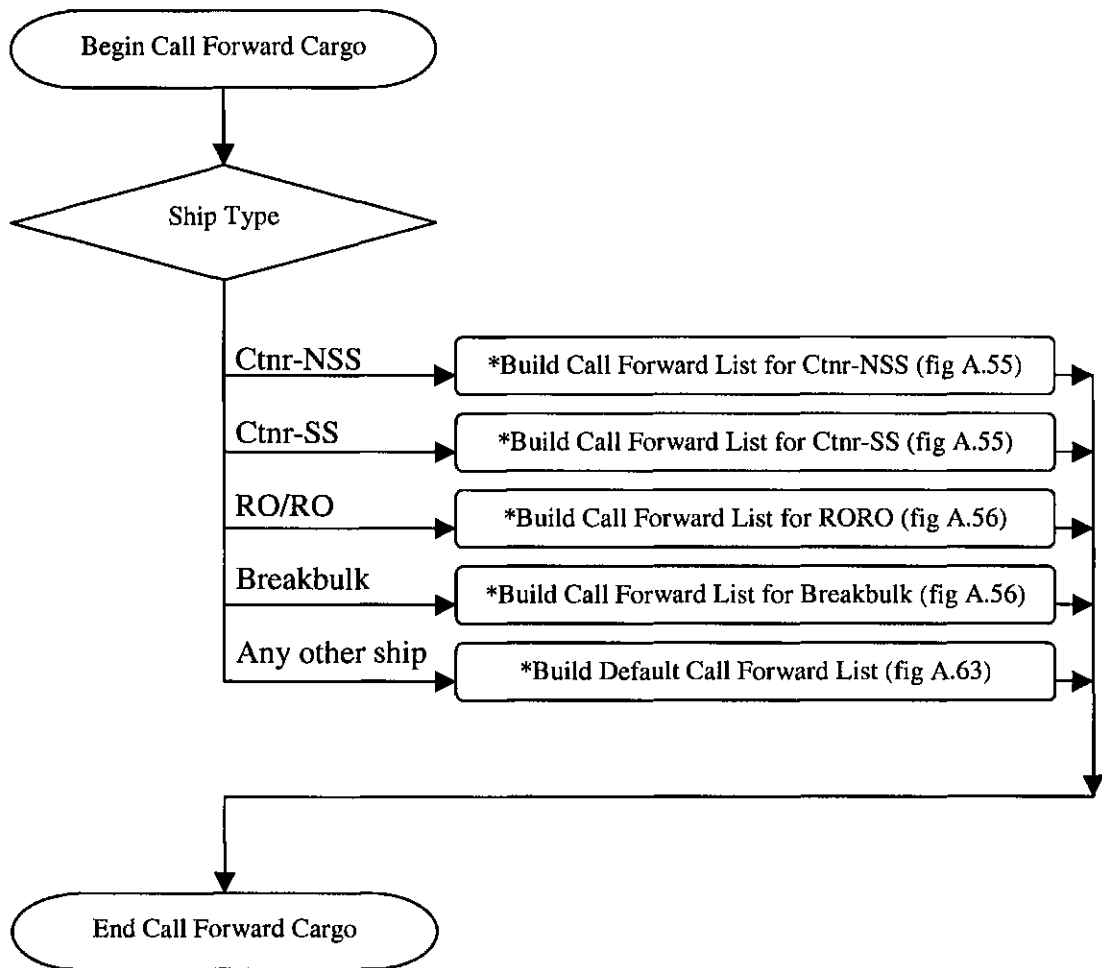


Fig. A.54

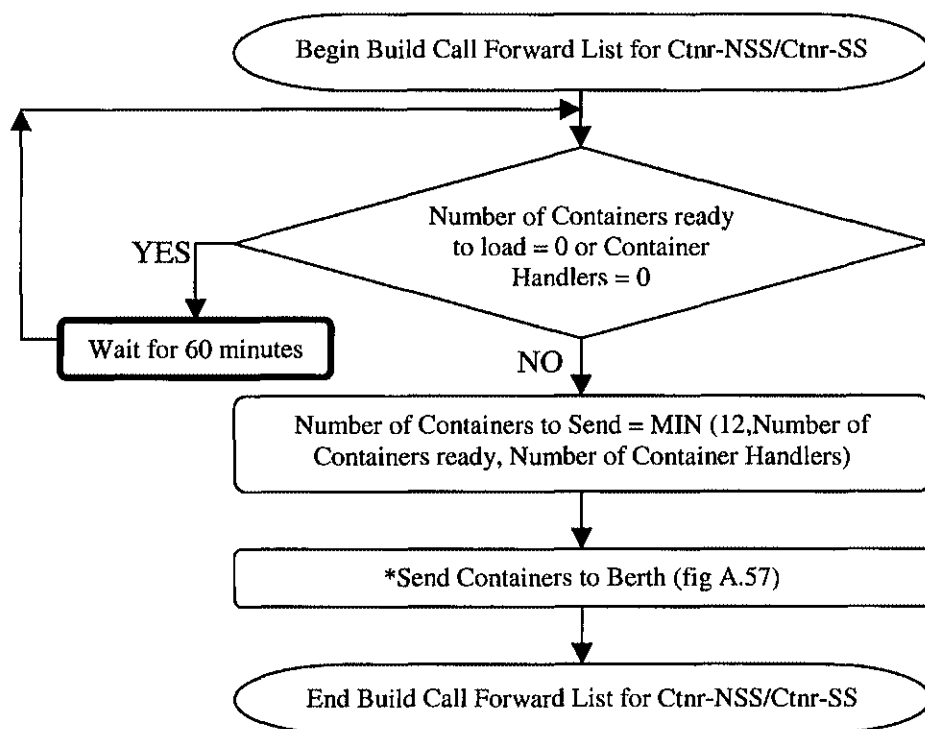


Fig. A.55

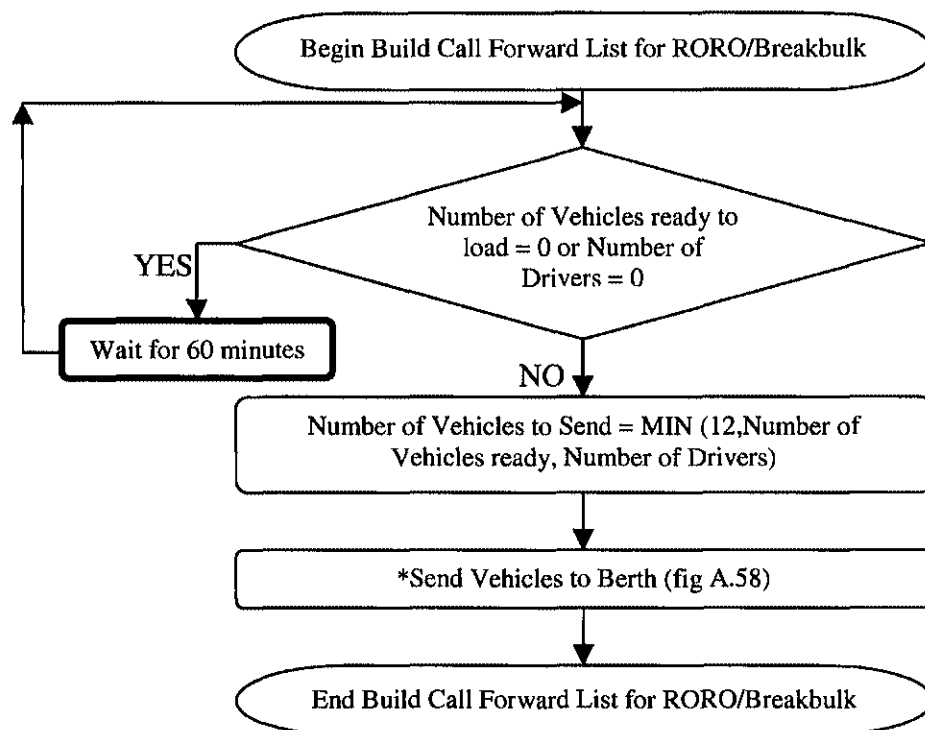


Fig. A.56

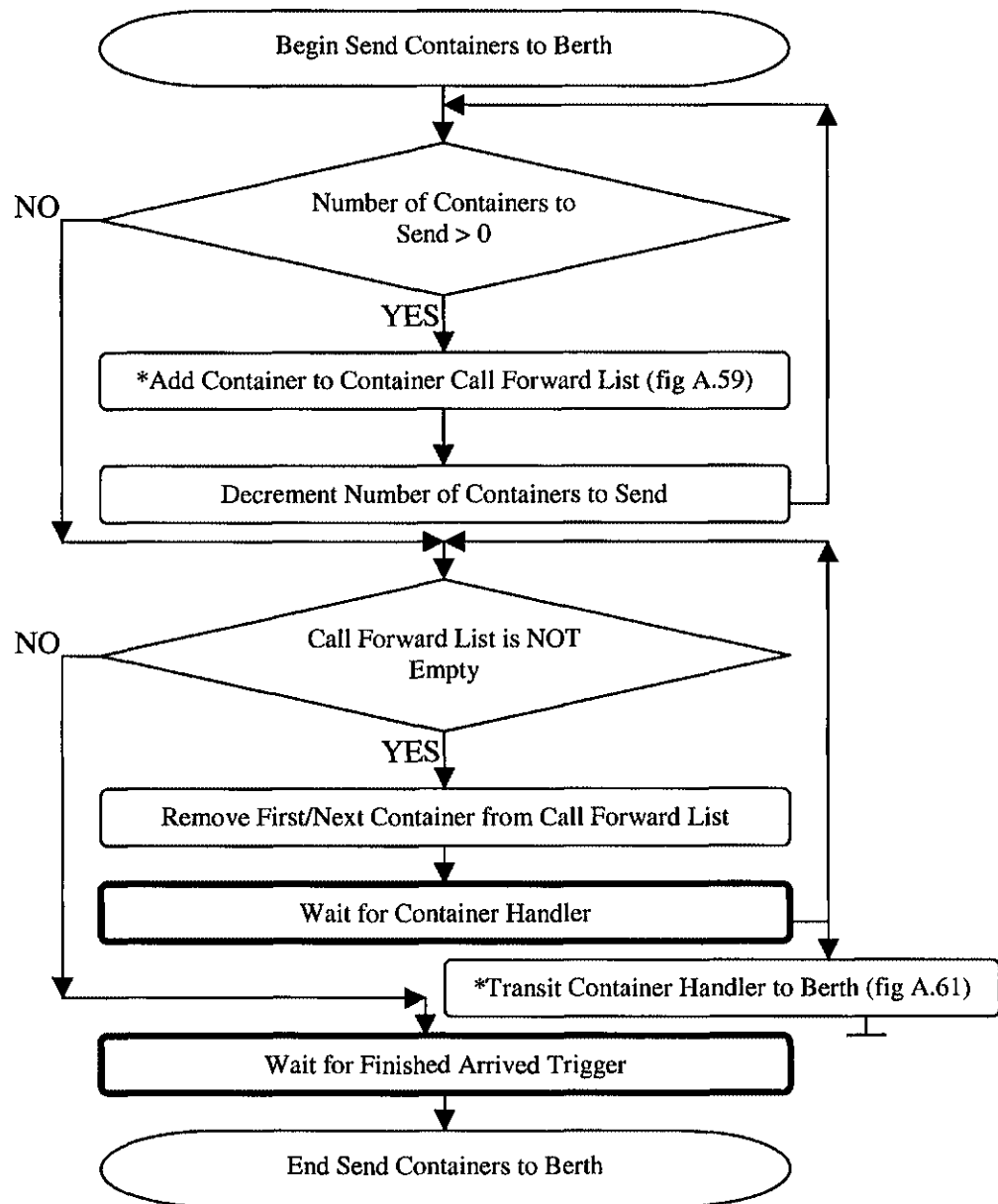


Fig. A.57

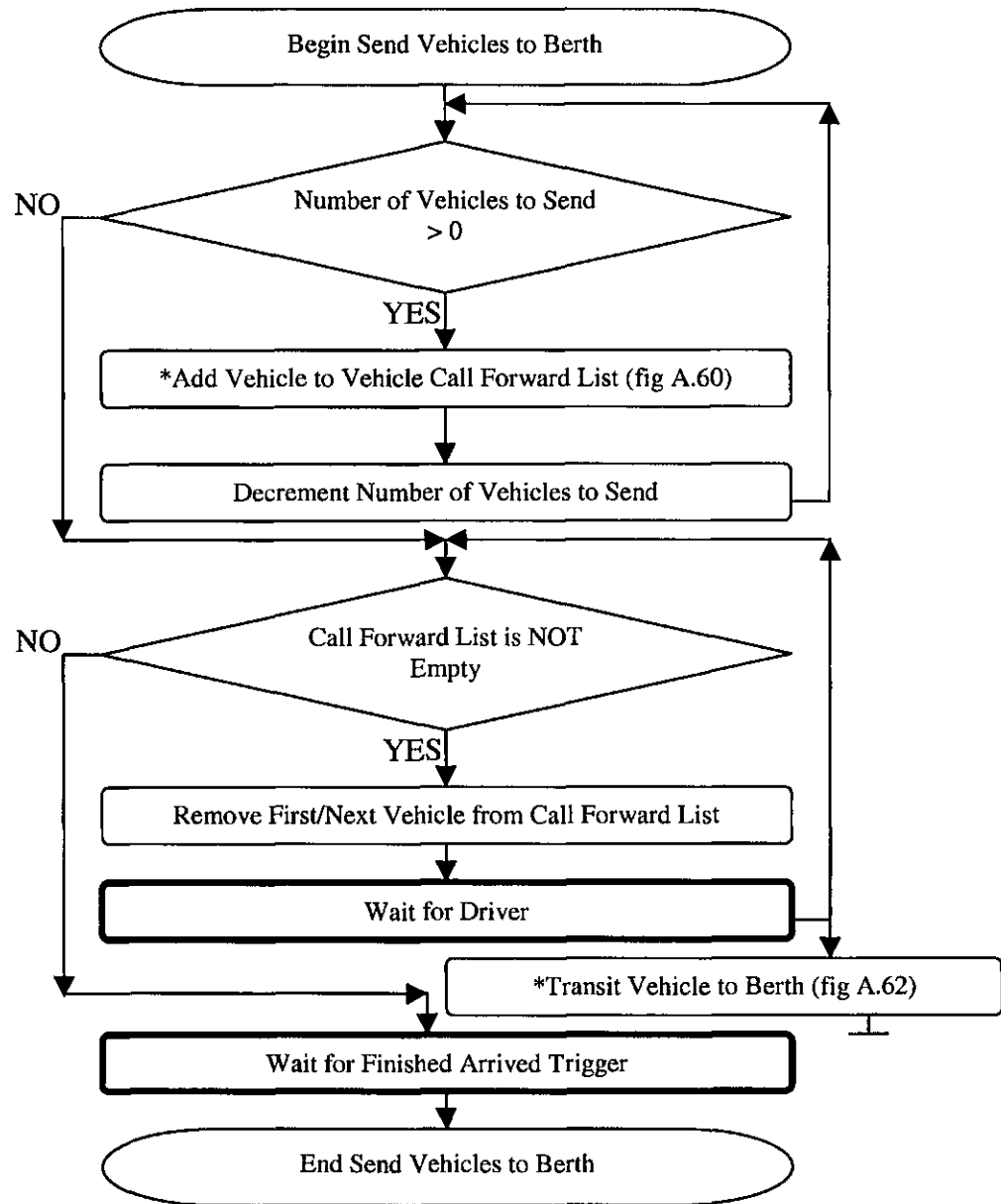


Fig. A.58



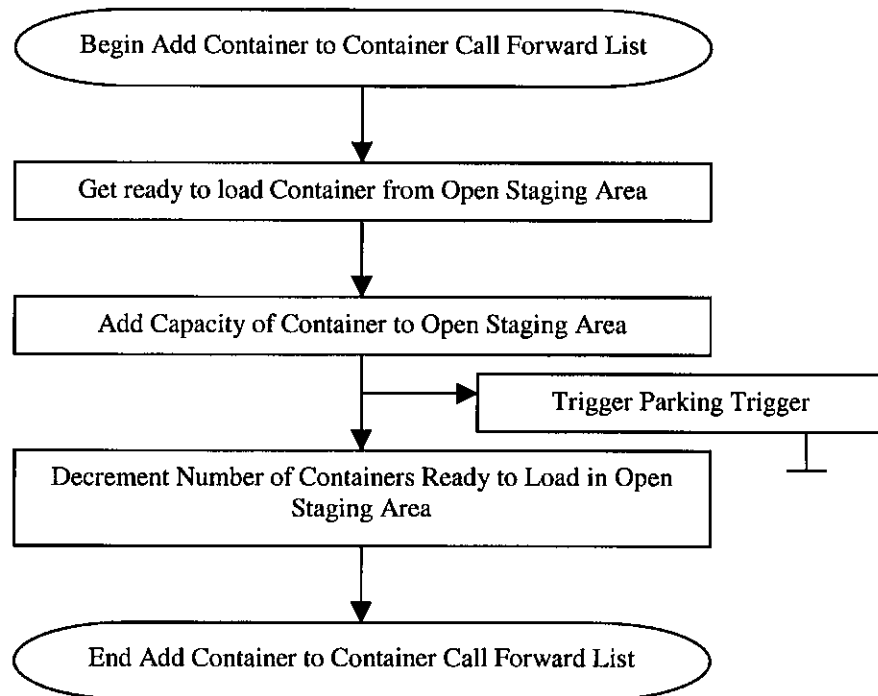


Fig. A.59

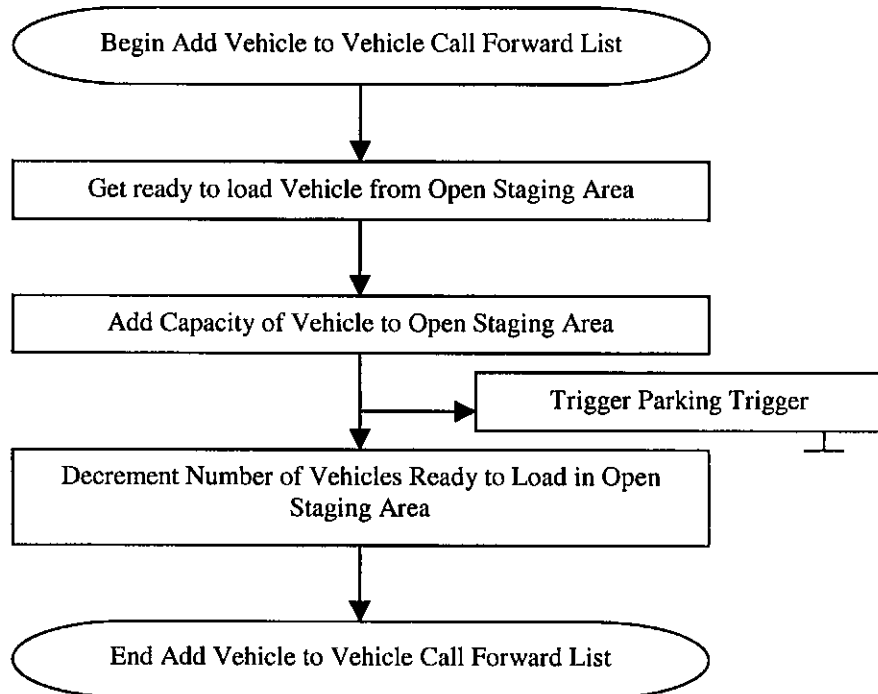


Fig. A.60

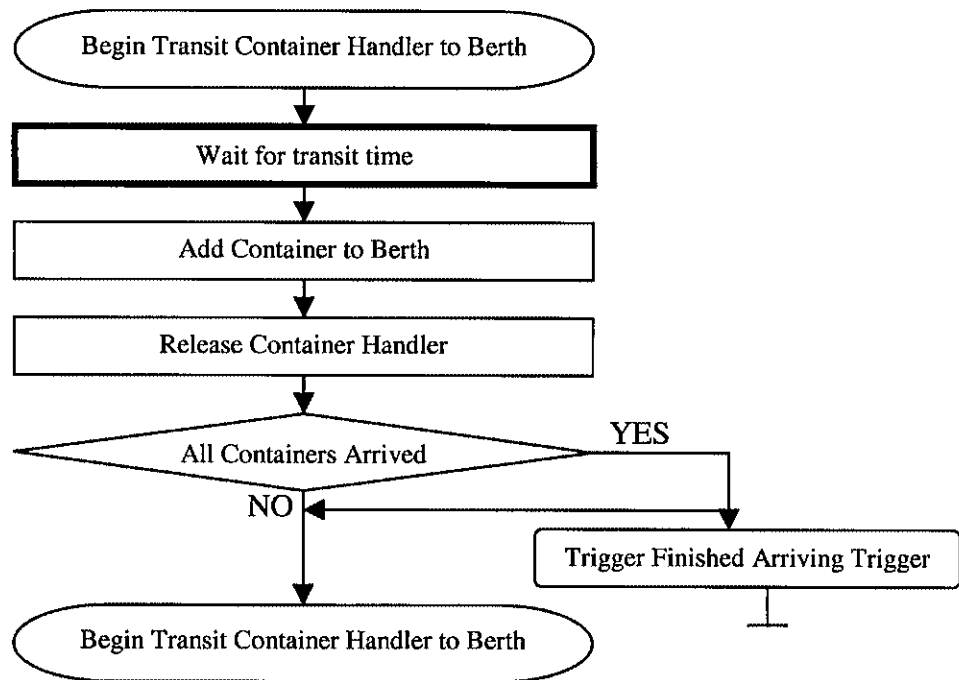


Fig. A.61

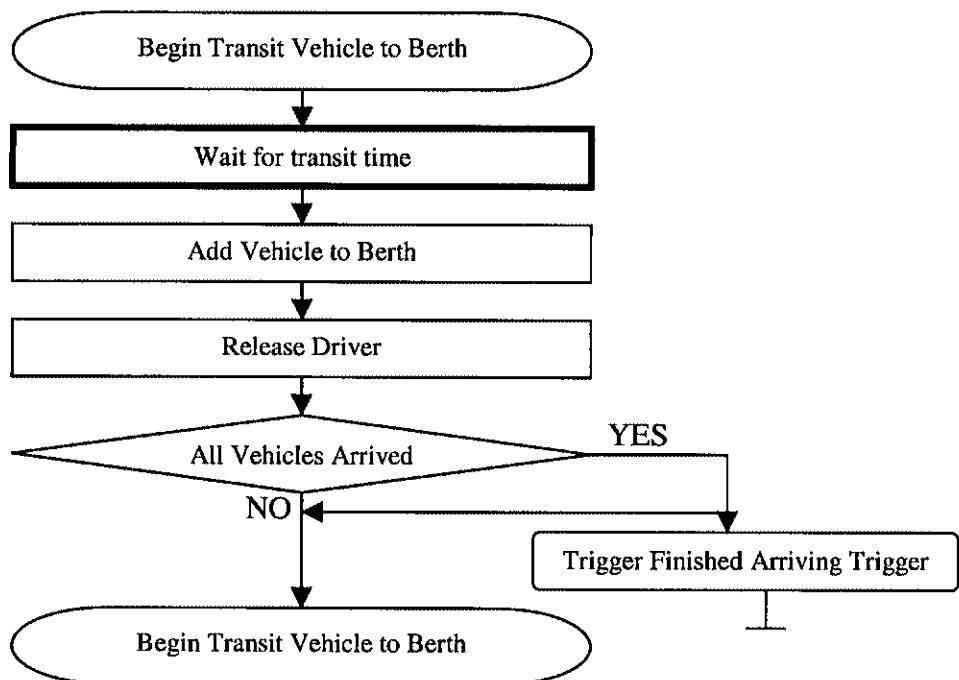


Fig. A.62

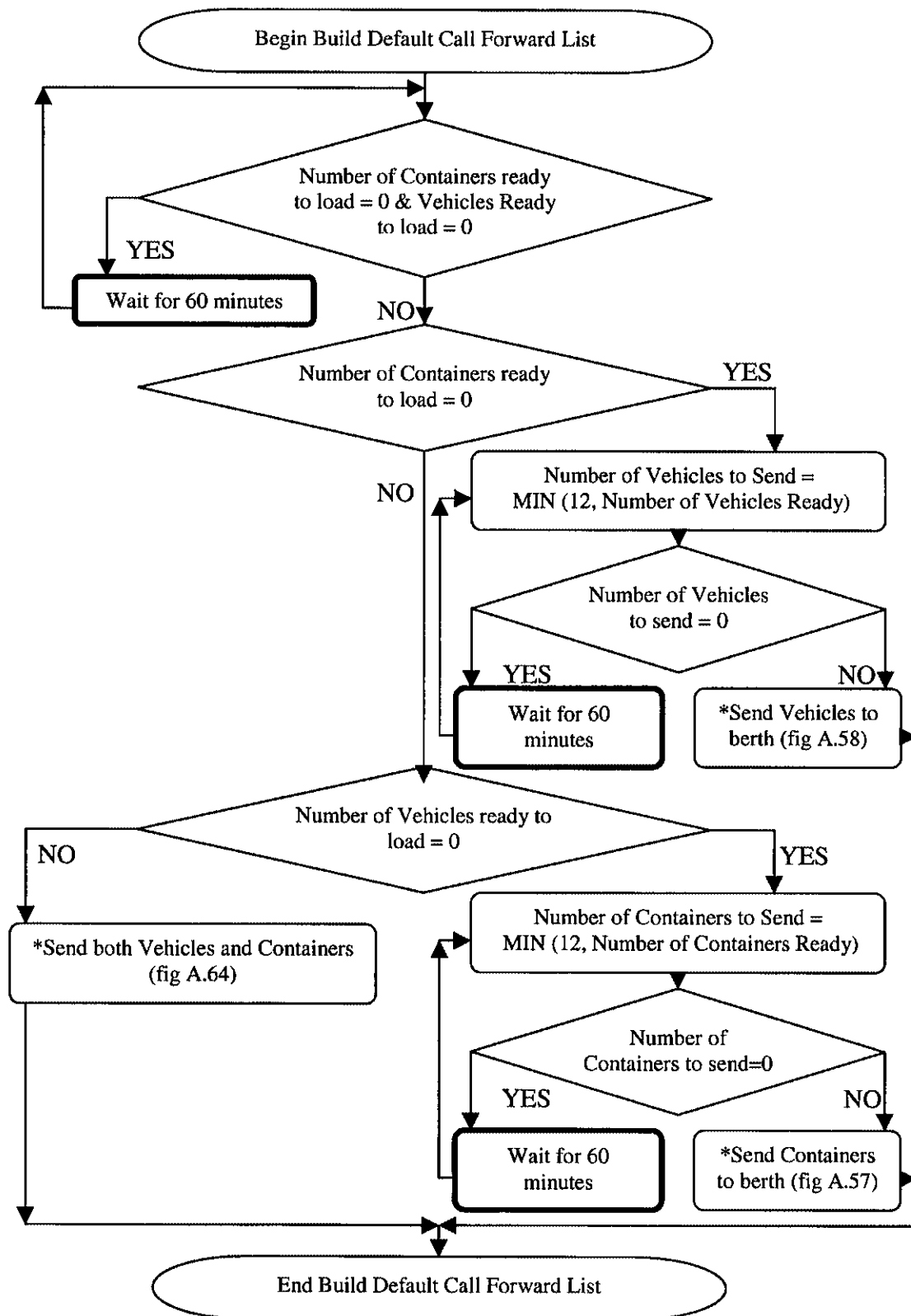


Fig. A.63

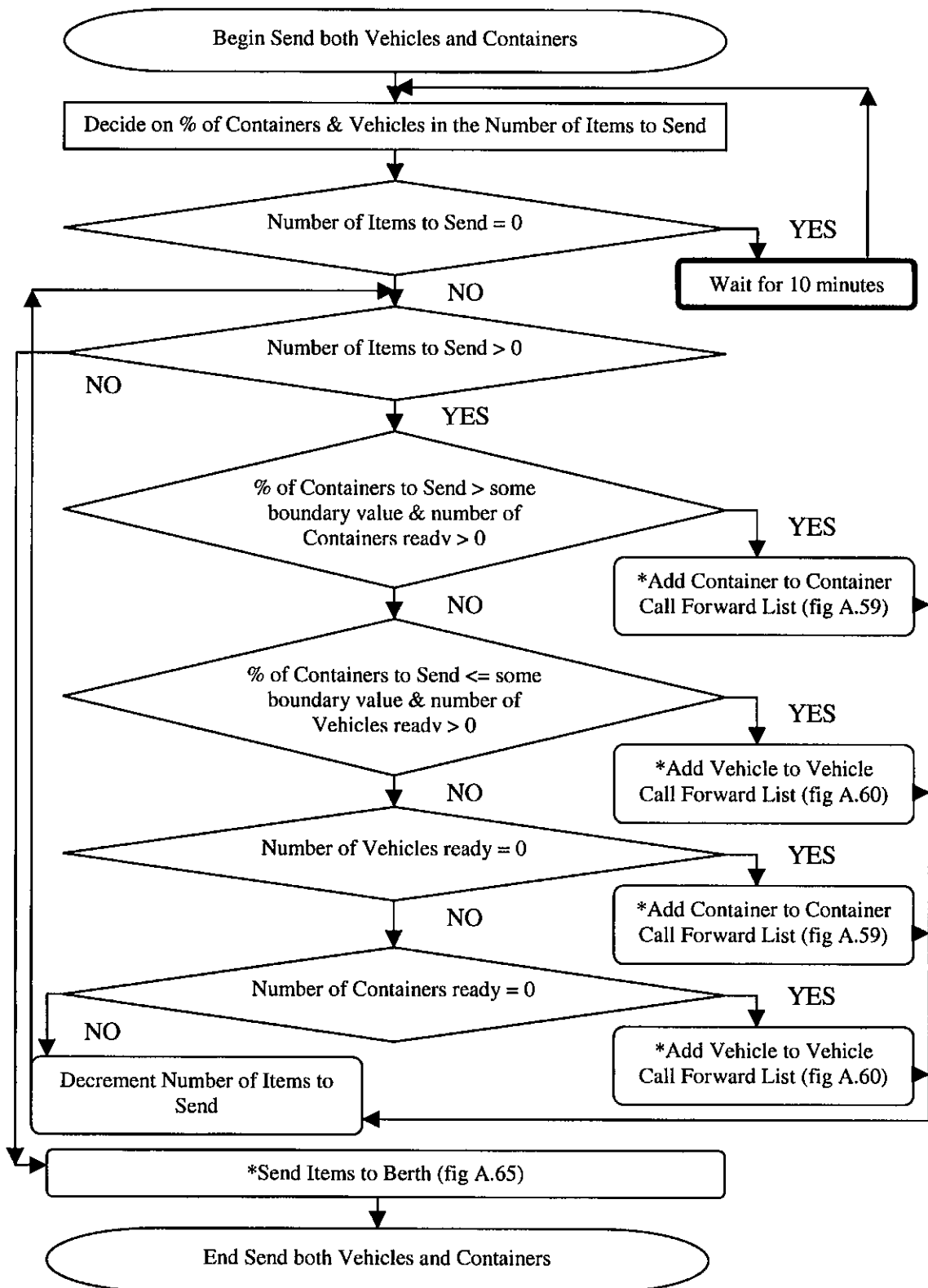


Fig. A.64

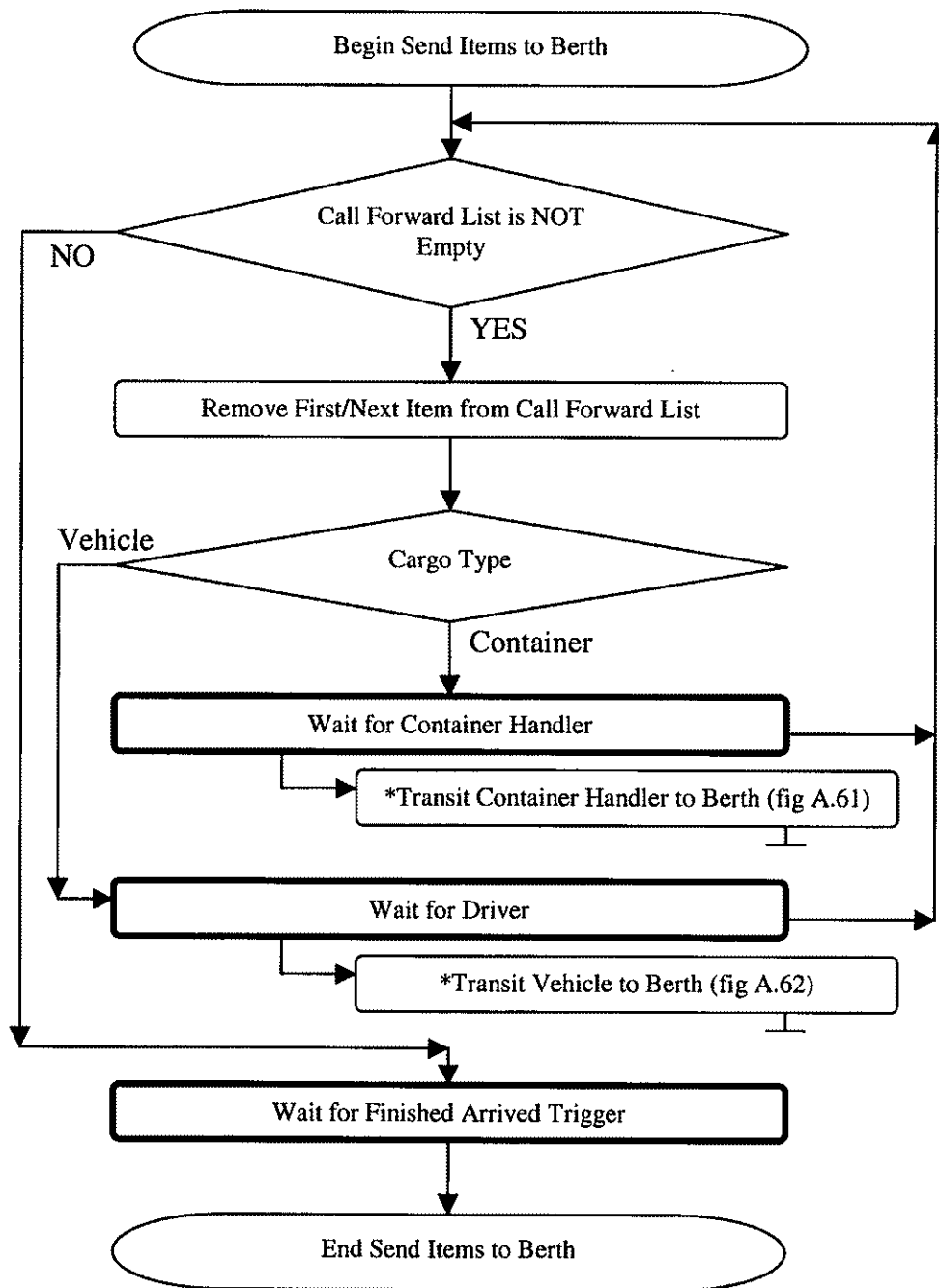


Fig. A.65

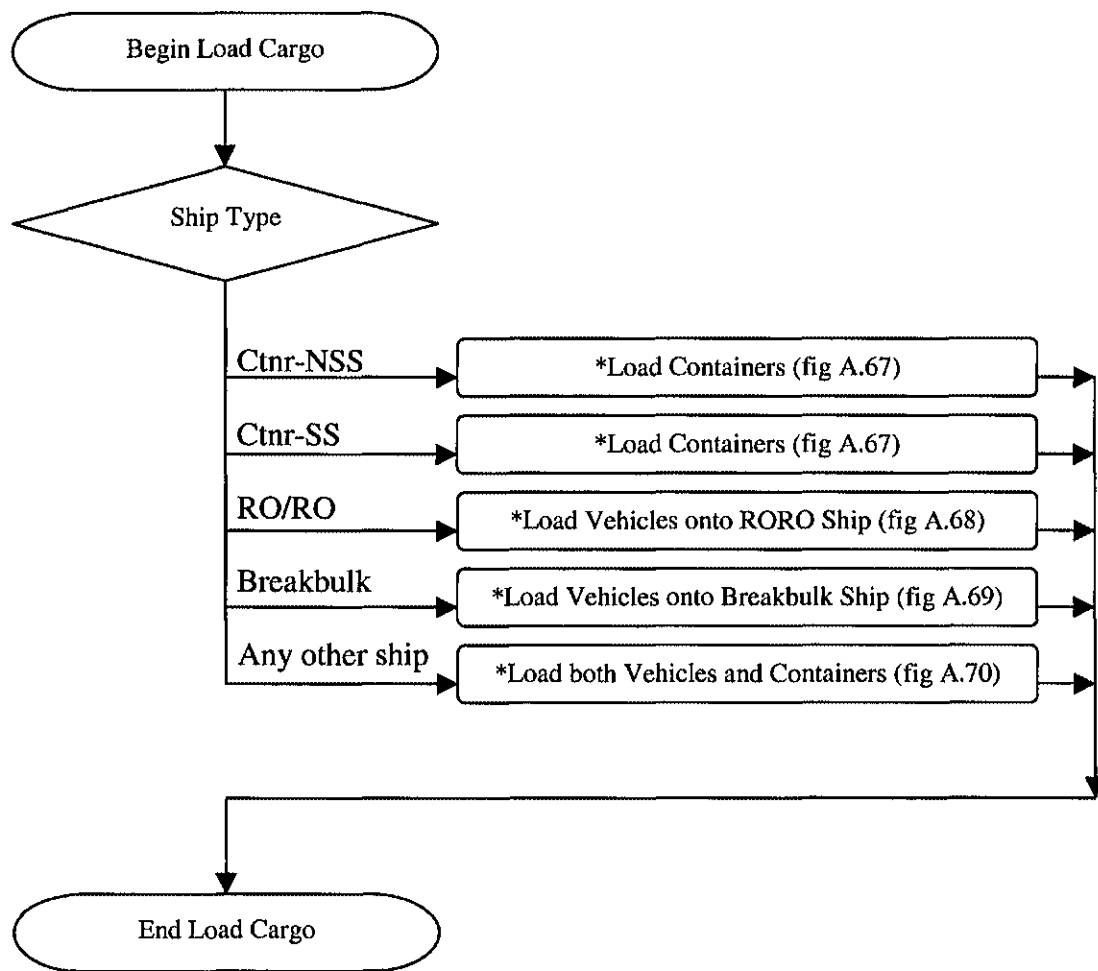


Fig. A.66

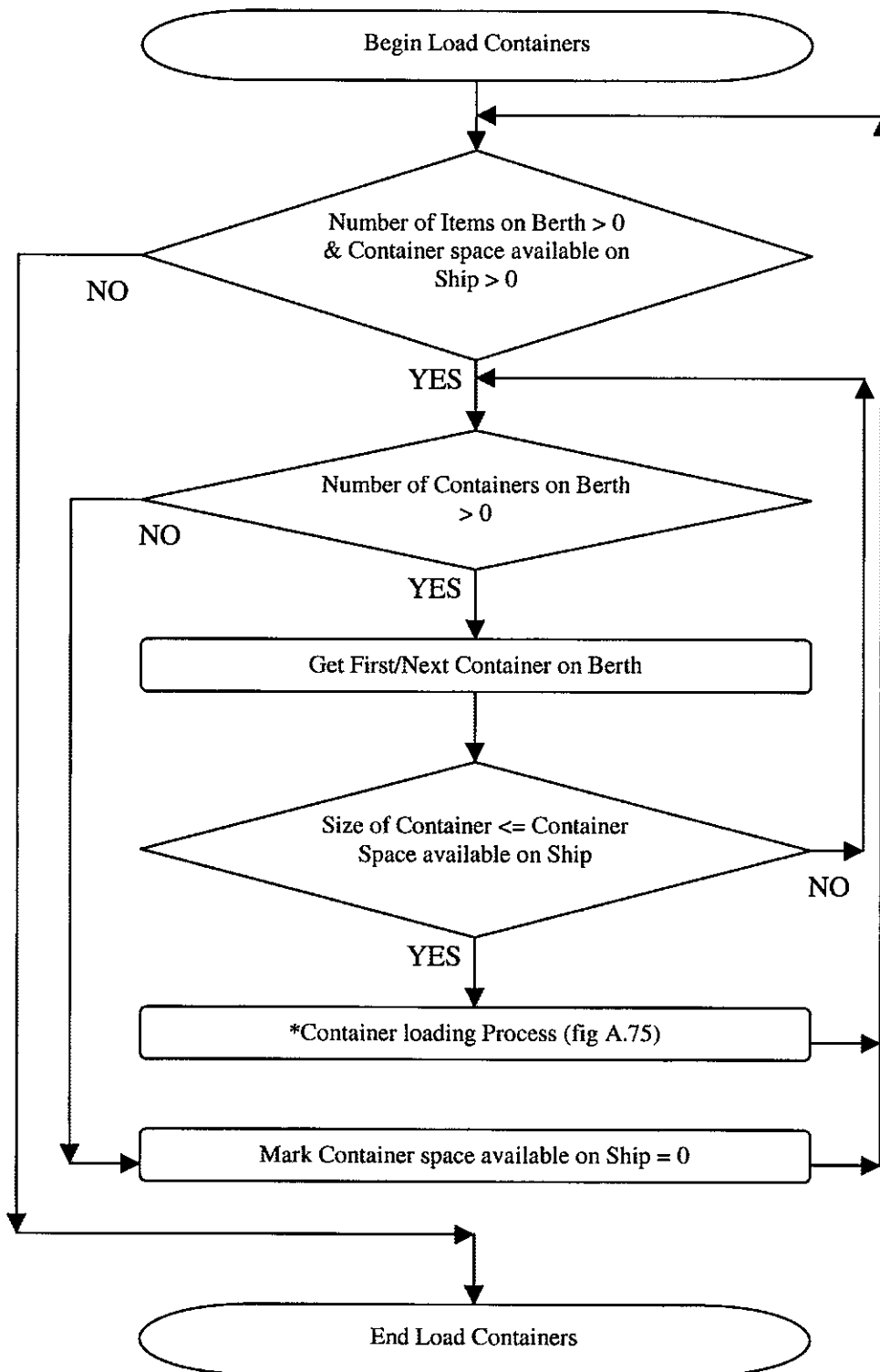


Fig. A.67

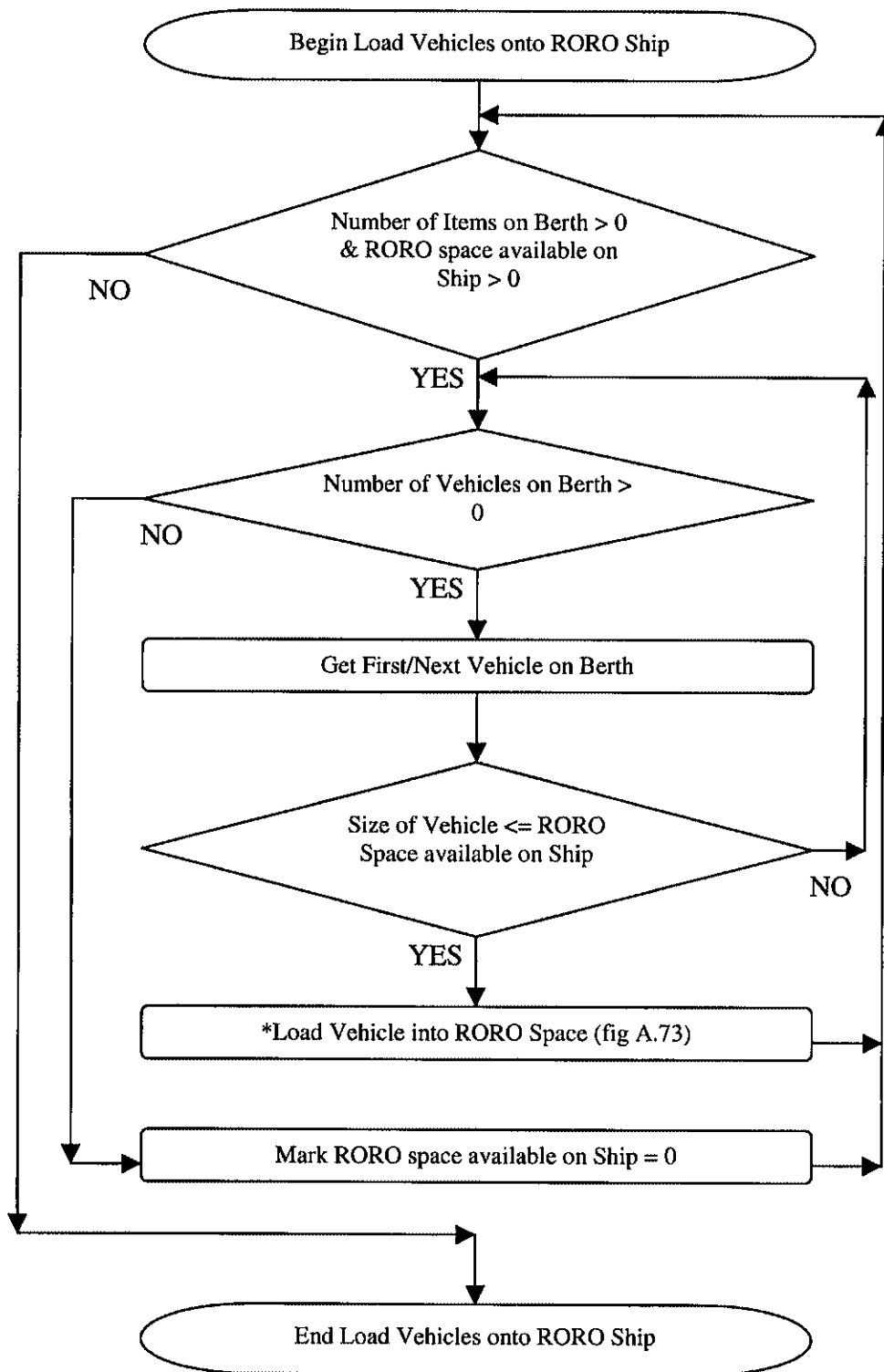


Fig. A.68



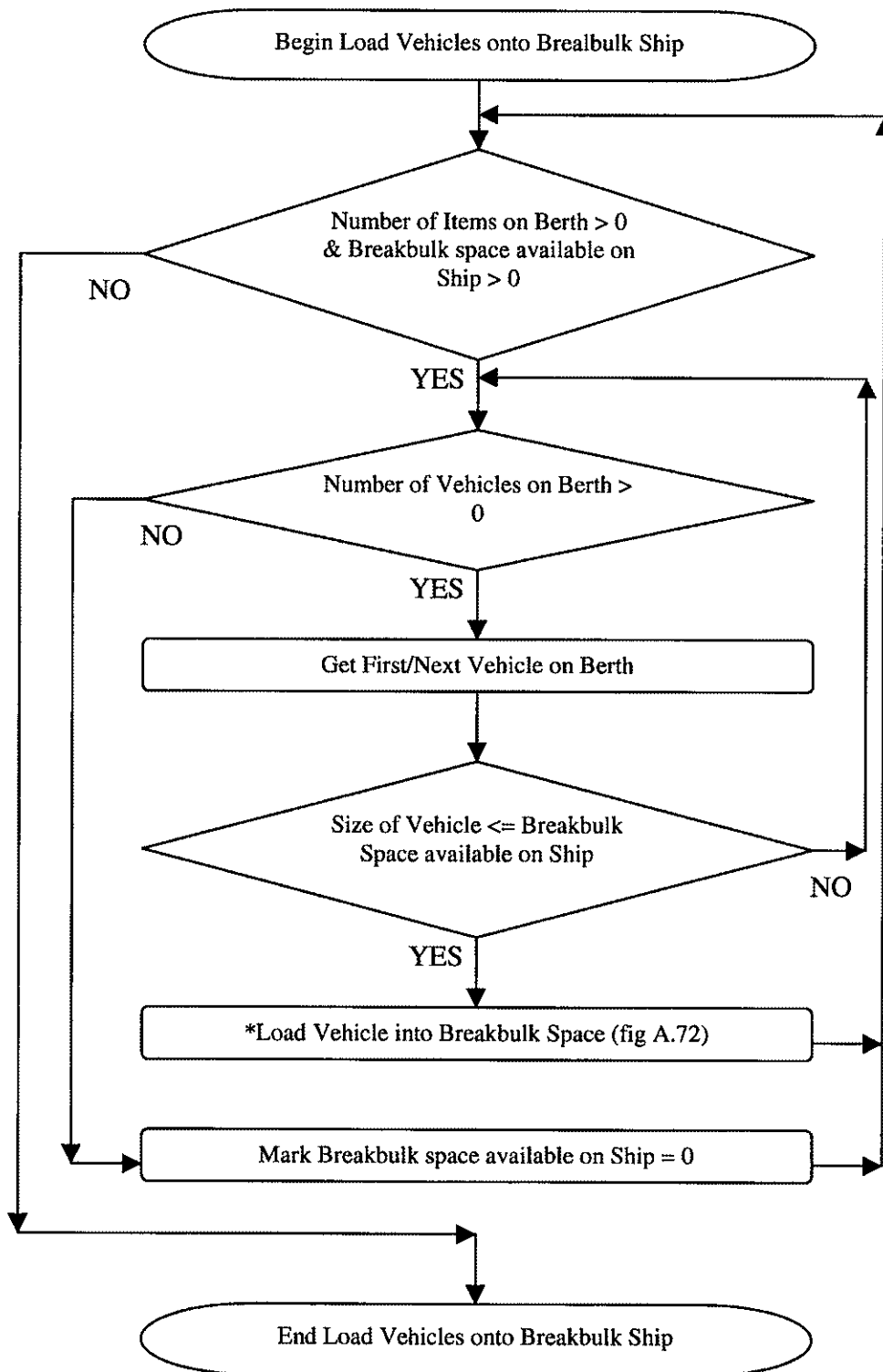


Fig. A.69

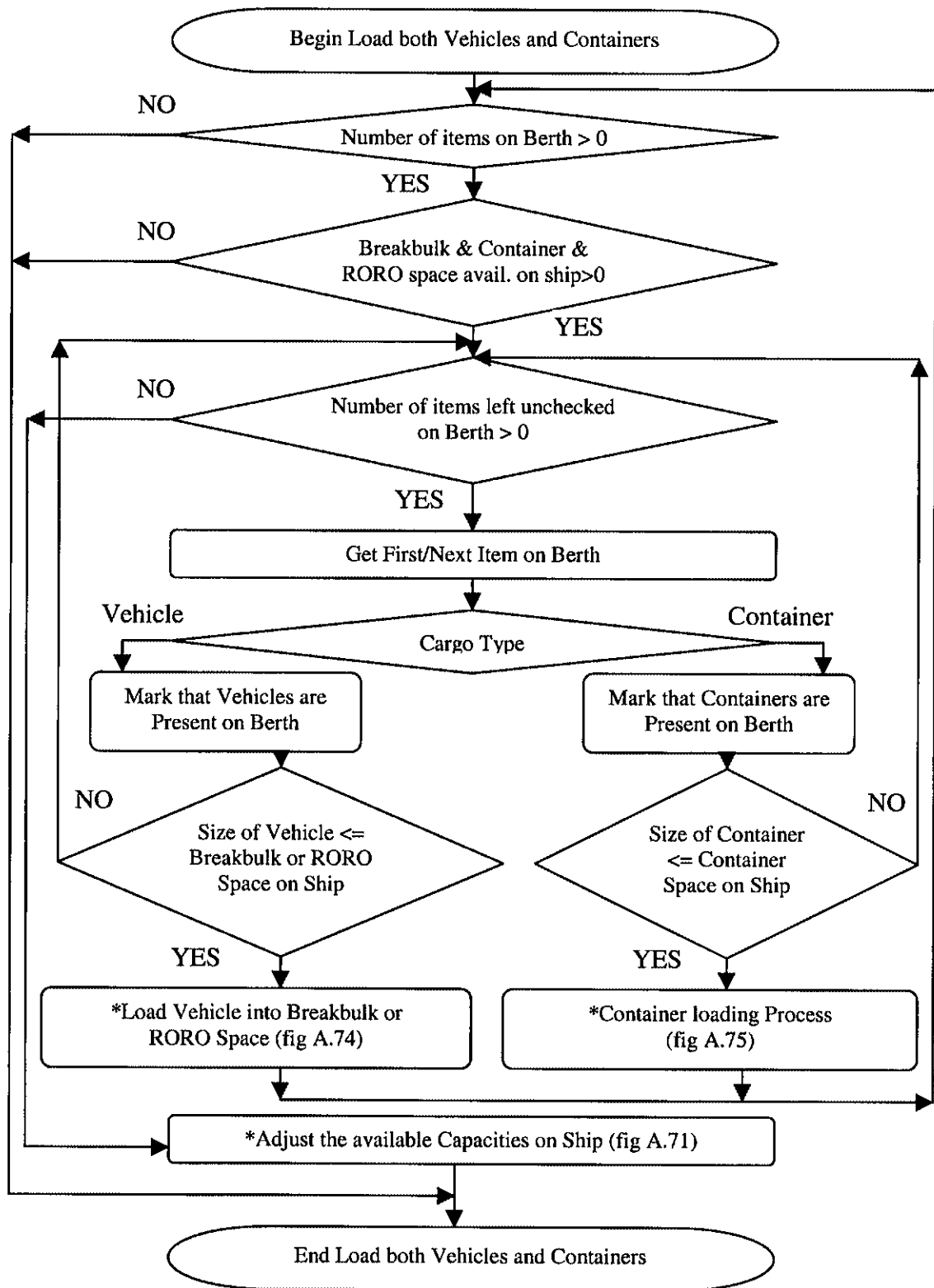


Fig. A.70

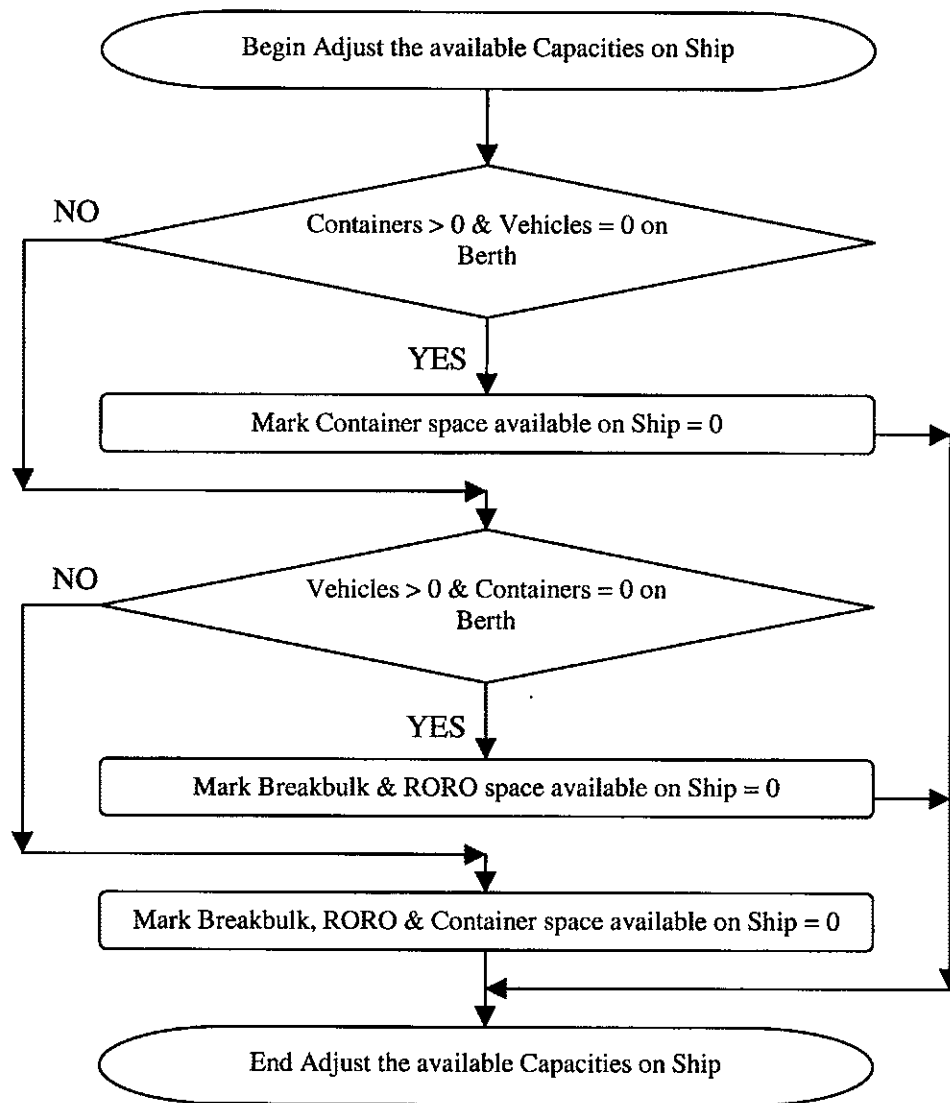


Fig. A.71

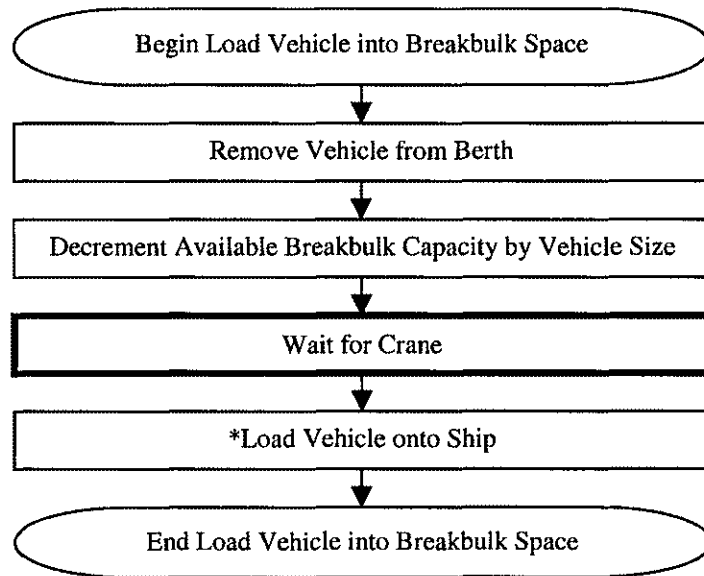


Fig. A.72

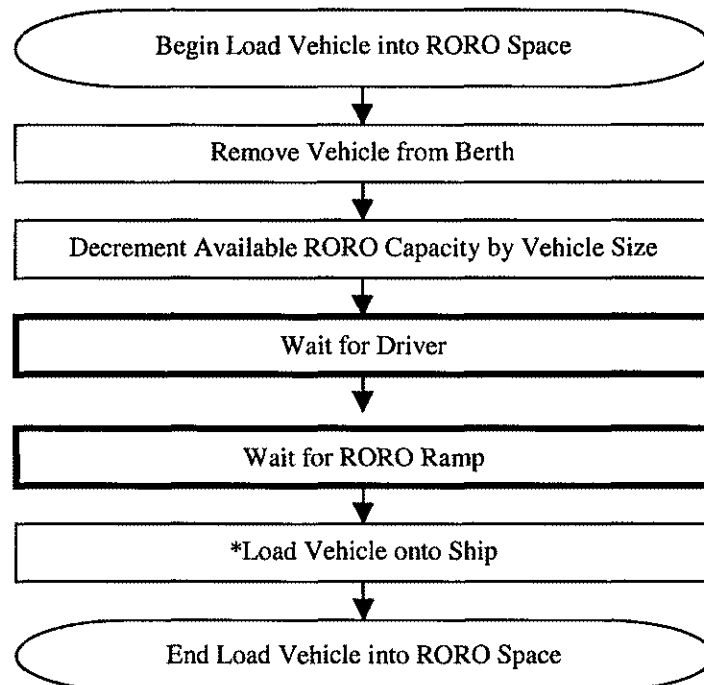


Fig. A.73

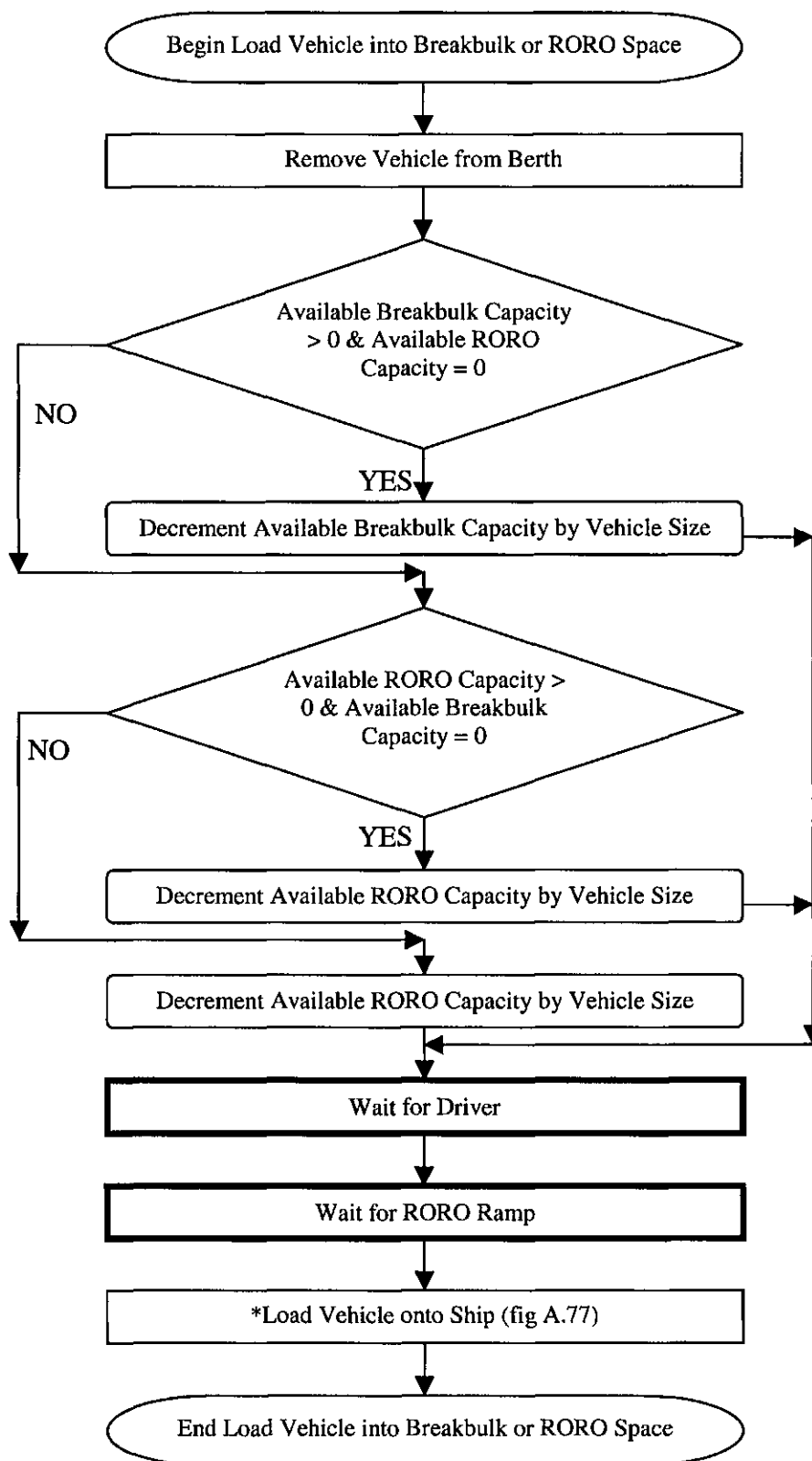


Fig. A.74

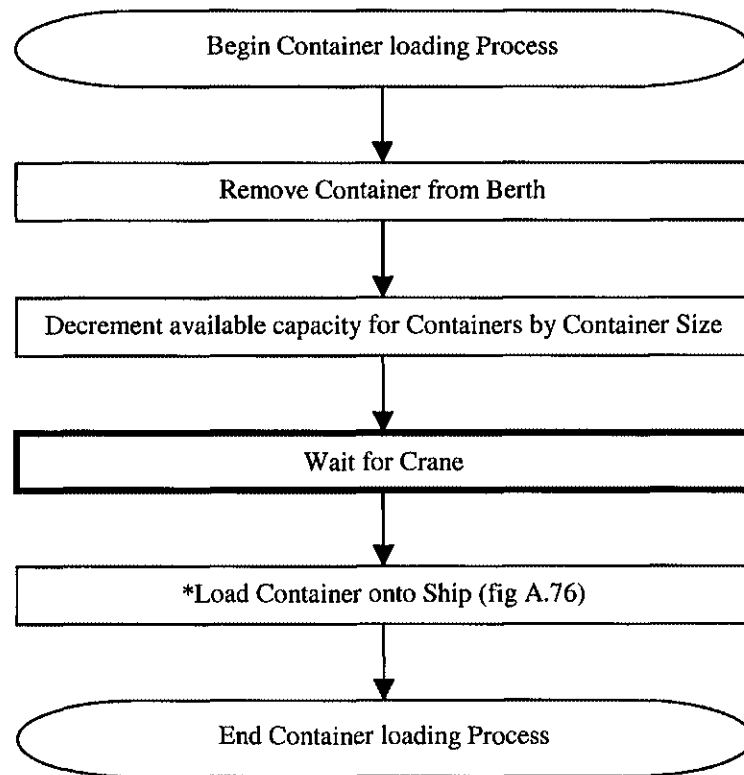


Fig. A.75

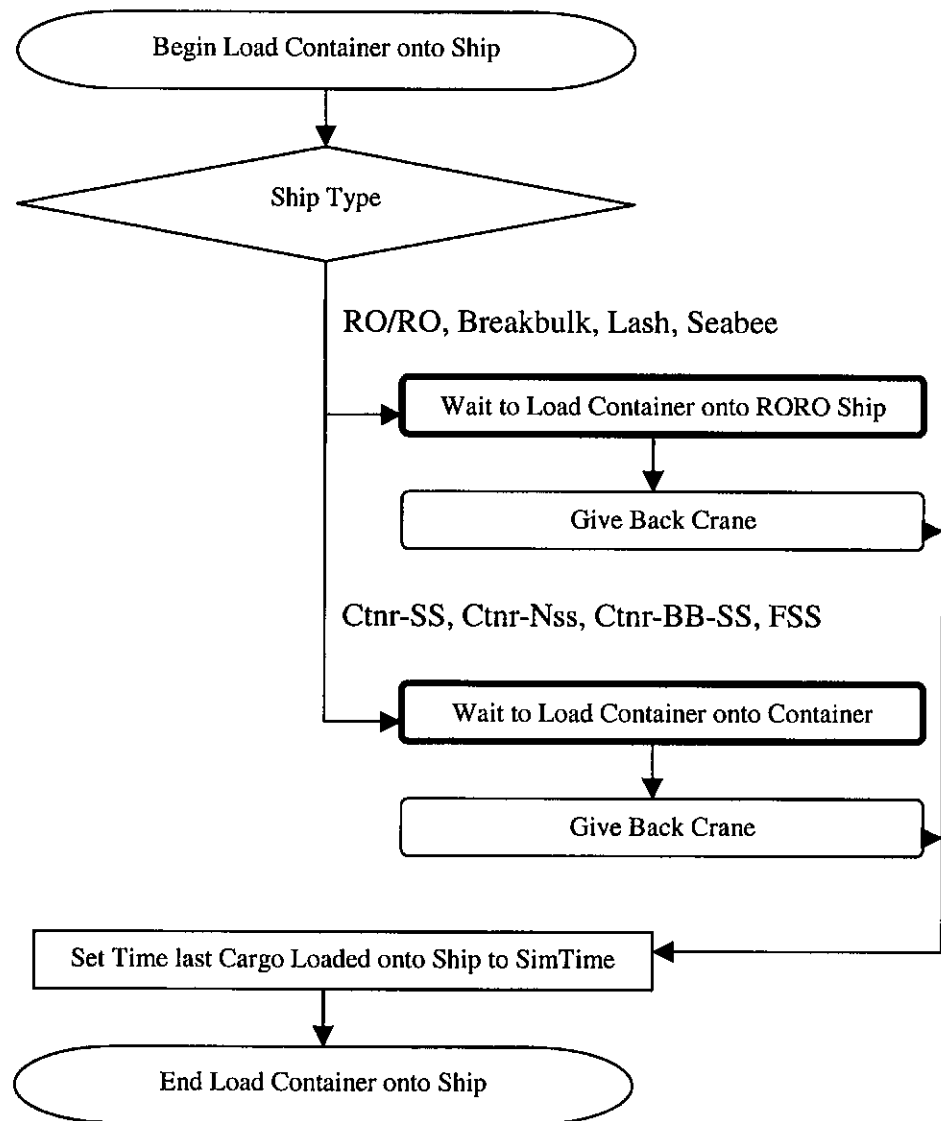


Fig. A.76

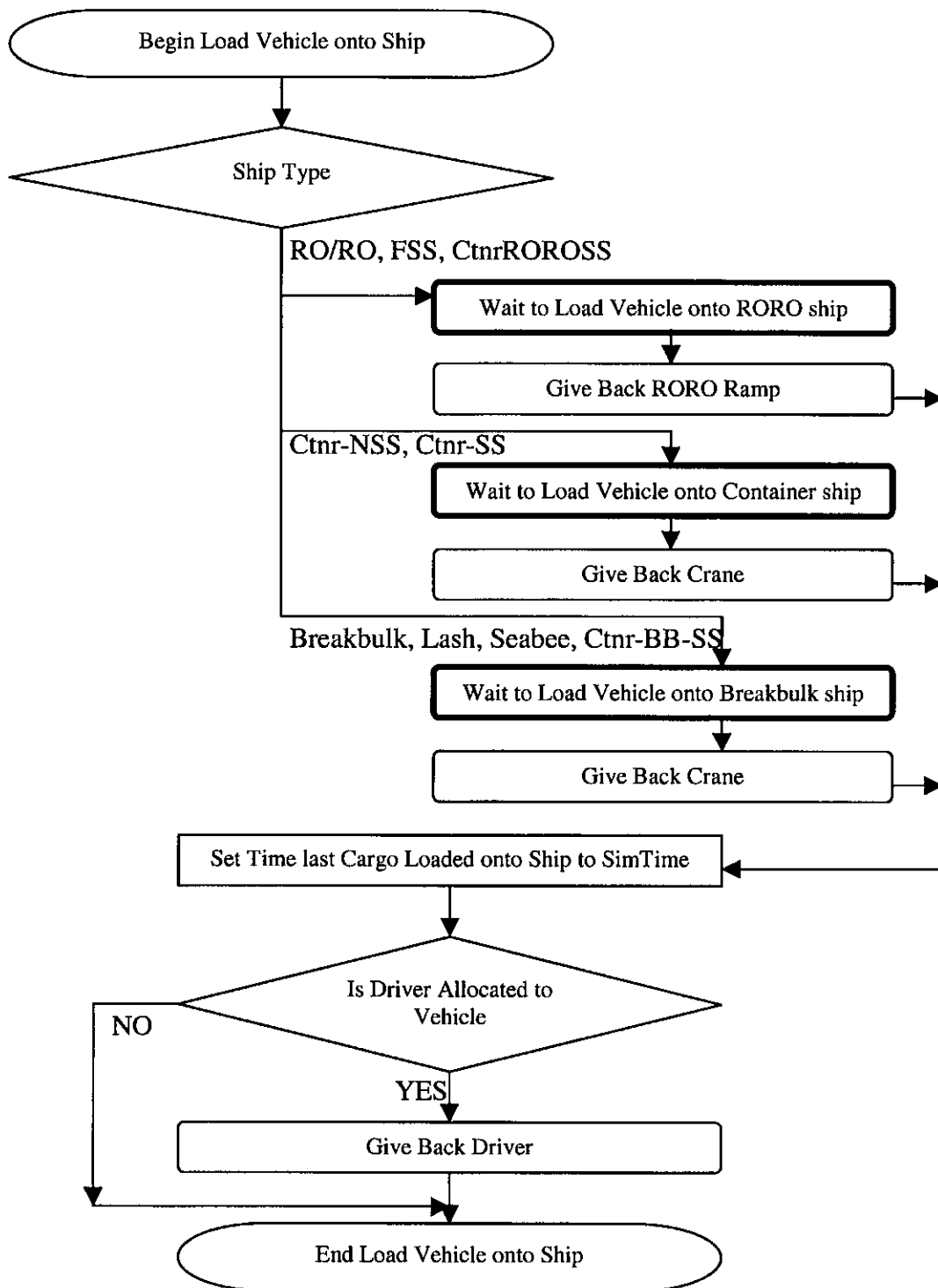


Fig. A.77



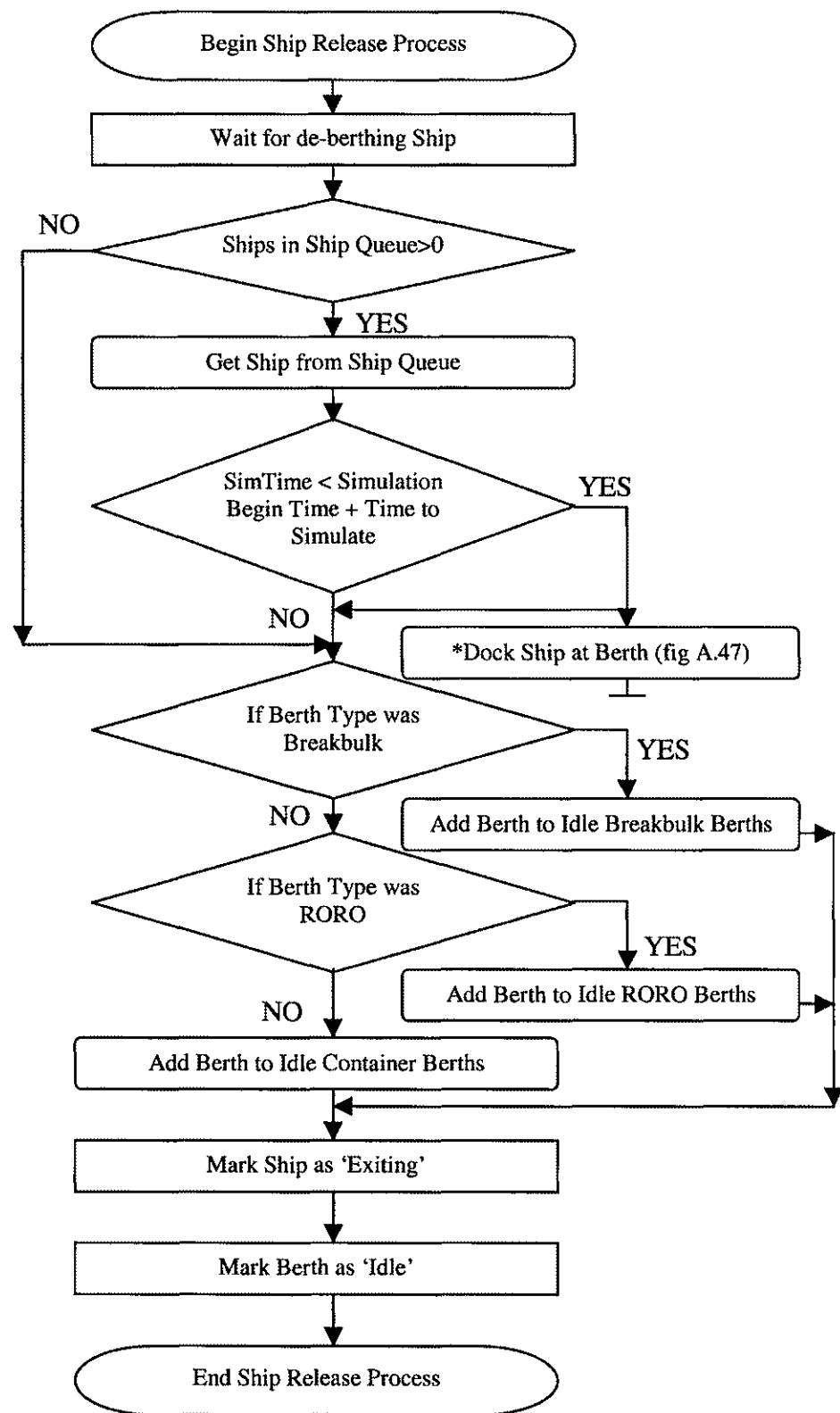


Fig. A.78

## APPENDIX B

### LIFE CYCLES OF GRAPHICAL WINDOW OBJECTS IN PORTSIM

The two important graphical window objects that are used extensively as a means of user interface are WindowObj and DialogBoxObj.

#### B.1 Life Cycle of instances of type WindowObj

This section shows the generic representation of life cycle of instances of type WindowObj used in PORTSIM.

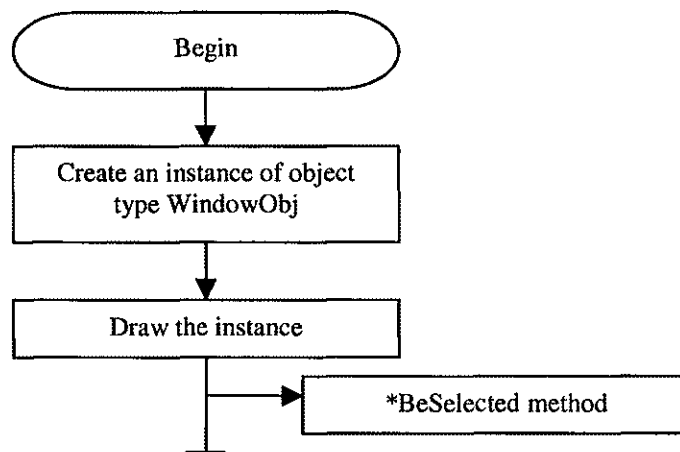


Fig. B.1 Generic representation of Life Cycle of instances of WindowObj

## B.2 Life Cycles of instances of type DialogBoxObj

There are two types of dialog box objects used in PORTSIM. One is of DialogBoxObj object type that SIMGRAPHICS provides and another is of object types that are derived from DialogBoxObj.

### *B.2.1. Life Cycles of instances of type DialogBoxObj*

This section shows the generic representation of the life cycles of instances of type DialogBoxObj object.

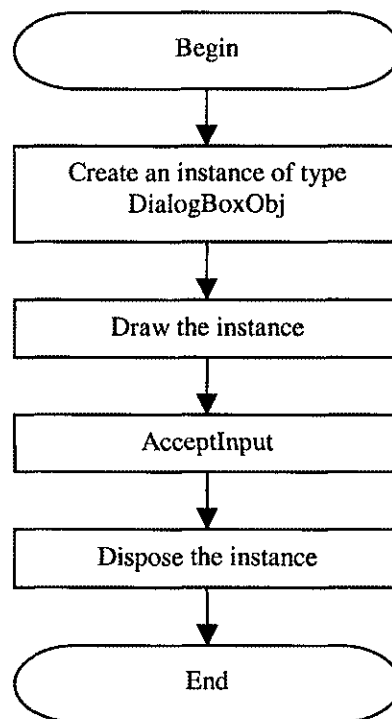


Fig. B.2 Life Cycle of instances of type DialogBoxObj with no processing depending on input

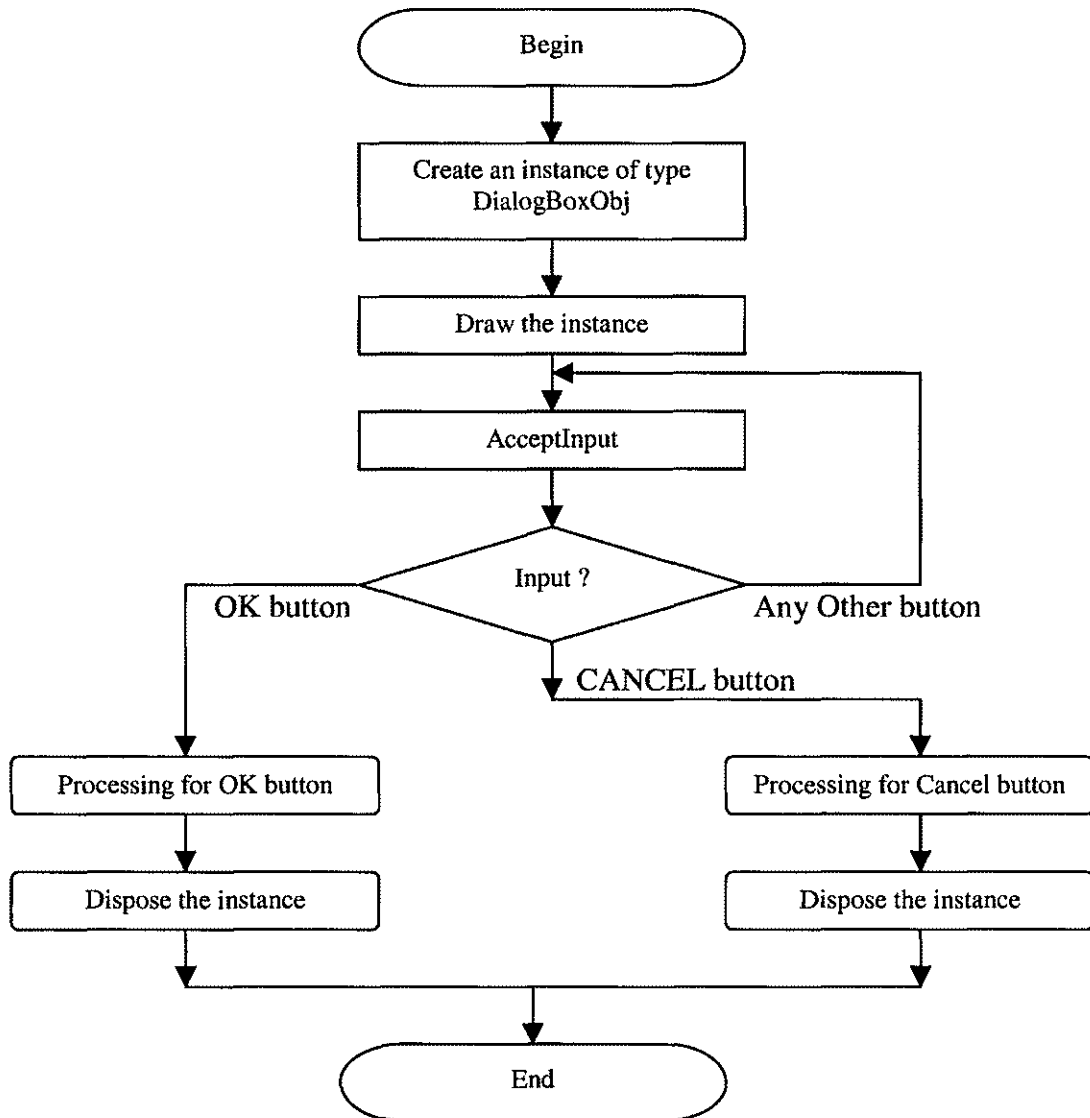


Fig. B.3 Life Cycle of instances of type DialogBoxObj with a loop and with unique processing for 'Ok' and 'Cancel' input

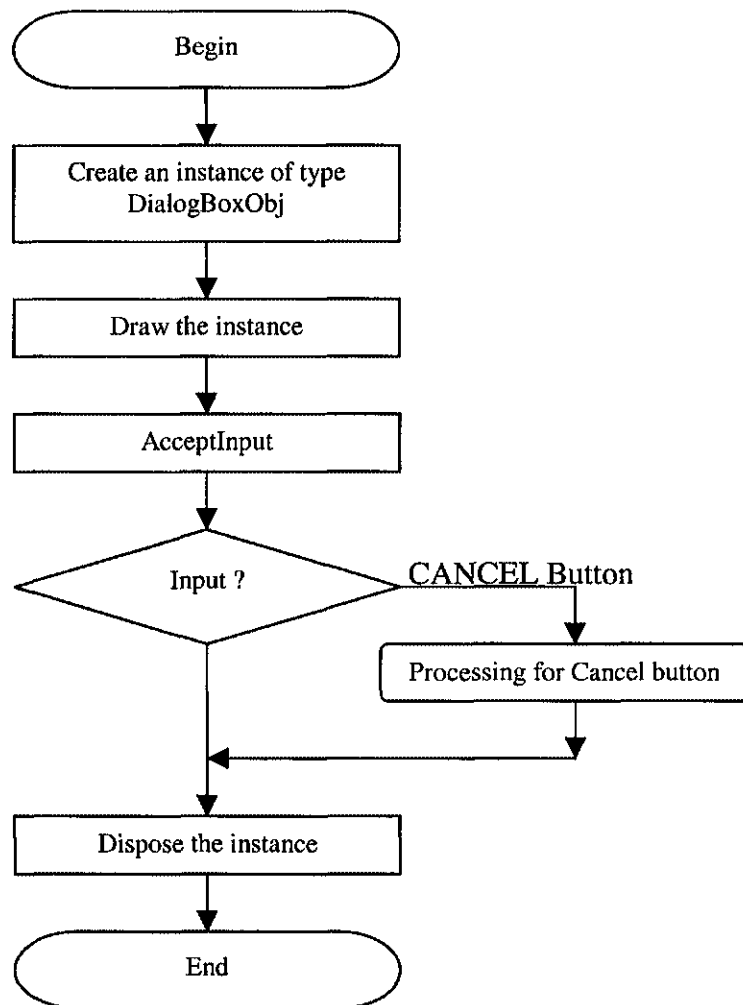


Fig. B.4 Life Cycle of instances of type DialogBoxObj with processing for 'Cancel' input alone

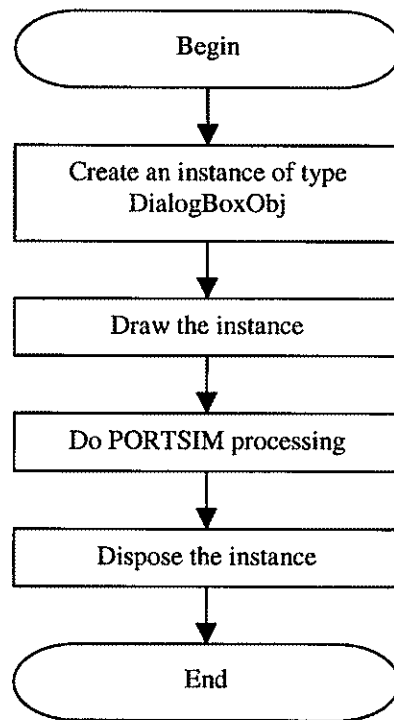


Fig. B.5 Life Cycle of instances of type DialogBoxObj with no provision for accepting input

### *B.2.2. Life Cycles of instances of Object types that are derived from DialogBoxObj*

This section shows the generic representation of the life cycles of instances of dialog boxes of object types that are derived from DialogBoxObj object.

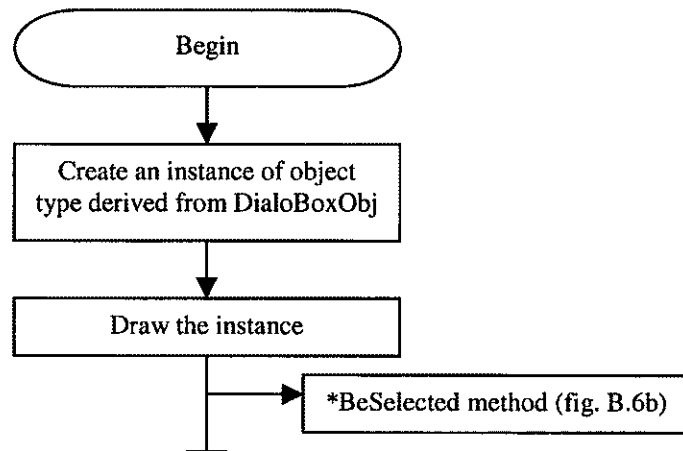


Fig. B.6a Life Cycle of instances of type derived from DialogBoxObj using BeSelected method to accept input

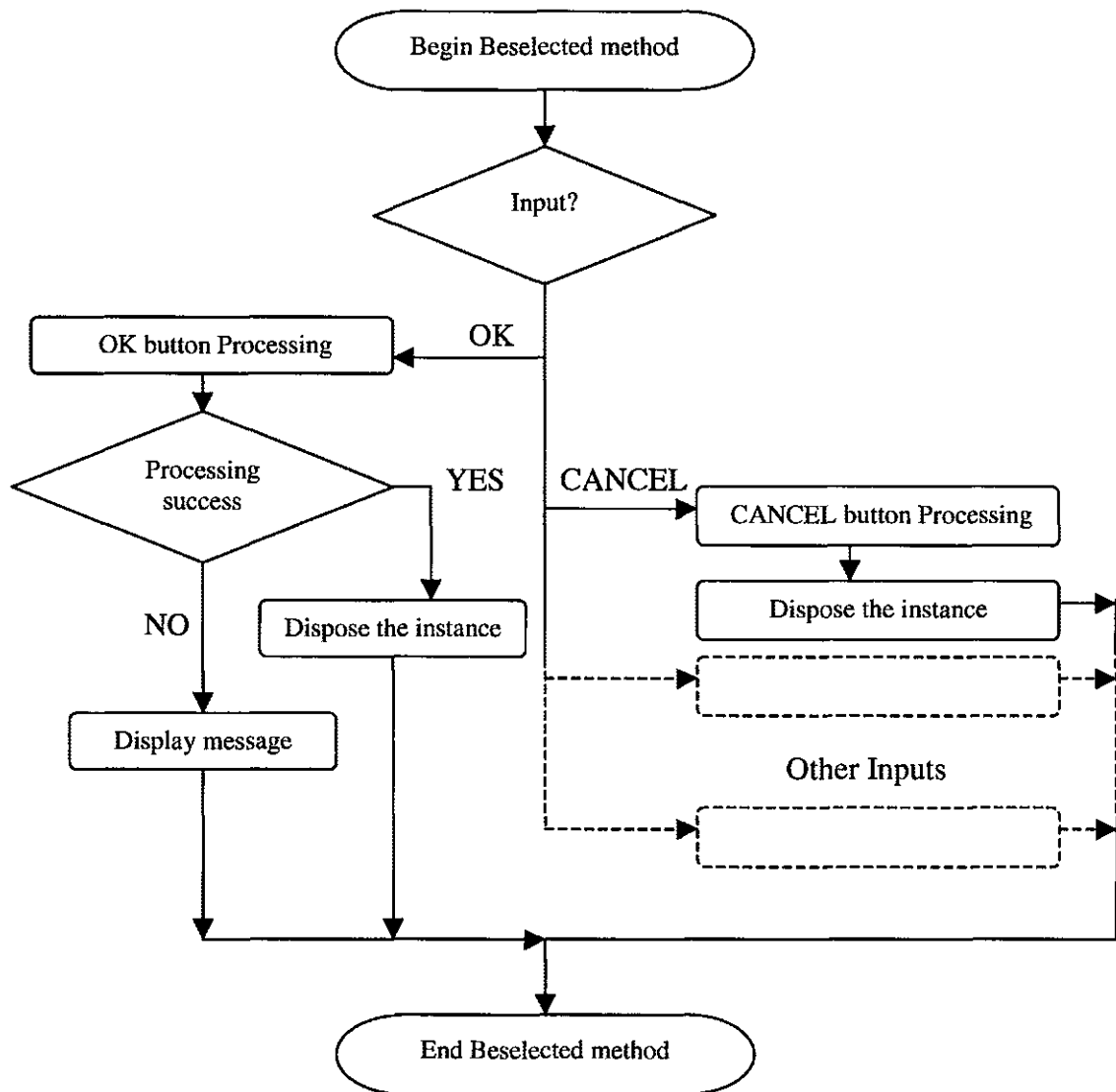


Fig. B.6b Generic representation of Program Flow within the BeSelected method of object types shown in Fig. B.6a



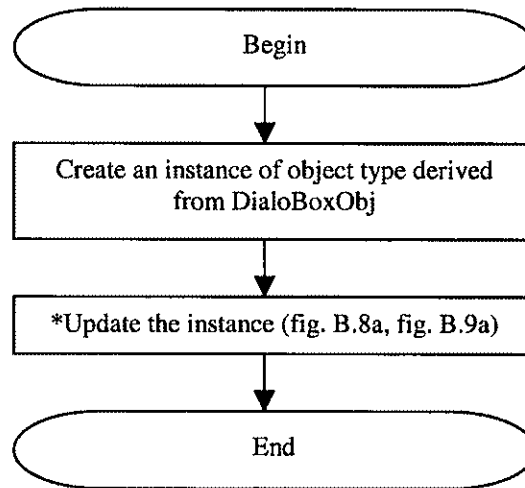


Fig. B.7 Life Cycle of instances of type derived from DialogBoxObj using 'Update' function

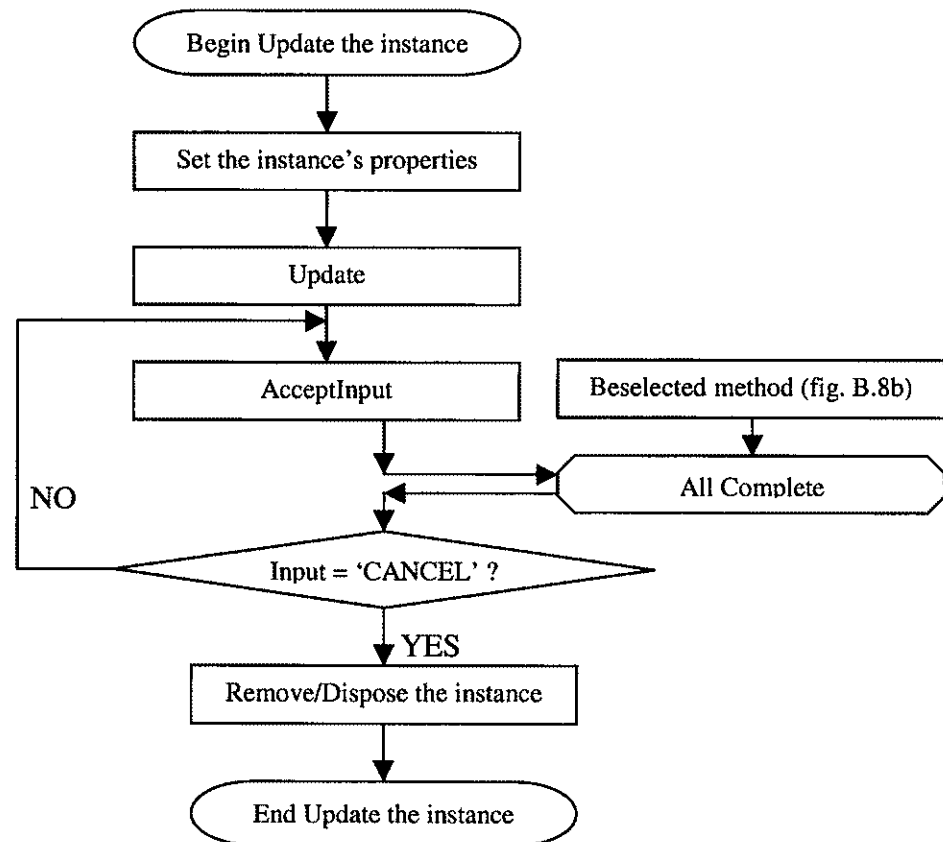


Fig. B.8a One of the generic representations of 'Update' functions shown in Fig. B.7

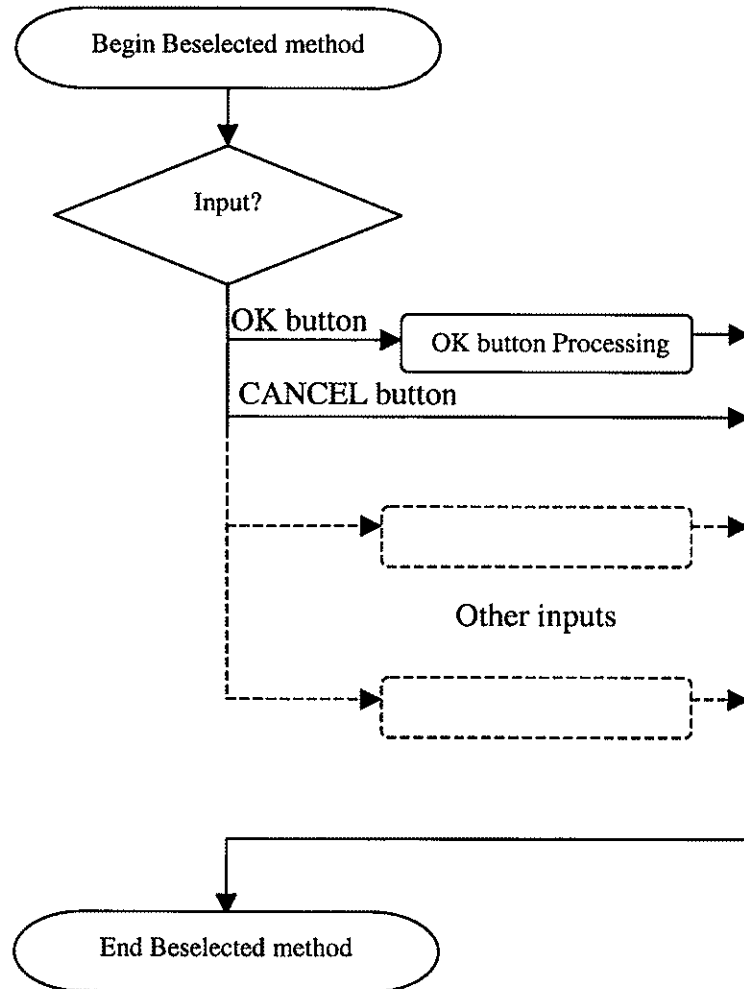


Fig. B.8b Generic representation of Program Flow within the BeSelected method of objects having 'Update' function as shown in Fig. B.8a

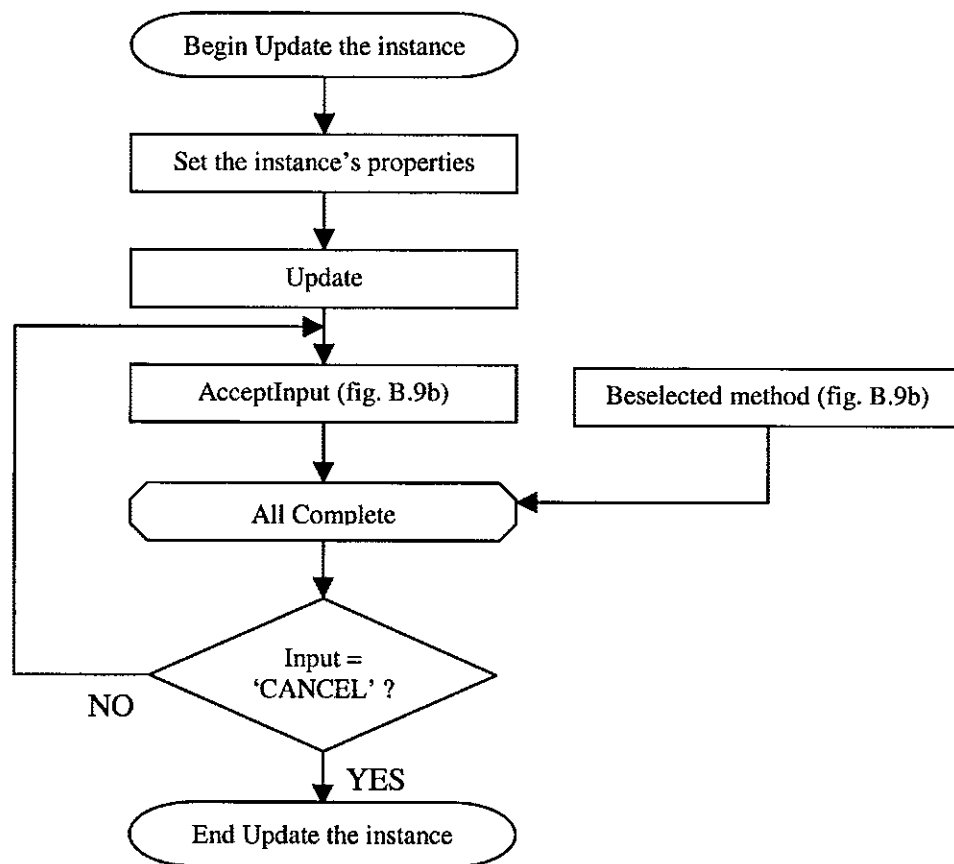


Fig. B.9a One of the generic representations of 'Update' functions shown in Fig. B.7

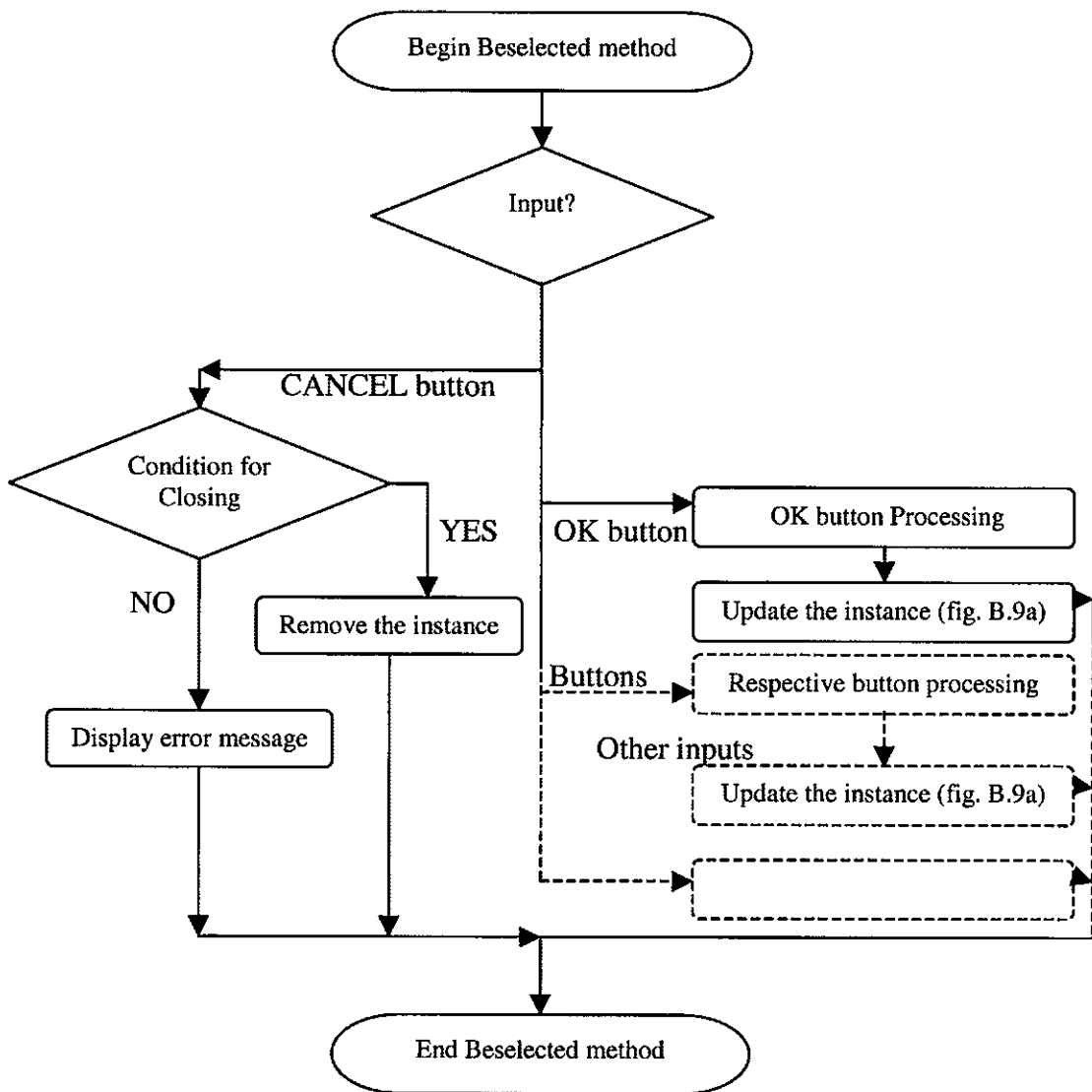


Fig. B.9b Generic representation of Program Flow within the BeSelected method of objects having 'Update' function as shown in Fig. B.9a

## APPENDIX C

### PORTSIM CODE FOR RELATED FLOW DIAGRAMS

The following Figures represent the actual PORTSIM code for modified flow diagrams, except Fig. C.1, of Appendix A and Appendix B in Sections 3 and 4.

```

TELL METHOD callForwardFlatcars(INinterchangeYard:interchangeYardObj;IN
ranGen : RandomObj);
VAR
  lo,hi    : REAL;
  temp     : REAL;
  stringOfFlatcars : trainObj;
  railcar   : railcarObj;
  tangentLengthAvail : INTEGER;
  len       : INTEGER;
  j         : INTEGER;
  tempTime  : REAL;
  svcTime   : REAL;
  typeNeeded : BOOLEAN;
  numFlatcarsCarryingVehicles : INTEGER;
  numFlatcarsCarryingContainers : INTEGER;
  cargo     : cargoObj;
BEGIN
  WHILE (SimTime() <= (simBeginTime + timeToSimulate))
  {If no flatcars in interchange yard,}
  {wait until a train has been classified in the
  interchange yard}
  {This signifies that railcars are available to call
  forward}
  IF (ordering = "Vehicles/Unit Equipment Only")
  IF (poeMode)
    {Vehicles/Unit Equipment only accepted}
    numFlatcarsCarryingVehicles := ASK interchangeYard TO
                                countFlatcarsCarryingVehicles();
    WHILE (numFlatcarsCarryingVehicles = 0)
      WAIT FOR trainClassifiedTrigger TO Fire;
      ON INTERRUPT
      END WAIT;
      numFlatcarsCarryingVehicles := ASK interchangeYard TO
                                countFlatcarsCarryingVehicles();
    END WHILE;
  END IF;
  ELSIF (ordering = "Containers Only")
  IF (poeMode)
    {Containers only accepted}
    numFlatcarsCarryingContainers := ASK interchangeYard TO
    countFlatcarsCarryingContainers();
    WHILE (numFlatcarsCarryingContainers = 0)
      WAIT FOR trainClassifiedTrigger TO Fire;
      ON INTERRUPT
      END WAIT;

```

```

        numFlatcarsCarryingContainers := ASK interchangeYard TO
        countFlatcarsCarryingContainers();
    END WHILE;
END IF;
ELSE
    {Both Vehicles/Unit Equipment and Containers are accepted}
END IF;
IF (status = "idle")
    {Set status to busy}
    ASK SELF TO setStatus("busy");
    tempTime := SimTime();
    svcTime := 0.0;
    {Build the appropriate stringOfFlatcars to send to spur}
    NEW(stringOfFlatcars);
    WHILE (stringOfFlatcars.numberIn = 0)
        ASK SELF TO buildStringOfFlatcar(interchangeYard,stringOfFlatcars);
        IF (stringOfFlatcars.numberIn = 0)
            {OUTPUT("Waiting for more items to be ready in
            emptyRailcarStorage");}
            WAIT DURATION 30.0;
            END WAIT;
        END IF;
    END WHILE;
    {Let the system know that more items can be put in interchangeYard}
    TELL IYTrigger TO Trigger;
    {Obtain a locomotive resource and transit it to the IY}
    WAIT FOR port.locomotiveQ TO Give(stringOfFlatcars,1);
    {Processing Time to switch locomotive to IY}
    IF (useProcessDistributionsForTransitTimes)
        lo := switchSpurToIYTime.actualValue-
            switchSpurToIYStdDev.actualValue;
        hi := switchSpurToIYTime.actualValue +
            switchSpurToIYStdDev.actualValue;
        temp := ranGen.UniformReal(lo,hi);
        IF (temp < 0.0)
            temp := 0.0;
        END IF;
    ELSE
        {Calculate the time based upon speed and distance of route}
        {NEED TO ADD THIS}
        {Until this is added, make transit time 0.0}
        temp := 0.0;
    END IF;
    WAIT DURATION temp;
    END WAIT;

    {Processing Time to couple flatcars at IY}
    lo := coupleAtIYTime.actualValue - coupleAtIYStdDev.actualValue;
    hi := coupleAtIYTime.actualValue + coupleAtIYStdDev.actualValue;
    temp := ranGen.UniformReal(lo,hi);
    IF (temp < 0.0)
        temp := 0.0;
    END IF;
    WAIT DURATION temp;
    END WAIT;

    IF (type = "OPEN STAGING") OR (type = "Open Staging") OR

```

```

    (type = "OPEN STORAGE") OR (type = "Open Storage")
{Processing Time to switch flatcars to spur}
IF (useProcessDistributionsForTransitTimes)
    lo := switchToSpurTime.actualValue -
        switchToSpurStdDev.actualValue;
    hi := switchToSpurTime.actualValue +
        switchToSpurStdDev.actualValue;
    temp := ranGen.UniformReal(lo,hi);
    IF (temp < 0.0)
        temp := 0.0;
    END IF;
ELSE
    {Calculate the time based upon speed and distance of route}
    {NEED TO ADD THIS}
    {Until this is added, make transit time 0.0}
    temp := 0.0;
END IF;
WAIT DURATION temp;
END WAIT;
ELSIF (type = "COVERED STAGING") OR (type = "Covered Staging")
{Processing Time to switch flatcars to dock}
IF (useProcessDistributionsForTransitTimes)
    lo := switchToDockTime.actualValue -
        switchToDockStdDev.actualValue;
    hi := switchToDockTime.actualValue +
        switchToDockStdDev.actualValue;
    temp := ranGen.UniformReal(lo,hi);
    IF (temp < 0.0)
        temp := 0.0;
    END IF;
ELSE
    {Calculate the time based upon speed and distance of route}
    {NEED TO ADD THIS}
    {Until this is added, make transit time 0.0}
    temp := 0.0;
END IF;
WAIT DURATION temp;
END WAIT;
ELSE
    {type = "APRON" or type = "Apron"}
    {Processing Time to switch flatcars to Berth}
    IF (useProcessDistributionsForTransitTimes)
        lo := switchToBerthTime.actualValue -
            switchToBerthStdDev.actualValue;
        hi := switchToBerthTime.actualValue +
            switchToBerthStdDev.actualValue;
        temp := ranGen.UniformReal(lo,hi);
        IF (temp < 0.0)
            temp := 0.0;
        END IF;
    ELSE
        {Calculate the time based upon speed and distance of route}
        {NEED TO ADD THIS}
        {Until this is added, make transit time 0.0}
        temp := 0.0;
    END IF;
    WAIT DURATION temp;

```

```

    END WAIT;
END IF;
{Processing Time to uncouple flatcars at spur}
lo := uncoupleAtSpurTime.actualValue -
    uncoupleAtSpurStdDev.actualValue;
hi := uncoupleAtSpurTime.actualValue +
    uncoupleAtSpurStdDev.actualValue;
temp := ranGen.UniformReal(lo,hi);
IF (temp < 0.0)
    temp := 0.0;
END IF;
WAIT DURATION temp;
END WAIT;
END WAIT;
{Make the locomotive available again}
ASK port.locomotiveQ TO TakeBack(stringOfFlatcars,1);
IF (poeMode)
    {Wait for vehicles/containers to be discharged from flatcars}
    WAIT FOR SELF TO dischargeFlatcarContents(stringOfFlatcars,ranGen);
    END WAIT;
    DISPOSE(stringOfFlatcars);
    {Cycle out empty flatcars}
    {NOTE: This code needs to be added}
END IF;
{Set status "idle" to signify that spur is ready for next string}
{of railcars}
ASK SELF TO setStatus("idle");
{Update operational svcTime parameters}
svcTime := SimTime() - tempTime;
totalBusyTime := totalBusyTime + svcTime;
IF svcTime > maxServiceTime
    maxServiceTime := svcTime;
END IF;

TELL spurDoneTrigger TO Trigger;
ELSE
    WAIT FOR spurDoneTrigger TO Fire;
    END WAIT;
END IF;
END WHILE;
END METHOD {callForwardFlatcars};

```

Fig. C.1 PORTSIM code for "callForwardFlatcars" method



```

TELL METHOD callForwardFlatcars(INinterchangeYard;interchangeYardObj;IN
ranGen : RandomObj);
VAR
  lo,hi    : REAL;
  temp     : REAL;
  stringOfFlatcars : trainObj;
  railcar   : railcarObj;
  tangentLengthAvail : INTEGER;
  len       : INTEGER;
  j         : INTEGER;
  tempTime  : REAL;
  svcTime   : REAL;
  typeNeeded : BOOLEAN;
  numFlatcarsCarryingVehicles : INTEGER;
  numFlatcarsCarryingContainers : INTEGER;
  cargo     : cargoObj;
BEGIN
  WHILE (SimTime() <= (simBeginTime + timeToSimulate))
  {If no flatcars in interchange yard,}
  {wait until a train has been classified in the
  interchange yard}
  {This signifies that railcars are available to call
  forward}
  IF (ordering = "Vehicles/Unit Equipment Only")
  IF (poeMode)
    {Vehicles/Unit Equipment only accepted}
    numFlatcarsCarryingVehicles := ASK interchangeYard TO
                                countFlatcarsCarryingVehicles();
    WHILE (numFlatcarsCarryingVehicles = 0)
      WAIT FOR trainClassifiedTrigger TO Fire;
      ON INTERRUPT
      END WAIT;
      numFlatcarsCarryingVehicles := ASK interchangeYard TO
                                countFlatcarsCarryingVehicles();
    END WHILE;
  END IF;
  ELSIF (ordering = "Containers Only")
  IF (poeMode)
    {Containers only accepted}
    numFlatcarsCarryingContainers := ASK interchangeYard TO
    countFlatcarsCarryingContainers();
    WHILE (numFlatcarsCarryingContainers = 0)
      WAIT FOR trainClassifiedTrigger TO Fire;
      ON INTERRUPT
      END WAIT;
      numFlatcarsCarryingContainers := ASK interchangeYard TO
    countFlatcarsCarryingContainers();
    END WHILE;
  END IF;
  ELSE
    {Both Vehicles/Unit Equipment and Containers are accepted}
  END IF;
  IF (status = "idle")
    {Set status to busy}
    {ASK SELF TO setStatus("busy");} {Removed}
    tempTime := SimTime();

```

```

svcTime := 0.0;
{Build the appropriate stringOfFlatcars to send to spur}
NEW(stringOfFlatcars);
WHILE (stringOfFlatcars.numberIn = 0)
  IF(status = "idle") {"IF" construct added}
    ASK SELF TO
      buildStringOfFlatcars(interchangeYard,stringOfFlatcars);
  END IF;
  IF (stringOfFlatcars.numberIn = 0)
    {OUTPUT("Waiting for more items to be ready in
      emptyRailcarStorage");}
    WAIT DURATION 30.0;
    END WAIT;
  END IF;
END WHILE;
{Let the system know that more items can be put in interchangeYard}
TELL IYTrigger TO Trigger;
ASK SELF TO setStatus("busy"); {Changed status here}
{Obtain a locomotive resource and transit it to the IY}
WAIT FOR port.locomotiveQ TO Give(stringOfFlatcars,1);
  {Processing Time to switch locomotive to IY}
  IF (useProcessDistributionsForTransitTimes)
    lo := switchSpurToIYTime.actualValue -
      switchSpurToIYStdDev.actualValue;
    hi := switchSpurToIYTime.actualValue +
      switchSpurToIYStdDev.actualValue;
    temp := ranGen.UniformReal(lo,hi);
    IF (temp < 0.0)
      temp := 0.0;
    END IF;
  ELSE
    {Calculate the time based upon speed and distance of route}
    {NEED TO ADD THIS}
    {Until this is added, make transit time 0.0}
    temp := 0.0;
  END IF;
  WAIT DURATION temp;
  END WAIT;

  {Processing Time to couple flatcars at IY}
  lo := coupleAtIYTime.actualValue - coupleAtIYStdDev.actualValue;
  hi := coupleAtIYTime.actualValue + coupleAtIYStdDev.actualValue;
  temp := ranGen.UniformReal(lo,hi);
  IF (temp < 0.0)
    temp := 0.0;
  END IF;
  WAIT DURATION temp;
  END WAIT;

  IF (type = "OPEN STAGING") OR (type = "Open Staging") OR
    (type = "OPEN STORAGE") OR (type = "Open Storage")
    {Processing Time to switch flatcars to spur}
    IF (useProcessDistributionsForTransitTimes)
      lo := switchToSpurTime.actualValue -
        switchToSpurStdDev.actualValue;
      hi := switchToSpurTime.actualValue +

```

```

        switchToSpurStdDev.actualValue;
temp := ranGen.UniformReal(lo,hi);
IF (temp < 0.0)
    temp := 0.0;
END IF;
ELSE
    {Calculate the time based upon speed and distance of route}
    {NEED TO ADD THIS}
    {Until this is added, make transit time 0.0}
    temp := 0.0;
END IF;
WAIT DURATION temp;
END WAIT;
ELSIF (type = "COVERED STAGING") OR (type = "Covered Staging")
    {Processing Time to switch flatcars to dock}
    IF (useProcessDistributionsForTransitTimes)
        lo := switchToDockTime.actualValue -
            switchToDockStdDev.actualValue;
        hi := switchToDockTime.actualValue +
            switchToDockStdDev.actualValue;
        temp := ranGen.UniformReal(lo,hi);
        IF (temp < 0.0)
            temp := 0.0;
        END IF;
    ELSE
        {Calculate the time based upon speed and distance of route}
        {NEED TO ADD THIS}
        {Until this is added, make transit time 0.0}
        temp := 0.0;
    END IF;
    WAIT DURATION temp;
    END WAIT;
ELSE
    {type = "APRON" or type = "Apron"}
    {Processing Time to switch flatcars to Berth}
    IF (useProcessDistributionsForTransitTimes)
        lo := switchToBerthTime.actualValue -
            switchToBerthStdDev.actualValue;
        hi := switchToBerthTime.actualValue +
            switchToBerthStdDev.actualValue;
        temp := ranGen.UniformReal(lo,hi);
        IF (temp < 0.0)
            temp := 0.0;
        END IF;
    ELSE
        {Calculate the time based upon speed and distance of route}
        {NEED TO ADD THIS}
        {Until this is added, make transit time 0.0}
        temp := 0.0;
    END IF;
    WAIT DURATION temp;
    END WAIT;
END IF;
{Processing Time to uncouple flatcars at spur}
lo := uncoupleAtSpurTime.actualValue -
    uncoupleAtSpurStdDev.actualValue;
hi := uncoupleAtSpurTime.actualValue +

```

```

        uncoupleAtSpurStdDev.actualValue;
    temp := ranGen.UniformReal(lo,hi);
    IF (temp < 0.0)
        temp := 0.0;
    END IF;
    WAIT DURATION temp;
    END WAIT;
END WAIT;
{Make the locomotive available again}
ASK port.locomotiveQ TO TakeBack(stringOfFlatcars,1);
IF (poeMode)
    {Wait for vehicles/containers to be discharged from flatcars}
    WAIT FOR SELF TO dischargeFlatcarContents(stringOfFlatcars,ranGen);
    END WAIT;
    DISPOSE(stringOfFlatcars);
    {Cycle out empty flatcars}
    {NOTE: This code needs to be added}
END IF;
{Set status "idle" to signify that spur is ready for next string}
{of railcars}
ASK SELF TO setStatus("idle");
{Update operational svcTime parameters}
svcTime := SimTime() - tempTime;
totalBusyTime := totalBusyTime + svcTime;
IF svcTime > maxServiceTime
    maxServiceTime := svcTime;
END IF;

    TELL spurDoneTrigger TO Trigger;
ELSE
    WAIT FOR spurDoneTrigger TO Fire;
    END WAIT;
END IF;
END WHILE;
END METHOD {callForwardFlatcars};

```

Fig. C.2 Modified PORTSIM code for "callForwardFlatcars" method

```

WAITFOR METHOD callForwardCargo (IN stream1 : RandomObj; IN myBerth :
    berthObj; IN vehStagingArea : stagingObj; IN contStagingArea :
    stagingObj; IN palletStagingArea : stagingObj);
{This method calls forward 12 items from the staging area at a time}
{If less than 12 items are available to be call forwarded from the
staging}
{area, then this method will take as many as are available.}
VAR
    cont                : containerObj;
    veh                 : vehicleObj;
    pallet              : palletObj;
    cargo               : ANYOBJ;
    tempveh             : vehicleObj;
    selectedRoute       : routeObj;
    tempreal            : REAL;
    i,j                 : INTEGER;
    cnt                 : INTEGER;
    numItems            : INTEGER; {# of items in staging available to be
                                   loaded}
    contPct,vehPct,palletPct : REAL;
    flag                : BOOLEAN;
    callForwardList     : myQueueObj;
BEGIN
CASE typedesc
    WHEN "Ctnr-NSS":      {Fig. 24}
        {Code before modification
        WHILE (contStagingArea.numContainersReady = 0) OR
            (numContainerHandlers = 0)
            WAIT DURATION 60.0
            END WAIT;
        END WHILE;}
        IF (contStagingArea.numContainersReady = 0) OR {"WHILE" construct}
            (numContainerHandlers = 0)                  {replaced with "IF"}
            WAIT DURATION 60.0
            END WAIT;
        END IF;
        {Checked again after waiting for 60 minutes}
        IF NOT ((contStagingArea.numContainersReady = 0) OR
            (numContainerHandlers = 0))
            numItems := contStagingArea.numContainersReady;
            IF (numItems > numContainerHandlers) OR (numItems > 12)
                IF (numContainerHandlers > 12)
                    numItems := 12;
                ELSE
                    numItems := numContainerHandlers;
                END IF;
            END IF;
        END IF;

        {Obtain the containers that have been dwelled and are ready to
        load,}
        {discharge them from the staging area, and then transit them to
        berth.}
        {Code for sending Containers to Berth}
        END IF;
    WHEN "Ctnr-SS":      {Fig. 24}
        {Code before modification
        WHILE (contStagingArea.numContainersReady = 0) OR

```

```

        (numContainerHandlers = 0)
    WAIT DURATION 60.0
    END WAIT;
END WHILE;}
IF (contStagingArea.numContainersReady = 0) OR {"WHILE" construct}
    (numContainerHandlers = 0)                                {replaced with "IF"}
    WAIT DURATION 60.0
    END WAIT;
END IF;
{Checked again after waiting for 60 minutes}
IF NOT ((contStagingArea.numContainersReady = 0) OR
    (numContainerHandlers = 0))
    numItems := contStagingArea.numContainersReady;
    IF (numItems > numContainerHandlers) OR (numItems > 12)
        IF (numContainerHandlers > 12)
            numItems := 12;
        ELSE
            numItems := numContainerHandlers;
        END IF;
    END IF;
END IF;

{Obtain the containers that have been dwelled and are ready to
load,}
{discharge them from the staging area, and then transit them to
berth.}
{Code for sending Containers to Berth}
END IF;
WHEN "RO/RO":                {Fig. 25}
    {Code before modification}
    WHILE (vehStagingArea.numVehiclesReady = 0) OR (numDrivers = 0)
        WAIT DURATION 60.0
        END WAIT;
    END WHILE;}
    {"WHILE" construct replaced with "IF"}
    IF (vehStagingArea.numVehiclesReady = 0) OR (numDrivers = 0)
        WAIT DURATION 60.0
        END WAIT;
    END IF;
    {Checked again after waiting for 60 units}
    IF NOT ((vehStagingArea.numVehiclesReady = 0) OR (numDrivers = 0))
        i := 0; numItems := 0;
        numItems := vehStagingArea.numVehiclesReady;
        IF (numItems > numDrivers) OR (numItems > 12)
            IF (numDrivers > 12)
                numItems := 12;
            ELSE
                numItems := numDrivers;
            END IF;
        END IF;
    END IF;
    {Obtain the vehicles that have been inspected and dwelled and are
    ready to load,}
    {discharge them from the staging area, and then transit them to
    berth.}
    {Code for sending Vehicles to Berth}
    END IF;
WHEN "Breakbulk":            {Fig. 25}
    {Code before modification}

```

```

WHILE (vehStagingArea.numVehiclesReady = 0) OR (numDrivers = 0)
  WAIT DURATION 60.0
  END WAIT;
END WHILE;}
{"WHILE" construct replaced with "IF"}
IF (vehStagingArea.numVehiclesReady = 0) OR (numDrivers = 0)
  WAIT DURATION 60.0
  END WAIT;
END IF;
{Checked again after waiting for 60 units}
IF NOT ((vehStagingArea.numVehiclesReady = 0) OR (numDrivers = 0))
  numItems := vehStagingArea.numVehiclesReady;
  IF (numItems > numDrivers) OR (numItems > 12)
    IF (numDrivers > 12)
      numItems := 12;
    ELSE
      numItems := numDrivers;
    END IF;
  END IF;
END IF;

{Obtain the vehicles that have been inspected and dwelled and are
 ready to load,}
{discharge them from the staging area, and then transit them to
 berth.}
{Code for sending Vehicles to Berth}
END IF;
OTHERWISE {Fig. 26}
{Need to have a mechanism for discharging both vehicles and
 containers.}
{Use a distribution that will discharge a percentage of vehicles
 equal to}
{the percentage of overall area in the ship it is supposed to load.}
{Code before modification}
WHILE ((vehStagingArea.numVehiclesReady = 0) AND
      (contStagingArea.numContainersReady = 0))
  WAIT DURATION 60.0
  END WAIT;
END WHILE;}
{"WHILE" construct replaced with "IF"}
IF ((vehStagingArea.numVehiclesReady = 0) AND
    (contStagingArea.numContainersReady = 0))
  WAIT DURATION 60.0
  END WAIT;
END IF;

{Checked again after waiting for 60 units}
IF NOT ((vehStagingArea.numVehiclesReady = 0) AND
      (contStagingArea.numContainersReady = 0))
  IF (contStagingArea.numContainersReady = 0)
    {OUTPUT("Trying to send vehicles only");}
    {Send only vehicles}
    {Limit the number of vehicles to call forward by:}
    {1) # of vehicles ready, 2) 12 maximum, 3) # of drivers}
    numItems := vehStagingArea.numVehiclesReady;
    numItems := MINOF(numItems, 12);
    {numItems := MINOF(numItems, numDrivers);}
    WHILE (numItems = 0)

```

```

    WAIT DURATION 60.0;
    END WAIT;
    {Limit the number of vehicles to call forward by:}
    {1) # of vehicles ready, 2) 12 maximum, 3) # of drivers}
    numItems := vehStagingArea.numVehiclesReady;
    numItems := MINOF(numItems,12);
    END WHILE;

    {Obtain the vehicles that have been inspected and dwelled and are
    ready to load,}
    {discharge them from the staging area, and then transit them to
    berth.}
    {Code for sending Vehicles to berth}
    ELSIF (vehStagingArea.numVehiclesReady = 0)
    {OUTPUT("Trying to send containers only");}
    {Limit the number of vehicles to call forward by:}
    {1) # of containers ready, 2) 12 maximum, 3) # of container
    handlers}
    numItems := contStagingArea.numContainersReady;
    numItems := MINOF(numItems,12);
    WHILE (numItems = 0)
    WAIT DURATION 60.0;
    END WAIT;
    numItems := contStagingArea.numContainersReady;
    numItems := MINOF(numItems,12);
    {numItems := MINOF(numItems,numContainerHandlers);}
    END WHILE;

    {Obtain the containers that have been dwelled and are ready to
    load,}
    {discharge them from the staging area, and then transit them to
    berth.}
    {Code for sending Containers to Berth}
    ELSE
    {Code for sending both containers and vehicles to the berth}
    END IF;
    END IF;
    END CASE;
END METHOD {callForwardCargo};

```

Fig. C.3 Modified PORTSIM code for "callForwardCargo"



```

DEFINITION MODULE aanimate;
.....
wrapDialogBoxObj = OBJECT(DialogBoxObj);
  OVERRIDE
    ASK METHOD BeClosed;
  END OBJECT {wrapDialogBoxObj};
.....
END MODULE {aaanimate}

IMPLEMENTATION MODULE aanimate;
.....
OBJECT wrapDialogBoxObj;
  ASK METHOD BeClosed;    {Fig. 27}
  BEGIN
    ASK SELF TO Erase;
  END METHOD;
END OBJECT;

OBJECT XXXXX;
  ASK METHOD XXXXX;
  VAR
    dialogBox      : DialogBoxObj;
    wdialogBox     : wrapDialogBoxObj;
    .....
  BEGIN
    .....
    NEW(wdialogBox);      {Fig. 28}
    ASK wdialogBox TO
      LoadFromLibrary(library, "transportModesErrorDialog");
    ASK portsimMainWin TO AddGraphic(wdialogBox);
    item := ASK wdialogBox TO AcceptInput();
    ASK portsimMainWin TO RemoveThisGraphic(wdialogBox);
    DISPOSE(wdialogBox);
    .....
  END METHOD;
END OBJECT;
.....
END MODULE.

```

Fig. C.4 PORTSIM code for Figures 27 and 28

```

IMPLEMENTATION MODULE aanimate;
.....
OBJECT XXXXX;
  ASK METHOD XXXXX;
  VAR
    dialogBox          : DialogBoxObj;
    wdialogBox2        : wrapDialogBoxObj;
    .....
  BEGIN
    .....
    NEW (wdialogBox2);
    ASK portsimMainWin TO AddGraphic(wdialogBox2);
    REPEAT
      item3 := ASK wdialogBox2 TO AcceptInput();
      IF item3 = NILOBJ
        EXIT;
      END IF;
    UNTIL (ASK item3 ReferenceName = "ok") OR
          (ASK item3 ReferenceName = "cancel");
    IF item3 <> NILOBJ
      IF (ASK item3 ReferenceName = "ok")
        .....
        ASK portsimMainWin TO RemoveThisGraphic(wdialogBox2);
        .....
      ELSE
        {Remove this dialog box}
        ASK portsimMainWin TO RemoveThisGraphic(wdialogBox2);
        DISPOSE(wdialogBox2);
      END IF;
    ELSE
      ASK portsimMainWin TO RemoveThisGraphic(wdialogBox2);
      DISPOSE(wdialogBox2);
    END IF;
    .....
  END METHOD;
END OBJECT;
.....
END MODULE.

```

Fig. C.5 PORTSIM code for Fig. 29

```

IMPLEMENTATION MODULE aanimate;
.....
OBJECT XXXXX;
  ASK METHOD XXXXX;
  VAR
    dialogBox      : DialogBoxObj;
    wdialogBox     : wrapDialogBoxObj;
    .....
  BEGIN
    .....
    NEW(wdialogBox);
    ASK wdialogBox TO LoadFromLibrary(library,"scenarioExistsDialog");
    ASK portsimMainWin TO AddGraphic(wdialogBox);
    ASK wdialogBox TO Draw;
    button := ASK wdialogBox TO AcceptInput();
    IF button <> NILOBJ
      IF ASK button ReferenceName = "cancel"
        flag := FALSE;
      END IF;
    ELSE
      flag := FALSE;
    END IF;
    DISPOSE(wdialogBox);
    .....
  END METHOD;
END OBJECT;
.....
END MODULE.

```

Fig. C.6 PORTSIM code for Fig. 30

```

DEFINITION MODULE aanimate;
.....
shipRemoveDialogBoxObj = OBJECT(DialogBoxObj);
OVERRIDE
  ASK METHOD BeSelected;
  ASK METHOD BeClosed;
END OBJECT {shipRemoveDialogBoxObj};
.....
END MODULE.

```

```

IMPLEMENTATION MODULE aanimate;
.....

OBJECT shipRemoveDialogBoxObj;
  ASK METHOD BeSelected;
  VAR
    .....
  BEGIN
    cnt := ASK LastPicked Id;

```

```

CASE cnt
  WHEN 1: {ok}
    .....
    ASK portsimMainWin TO RemoveThisGraphic(SELF);
    DISPOSE(SELF);
  WHEN .....
  WHEN 2: {cancel}
    ASK portsimMainWin TO RemoveThisGraphic(SELF);
    DISPOSE(SELF);
  OTHERWISE
  END CASE;
END METHOD {BeSelected};

ASK METHOD BeClosed;    {Fig. 31}
BEGIN
  ASK portsimMainWin TO RemoveThisGraphic(SELF);
  DISPOSE(SELF);
END METHOD;
END OBJECT {shipRemoveDialogBoxObj};

OBJECT scenarioDialogBoxObj;
ASK METHOD BeSelected;
VAR
  .....
BEGIN
  CASE ASK LastPicked Id
    WHEN .....
    WHEN 12: {removeShipButton}
      NEW(shipRemoveDialogBox);    {Fig. 32}
      ASK shipRemoveDialogBox TO
        LoadFromLibrary(library,"shipRemoveDialog");
      {Initialize ship list}
      shipTable := ASK SELF Child("shipTable",10);
      shipListBox := ASK shipRemoveDialogBox Child("shipList",5);
      NEW(shipTableItem);
      shipTableItem := ASK shipTable FirstGraphic();
      WHILE shipTableItem <> NILOBJ
        NEW(shipListBoxItem);
        ASK shipListBoxItem TO SetText(shipTableItem.Label);
        ASK shipListBox TO AddGraphic(shipListBoxItem);
        shipTableItem := ASK shipTable NextGraphic(shipTableItem);
      END WHILE;
      ASK portsimMainWin TO AddGraphic(shipRemoveDialogBox);
      ASK shipRemoveDialogBox TO Draw;
    WHEN .....
  END CASE;
END METHOD;
END OBJECT {scenarioDialogBoxObj};
.....
END MODULE.

```

Fig. C.7 PORTSIM code for Figures 31 and 32

```

DEFINITION MODULE aanimate;
.....
gateParamsDialogBoxObj = OBJECT(DialogBoxObj);
  ASK METHOD updateGateParamsDialog(IN Num : INTEGER);
  OVERRIDE
    ASK METHOD BeSelected;
    ASK METHOD BeClosed;
  END OBJECT {gateParamsDialogBoxObj};
.....
END MODULE.

IMPLEMENTATION MODULE aanimate;
.....
OBJECT gateParamsDialogBoxObj;
  ASK METHOD updateGateParamsDialog(IN Num : INTEGER);
  VAR
    .....
  BEGIN
    .....
    ASK SELF TO Update;
    REPEAT
      item := ASK SELF TO AcceptInput();
      IF item = NILOBJ
        EXIT;
      END IF;
    UNTIL (ASK item ReferenceName = "cancel");
    IF (ASK portsimMainWin IncludesGraphic(SELF))
      ASK portsimMainWin TO RemoveThisGraphic(SELF);
    END IF;
  END METHOD;

  ASK METHOD BeClosed;
  BEGIN
    ASK SELF TO Erase;
  END METHOD;
END OBJECT {gateParamsDialogBoxObj};
.....
END MODULE.

```

Fig. C.8 PORTSIM code for Fig. 33

```
DEFINITION MODULE aanimate;
```

```
.....
stagingParamsDialogBoxObj = OBJECT(DialogBoxObj);
ASK METHOD updateStagingParamsDialog(IN Num : INTEGER);
OVERRIDE
  ASK METHOD BeSelected;
  ASK METHOD BeClosed;
END OBJECT {stagingParamsDialogBoxObj};
```

```
.....
END MODULE.
```

```
IMPLEMENTATION MODULE aanimate;
```

```
.....
OBJECT stagingParamsDialogBoxObj;
ASK METHOD updateStagingParamsDialog(IN Num : INTEGER);
VAR
  .....
BEGIN
  .....
  REPEAT
    item := ASK SELF TO AcceptInput();
    IF item = NILOBJ
      {Ensure at least one staging area is available before allowing
       user to exit}
      IF (availStagingAreas.numberIn < 0)
        NEW(wdialogBox);
        ASK wdialogBox TO
          LoadFromLibrary(library, "selectStagingErrorDialog");
        ASK portsimMainWin TO AddGraphic(wdialogBox);
        item := ASK wdialogBox TO AcceptInput();
        ASK portsimMainWin TO RemoveThisGraphic(wdialogBox);
        DISPOSE(wdialogBox);
      ELSE
        IF (ASK portsimMainWin IncludesGraphic(SELF))
          ASK portsimMainWin TO RemoveThisGraphic(SELF);
        END IF;
      END IF;
    END IF;
    EXIT;
  END IF;
  UNTIL (ASK item ReferenceName = "cancel");
END METHOD {updateStagingParamsDialog};
```

```
ASK METHOD BeSelected;
```

```
VAR
```

```
.....
BEGIN
  CASE name
    WHEN "ok":
      .....
    WHEN "cancel":
      {Ensure at least one staging area is available before allowing
       user to exit}
      IF (availStagingAreas.numberIn < 0)
        NEW(wdialogBox);
        ASK wdialogBox TO
          LoadFromLibrary(library, "selectStagingErrorDialog");
        ASK portsimMainWin TO AddGraphic(wdialogBox);
```

```

        item := ASK wdialogBox TO AcceptInput();
        ASK portsimMainWin TO RemoveThisGraphic(wdialogBox);
        DISPOSE(wdialogBox);
    ELSE
        IF (ASK portsimMainWin IncludesGraphic(SELF))
            ASK portsimMainWin TO RemoveThisGraphic(SELF);
        END IF;
    END IF;
    WHEN .....
    END CASE;
END METHOD {BeSelected};

ASK METHOD BeClosed;
BEGIN
    ASK SELF TO Erase;
END METHOD;
END OBJECT {stagingParamsDialogBoxObj};
.....
END MODULE.

```

Fig. C.9 PORTSIM code for Fig. 34

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